## Lecture 3

Tuesday 03-02-2008
14.00-17.30

## Linear Time-Invariant Systems, Part 2

Lecturer: Paolo Rapisarda

1. Part I:

- Inputs and outputs;
- Autonomous behaviors;
- Input-output representations.


## 2. Part II:

- Controllability;
- Image representations;
- Complementability and decomposition of behaviors;
- Observability.


## Inputs and outputs

Recall that $P\left(\frac{d}{d t}\right): \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{W}}\right) \rightarrow \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{g}}\right)$ is surjective $\Leftrightarrow$ $P(\xi)$ has full row rank as a polynomial matrix

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Equivalently: $P(\lambda)$ has full column rank for all $\lambda \in \mathbb{C}$

## Reprise: surjectivity and injectivity of differential operators

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Equivalently, $P$ admits a left inverse on $\mathscr{C}^{\infty}(\mathbb{R})$ :

$$
P=U\left[\begin{array}{c}
I_{\mathrm{m}} \\
0
\end{array}\right] V \Longrightarrow V^{-1}\left[\begin{array}{ll}
I_{\mathrm{m}} & 0
\end{array}\right] U^{-1} \text { is left inverse }
$$

## Free variables

Given $\mathscr{B} \in \mathscr{L}^{\mathrm{w}}$ and $I:=\left\{i_{1}, \ldots, i_{\mathrm{k}}\right\} \subseteq\{1, \ldots, \mathrm{w}\}$, let

$$
\begin{aligned}
\Pi_{I} \mathscr{B}:=\quad & \left\{\left(\hat{w}_{i_{1}}, \ldots, \hat{w}_{i_{\mathrm{k}}}\right) \in \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{k}}\right) \mid \exists w \in \mathscr{B}\right. \\
& \text { s.t. } \left.w=\left(w_{1}, \ldots, \hat{w}_{i_{1}}, \ldots, \hat{w}_{i_{\mathrm{k}}}, \ldots, w_{\mathrm{w}}\right)\right\}
\end{aligned}
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projection of $\mathscr{B}$ onto variables $w_{i_{j}}, j=1, \ldots, \mathrm{k}$

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& \text { s.t. } \left.w=\left(w_{1}, \ldots, w_{i_{1}}, \ldots, w_{i_{\mathrm{k}}}, \ldots, w_{\mathrm{w}}\right) \in \mathscr{B}\right\}
\end{aligned}
$$

projection of $\mathscr{B}$ onto variables $w_{i_{j}}, j=1, \ldots, \mathrm{k}$
Variables $w_{i_{j}}, j=1, \ldots, \mathrm{k}$ are free if

$$
\Pi_{I} \mathscr{B}=\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{k}}\right)
$$

## Free variables

Example:

$$
p_{1}\left(\frac{d}{d t}\right) w_{1}+p_{2}\left(\frac{d}{d t}\right) w_{2}+p_{3}\left(\frac{d}{d t}\right) w_{3}=0
$$

Assume $p_{i} \neq 0, i=1,2,3$.
Let $I=\{1\}$; since $\left[p_{2}(\xi) \quad p_{3}(\xi)\right]$ is full row rank, for every $w_{1} \in \mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R})$ there exist $w_{2}, w_{3}$ satisfying equation.
$w_{1}$ is free.

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$w_{1}$ is free.
$\left(w_{1}, w_{2}\right)$ (and $\left(w_{2}, w_{3}\right)$, and $\left.\left(w_{1}, w_{3}\right)\right)$ are also free.

Example: In $\frac{d}{d t} x=A x+B u$, the variable $u$ is free.

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## Behavior is

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\mathscr{B}=\operatorname{ker}\left[\begin{array}{ll}
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\end{array}\right] \leadsto R(\xi)=\left[\begin{array}{ll}
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Let $I=\left\{i_{1}, \ldots, i_{\mathrm{k}}\right\} \subseteq\{1, \ldots, \mathrm{w}\}$. The variables $w_{i_{1}}, \ldots, w_{i_{\mathrm{k}}}$ form a maximally free set if
they are free; and
for every $I^{\prime}=\left\{i_{1}^{\prime}, \ldots, i_{\mathrm{k}}^{\prime}\right\} \subset\left\{1, \ldots\right.$, w such that $I \subset I^{\prime}$ it holds

$$
\Pi_{I^{\prime}} \mathscr{B} \subset \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\left|I^{\prime}\right|}\right)
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## Maximally free sets

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\Pi_{I^{\prime}} \mathscr{B} \subsetneq \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\left|I^{\prime}\right|}\right)
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Maximally free set: every variable in it is free, but any additional variable is not

## Maximally free sets

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Maximally free sets are nonunique!

## Inputs and outputs

Theorem: Let $\mathscr{B} \in \mathscr{L}^{\text {w }}$. Assume (if necessary, after permutation of the variables) $w$ partitioned as

$$
w=\left[\begin{array}{l}
w_{1} \\
w_{2}
\end{array}\right]
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with $w_{1}$ maximally free. Then $w_{1}$ are inputs and $w_{2}$ outputs.

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Example: for $p_{1}\left(\frac{d}{d t}\right) w_{1}+p_{2}\left(\frac{d}{d t}\right) w_{2}+p_{3}\left(\frac{d}{d t}\right) w_{3}=0$ and assuming $p_{i} \neq 0$ for $i=1,2,3$, we can choose

- $\left\{w_{1}, w_{2}\right\}$ or
- $\left\{w_{2}, w_{3}\right\}$ or
- $\left\{w_{1}, w_{3}\right\}$
as inputs.

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Remarks
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Nonunicity of i/o partition is not an issue.

## Remarks

Nonunicity of i/o partition is not an issue. Consider (linear) resistors:

$$
\mathscr{B}=\{(V, I) \mid V=R \cdot I\}
$$

Is it voltage- or current-controlled? Consider


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Nonunicity of i/o partition is not an issue.
'Causality' an issue? What about

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w_{1}=\frac{d}{d t} w_{2} ?
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Don't $w_{1}$ and $w_{2}$ 'happen' at the same time?

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Don't $w_{1}$ and $w_{2}$ 'happen' at the same time?
'Smoothness', meaning

$$
\begin{aligned}
(u, y) \in \mathscr{B} & \text { and } \quad u k \text {-times differentiable } \\
& \Longrightarrow \quad y k \text {-times differentiable }
\end{aligned}
$$

if and only if $P^{-1} Q$ is proper. Strict properness $\Leftrightarrow$

$$
\begin{aligned}
(u, y) \in \mathscr{B} & \text { and } \\
& u k \text {-times differentiable } \\
& \Longrightarrow y(k+1) \text {-times differentiable }
\end{aligned}
$$

## Causality in discrete-time systems

Consider a linear $\mathscr{B} \subset\left(\mathbb{R}^{\mathrm{w}_{1}+\mathrm{w}_{2}}\right)^{\mathbb{Z}}$. Let $\mathscr{B}_{1}=\Pi_{\mathrm{w}_{1}} \mathscr{B}$.
$w_{2}$ does not anticipate $w_{1} \Leftrightarrow$

$$
w_{1} \in \mathscr{B}_{1} \quad \text { and } \quad w_{1 \mid \mathbb{Z}_{-}}=0
$$

$\Longrightarrow \quad$ exists $w_{2}^{\prime}$ s.t. $w_{2 \mid \mathbb{Z}_{-}}=0$ and $\left(w_{1}, w_{2}\right) \in \mathscr{B}$

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\end{array}
$$

Theorem: Assume $w_{2}$ is output and $w_{1}$ is input, and let

$$
P(\sigma) w_{2}=Q(\sigma) w_{1}
$$

be an i/o representation of $\mathscr{B}$. Then $w_{2}$ does not anticipate $w_{1}$ $\Leftrightarrow P^{-1} Q$ is proper.

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with $P$ square and nonsingular. Then $y$ is output and $u$ is input.

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with $P$ square and nonsingular. Then $y$ is output and $u$ is input.

Surjectivity of $P\left(\frac{d}{d t}\right) \Longrightarrow u$ is free.
$u$ maximally free: add one component of $y$ to those of $u$, resulting set satisfies differential equation $\Longrightarrow$ it is not free.

## Input-output representations

Theorem: Let $\mathscr{B} \in \mathscr{L}^{w}$. There exists (possibly after permuting components) a partition of $w=(u, y)$ and $P \in \mathbb{R}^{\mathrm{y} \times \mathrm{y}}[\xi]$ nonsingular, $Q \in \mathbb{R}^{\mathrm{y} \times \mathrm{u}}[\xi]$ such that

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The partition can be chosen so that $P^{-1} Q$ is proper.
Proof: Use minimal kernel representation $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$. $R$ of full row rank $\Longrightarrow$ exists nonsingular submatrix $P$.
For $P^{-1} Q$ proper, select $P$ to be a maximal determinantal degree (nonsingular) submatrix of $R$.


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R(\xi)=\left[\begin{array}{ccc}
m_{1} \xi^{2}+c_{1} \xi+k_{1} & -c_{1} \xi-k_{1} & -1 \\
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$w_{1}$ and $w_{2}$ outputs, $F$ input; $P^{-1} Q$ strictly proper

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$w_{1}$ and $F$ outputs, $w_{2}$ input; $P^{-1} Q$ not proper

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$\mathrm{p}(\mathscr{B})$ equals $\operatorname{rank}(R)$ for every $R$ such that $\operatorname{ker} R\left(\frac{d}{d t}\right)=\mathscr{B}$;

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Number of inputs fixed, input cardinality $\mathrm{m}(\mathscr{B})$;

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Number of inputs fixed, input cardinality $\mathrm{m}(\mathscr{B})$;
$\mathrm{m}(\mathscr{B})$ equals $\mathrm{w}(\mathscr{B})-\operatorname{rank}(R)$ for every $R$ such that ker $R\left(\frac{d}{d t}\right)=\mathscr{B}$.

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$\mathrm{m}(\mathscr{B})$ equals $\mathrm{w}(\mathscr{B})-\operatorname{rank}(R)$ for every $R$ such that ker $R\left(\frac{d}{d t}\right)=\mathscr{B}$.

In discrete-time case, there always exists a causal input-output partition!

## Autonomous behaviors

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No inputs: autonomous systems
```

Recall that $\mathscr{B}$ is autonomous if

$$
\begin{array}{rll}
w_{1}, w_{2} \in \mathscr{B} & \text { and } & \left.w_{1}\right|_{(-\infty, 0]}=\left.w_{2}\right|_{(-\infty, 0]} \\
& \Longrightarrow & w_{1}=w_{2}
\end{array}
$$

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## Equivalent with

$\mathrm{m}(\mathscr{B})=0$ (no inputs);
there exists $R \in \mathbb{R}^{\mathrm{w} \times \mathrm{w}}[\xi]$ nonsingular such that
$\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$

## Example: a mechanical system

$$
m_{c_{1}}^{m_{1}}
$$

## Example: a mechanical system



## Classical mechanics:

 motion depends only on 'initial conditions'
## Example: a mechanical system

$$
m_{c_{1}}^{m_{1}}
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\end{array}\right]
$$

$R$ nonsingular $\sim$ autonomous system

## Example: state-space systems

Let $(A, C)$ observable and consider

$$
\mathscr{B}:=\left\{y \in \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{y}}\right) \mid \exists x \text { s.t. } \frac{d}{d t} x=A x, y=C x\right\}
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$\mathscr{B}$ is autonomous: there are no free variables in $y$.

For autonomous $\mathscr{B}, \operatorname{det}(R)$ is invariant for all $R$ such that $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$.

## Remarks

For autonomous $\mathscr{B}, \operatorname{det}(R)$ is invariant for all $R$ such that $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$.

Theorem: Let $\mathscr{B} \in \mathscr{L}^{\text {w }}$ be autonomous. Then $\mathscr{B}$ is a finite-dimensional subspace of $\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{w}\right)$.

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Theorem: Let $\mathscr{B} \in \mathscr{L}^{\mathrm{w}}$ be autonomous. Then $\mathscr{B}$ is a finite-dimensional subspace of $\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{W}\right)$.
Proof: Take $R$ s.t. $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$, w.l.o.g. minimal. Compute Smith form $R=U \Delta V$ :

$$
R\left(\frac{d}{d t}\right) w=0 \Longleftrightarrow \Delta\left(\frac{d}{d t}\right) \underbrace{V\left(\frac{d}{d t}\right) w}_{=: w^{\prime}}=0
$$

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Theorem: Let $\mathscr{B} \in \mathscr{L}^{\text {w }}$ be autonomous. Then $\mathscr{B}$ is a finite-dimensional subspace of $\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{W}\right)$.
Proof: Now $\Delta\left(\frac{d}{d t}\right) \underbrace{V\left(\frac{d}{d t}\right) w}_{=: w^{\prime}}=0$ implies

$$
w^{\prime}=\operatorname{col}\left(w_{i}^{\prime}\right)_{i=1, \ldots, \mathrm{~W}} \in \operatorname{ker} \Delta\left(\frac{d}{d t}\right) \Leftrightarrow w_{i}^{\prime} \in \operatorname{ker} \delta_{i}\left(\frac{d}{d t}\right)
$$

with $\delta_{i}$ the $i$-th invariant polynomial. Scalar case.

For autonomous $\mathscr{B}, \operatorname{det}(R)$ is invariant for all $R$ such that $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$.

Theorem: Let $\mathscr{B} \in \mathscr{L}^{\mathrm{w}}$ be autonomous. Then $\mathscr{B}$ is a finite-dimensional subspace of $\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{W}\right)$.

Proof: Now $\Delta\left(\frac{d}{d t}\right) \underbrace{V\left(\frac{d}{d t}\right) w}_{=: w^{\prime}}=0$ implies

$$
w^{\prime}=\operatorname{col}\left(w_{i}^{\prime}\right)_{i=1, \ldots, \mathrm{w}} \in \operatorname{ker} \Delta\left(\frac{d}{d t}\right) \Leftrightarrow w_{i}^{\prime} \in \operatorname{ker} \delta_{i}\left(\frac{d}{d t}\right)
$$

with $\delta_{i}$ the $i$-th invariant polynomial. Scalar case. Set of solutions of linear differential equation is finite-dimensional. Also $w$ is!

## On autonomous system trajectories

## Scalar case:

$$
p\left(\frac{d}{d t}\right) w=0 \Longleftrightarrow w(t)=\sum_{i=1}^{n} \sum_{j=0}^{n_{i}} \alpha_{i j} t^{j} e^{\lambda_{i} t}
$$

where

- $n$ is number of distinct roots of $p(\xi)$;
- $\quad \lambda_{i}$ is $i$-th root of $p(\xi)$;
$n_{i}$ multiplicity of $\lambda_{i} ;$
$\alpha_{i j} \in \mathbb{C}$.


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- $\quad \lambda_{i}$ is $i$-th root of $p(\xi)$;
- $n_{i}$ multiplicity of $\lambda_{i}$;
- $\quad \alpha_{i j} \in \mathbb{C}$.
$\lambda_{i}$ are the characteristic frequencies of $p$.


## On autonomous system trajectories

For $\mathrm{w}>1$, resort to Smith form $R=U \Delta V$ :

$$
R\left(\frac{d}{d t}\right) w=0 \Longleftrightarrow \Delta\left(\frac{d}{d t}\right) \underbrace{V\left(\frac{d}{d t}\right) w}_{=: w^{\prime}}=0
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w^{\prime}=\operatorname{col}\left(w_{i}^{\prime}\right)_{i=1, \ldots, \mathrm{w}} \in \operatorname{ker} \Delta\left(\frac{d}{d t}\right) \Leftrightarrow w_{i}^{\prime} \in \operatorname{ker} \delta_{i}\left(\frac{d}{d t}\right)
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\end{gathered}
$$

with $\delta_{i}$ the $i$-th invariant polynomial. Scalar case!
Assume for simplicity all roots of $\operatorname{det}(R)$ are simple:

$$
w=V\left(\frac{d}{d t}\right)^{-1} w^{\prime} \Longleftrightarrow w(t)=\sum_{i=1}^{n} \alpha_{i} e^{\lambda_{i} t}
$$

with $\alpha_{i} \in \mathbb{C}^{\mathrm{W}}$ such that $R\left(\lambda_{i}\right) \alpha_{i}=0, i=1, \ldots, n$.

## Remarks

Linear combinations of polynomial exponential vector trajectories

$$
\sum_{i=1}^{n} \sum_{j=0}^{n_{i}} \alpha_{i j} t^{j} e^{\lambda_{i} t}
$$

with $\alpha_{i j} \in \mathbb{C}^{\mathrm{w}}$.

## Remarks

Linear combinations of polynomial exponential vector trajectories

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\sum_{i=1}^{n} \sum_{j=0}^{n_{i}} \alpha_{i j} t^{j} e^{\lambda_{i} t}
$$

with $\alpha_{i j} \in \mathbb{C}^{\mathrm{w}}$.

Characteristic frequencies $\lambda_{i}$ are roots of $\operatorname{det}(R)$. Together with corresponding multiplicities, they determine $\mathscr{B}$ uniquely.
$\mathscr{B} \in \mathscr{L}^{\mathrm{w}}$ is asymptotically stable $\Leftrightarrow$

$$
w \in \mathscr{B} \Longrightarrow \lim _{t} \rightarrow \infty w(t)=0
$$

Note: asymptotic stability implies autonomy.

## Stability

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Note: asymptotic stability implies autonomy.
Theorem: $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$ is asymptotically stable $\Leftrightarrow \boldsymbol{\operatorname { r a n k }}(R(\lambda))=\mathrm{w}(\mathscr{B})$ for all $\lambda \in \mathbb{C}$ s.t. $\operatorname{Re}(\lambda) \geq 0$.
$\mathscr{B} \in \mathscr{L}^{W}$ is stable $\Leftrightarrow$

$$
w \in \mathscr{B} \Longrightarrow w_{\mathbb{R}_{+}} \text {is bounded. }
$$

Note: stability implies autonomy.
Theorem: $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$ is stable $\Leftrightarrow$

1. $\boldsymbol{\operatorname { r a n k }}(R(\lambda))=\mathrm{w}(\mathscr{B})$ for all $\lambda \in \mathbb{C}$ s.t. $\operatorname{Re}(\lambda)>0$;
2. For all $\omega \in \mathbb{R}, w(\mathscr{B})-\operatorname{rank}(R(i \omega))$ equals the multiplicity of $i \omega$ as a root of $\operatorname{det}(R)$.

## Stability

$\mathscr{B} \in \mathscr{L}^{W}$ is stable $\Leftrightarrow$

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w \in \mathscr{B} \Longrightarrow w_{\mathbb{R}_{+}} \text {is bounded. }
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Note: stability implies autonomy.

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& \text { Theorem: } \mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right) \text { is stable } \Leftrightarrow \\
& \text { 1. } \operatorname{rank}(R(\lambda))=w(\mathscr{B}) \text { for all } \lambda \in \mathbb{C} \text { s.t. } \operatorname{Re}(\lambda)>0 \text {; } \\
& \text { 2. For all } \omega \in \mathbb{R}, w(\mathscr{B})-\operatorname{rank}(R(i \omega)) \text { equals the } \\
& \text { multiplicity of } i \omega \text { as a root of } \operatorname{det}(R) \text {. }
\end{aligned}
$$

Stability=roots in closed left half-plane, and semisimplicity.

End of Part I

## Controllability

## Controllability

$\mathscr{B}$ controllable $\Leftrightarrow$ for all $w_{1}, w_{2} \in \mathscr{B}$ there exists $w \in \mathscr{B}$ and $T \geq 0$ such that

$$
w(t)=\left\{\begin{array}{lll}
w_{1}(t) & \text { for } & t<0 \\
w_{2}(t) & \text { for } & t \geq T
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Past of any trajectory can be "patched up" with future of any trajectory

## Examples

$$
r\left(\frac{d}{d t}\right) w=0
$$

where $0 \neq r \in \mathbb{R}[\xi]$ has degree $n$.
System autonomous: every solution uniquely determined by 'initial conditions' $\frac{d^{i} w}{d t^{i}}(t), i=0, \ldots, n-1$, so no patching possible among $d \Leftrightarrow$ erent trajectories.

Past of trajectory uniquely determines its future.

```
Examples
```

Classical state-space system

$$
\begin{aligned}
\frac{d}{d t} x & =A x+B u \\
y & =C x+D u
\end{aligned}
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\mathscr{B}_{s}:=\left\{(u, y, x) \mid \text { s.t. } \frac{d}{d t} x=A x+B u, y=C x+D u\right\} \\
\mathscr{B}:=\left\{(u, y) \mid \exists x \text { s.t. } \frac{d}{d t} x=A x+B u, y=C x+D u\right\} \\
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$\mathscr{B}_{s}$ controllable $\Leftrightarrow \mathscr{B}_{x}$ controllable $\Longrightarrow \mathscr{B}$ controllable.

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$\mathscr{B}_{s}$ controllable $\Leftrightarrow \mathscr{B}_{x}$ controllable $\Longrightarrow \mathscr{B}$ controllable. If $x$ minimal, then $\mathscr{B}$ controllable $\Longrightarrow \mathscr{B}_{s}$ controllable.

## Examples

## Classical state-space system

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'State point-controllability": for all $x_{1}, x_{2} \in \mathbb{R}^{\mathrm{n}} \exists x \in \mathscr{B}_{x}$ and $T \geq 0$ s.t. $x(0)=x_{0}$ and $x(T)=x_{1}$.

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If $x$ minimal, then $\mathscr{B}$ controllable $\Leftrightarrow \mathscr{B}_{s}$ controllable $\Longleftrightarrow \mathscr{B}_{s}$ state point-controllable.

Theorem: $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$ is controllable $\Leftrightarrow$
$\operatorname{rank}(R(\lambda))$ is constant for all $\lambda \in \mathbb{C}$

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Proof: Compute Smith form

$$
R=U\left[\begin{array}{ll}
\Delta & 0 \\
0 & 0
\end{array}\right] V \in \mathbb{R}^{\mathrm{p} \times \mathrm{w}}[\xi]
$$

$U\left(\frac{d}{d t}\right), V\left(\frac{d}{d t}\right)$ bijective $\Longrightarrow \operatorname{ker} R\left(\frac{d}{d t}\right)$ controllable $\Leftrightarrow$ ker $\Delta\left(\frac{d}{d t}\right)$ is.

## Algebraic characterization of controllability

Theorem: $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$ is controllable $\operatorname{rank}(R(\lambda))$ is constant for all $\lambda \in \mathbb{C}$

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Change variables $w \leadsto w^{\prime}:=V\left(\frac{d}{d t}\right) w$, define $\mathscr{B}^{\prime}:=V\left(\frac{d}{d t}\right) \mathscr{B}=\operatorname{ker} \Delta\left(\frac{d}{d t}\right)$.

## Algebraic characterization of controllability

Theorem: $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$ is controllable
$\operatorname{rank}(R(\lambda))$ is constant for all $\lambda \in \mathbb{C}$
Proof: Last $\mathrm{p}-\boldsymbol{\operatorname { r a n k }}(R)$ trajectories of $\mathscr{B}^{\prime}=\operatorname{ker} \Delta\left(\frac{d}{d t}\right)$ are free, since equations are $0 \cdot w_{i}^{\prime}=0$.
First $\operatorname{rank}(R)$ equations are

$$
\delta_{i}\left(\frac{d}{d t}\right) w_{i}^{\prime}=0
$$

with $\delta_{i} i$-th invariant polynomial of $R$.
Evidently, $w_{i}^{\prime}$ controllable if and only if $\delta_{i}=1$.


Case 1: $C R_{C} \neq \frac{L}{R_{L}}$

$$
\begin{aligned}
\left(\frac{R_{C}}{R_{L}}\right. & \left.+\left(1+\frac{R_{C}}{R_{L}}\right) C R_{C} \frac{d}{d t}+C R_{C} \frac{L}{R_{L}} \frac{d^{2}}{d t^{2}}\right) V \\
& =\left(1+C R_{C} \frac{d}{d t}\right)\left(1+\frac{L}{R_{L}} \frac{d}{d t}\right) R_{C} I
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¿Is system controllable?

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Are there common roots among the two polynomials?

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Are there common roots among the two polynomials?
No $\Longrightarrow$ system is controllable

## Example



Case 2: $C R_{C}=\frac{L}{R_{L}}$

$$
\left(\frac{R_{C}}{R_{L}}+C R_{C} \frac{d}{d t}\right) V=\left(1+C R_{C} \frac{d}{d t}\right) R_{C} I
$$

## Example



Case 2: $C R_{C}=\frac{L}{R_{L}}$

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¿Is system controllable?

## Example



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\left(\frac{R_{C}}{R_{L}}+C R_{C} \frac{d}{d t}\right) V=\left(1+C R_{C} \frac{d}{d t}\right) R_{C} I
$$

¿Is system controllable?

$$
\left[\begin{array}{ll}
\frac{R_{C}}{R_{L}}+C R_{C} \xi & -\left(1+C R_{C} \xi\right) R_{C}
\end{array}\right]
$$

Are there common roots among the two polynomials?

## Example



## Case 2: $C R_{C}=\frac{L}{R_{L}}$

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\left(\frac{R_{C}}{R_{L}}+C R_{C} \frac{d}{d t}\right) V=\left(1+C R_{C} \frac{d}{d t}\right) R_{C} I
$$

¿Is system controllable?
If $R_{C}=R_{L} \Longrightarrow$ system is not controllable

## Remarks

$\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$, with $R \in \mathbb{R}^{\mathrm{w} \times \mathrm{w}}[\xi]$ nonsingular, is controllable $\Longleftrightarrow R$ is unimodular $\Longleftrightarrow \mathscr{B}=\{0\}$

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$\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$, with $R \in \mathbb{R}^{\mathrm{w} \times \mathrm{W}}[\xi]$ nonsingular, is controllable $\Longleftrightarrow R$ is unimodular $\Longleftrightarrow \mathscr{B}=\{0\}$

Rank constancy test generalization of 'Hautus test' for state-space systems.

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Rank constancy test generalization of 'Hautus test' for state-space systems.

Trajectory-, not representation -based definition as in state-space framework.

## Image representations

## Image representations and controllability

Theorem: There exists $M \in \mathbb{R}^{w \times \bullet}[\xi]$ such that $\mathscr{B}=\operatorname{im} M\left(\frac{d}{d t}\right) \Leftrightarrow \mathscr{B}$ is controllable.

## Image representations and controllability

## Theorem: There exists $M \in \mathbb{R}^{w \times \bullet}[\xi]$ such that

 $\mathscr{B}=\operatorname{im} M\left(\frac{d}{d t}\right) \Leftrightarrow \mathscr{B}$ is controllable.Only if: Full behavior is controllable, since has kernel representation induced by

$$
\left[\begin{array}{ll}
I_{\mathrm{w}} & -M(\xi)
\end{array}\right]
$$

with constant rank over $\mathbb{C}$.

## Image representations and controllability

## Theorem: There exists $M \in \mathbb{R}^{\mathrm{w} \times} \cdot[\xi]$ such that

 $\mathscr{B}=\operatorname{im} M\left(\frac{d}{d t}\right) \Leftrightarrow \mathscr{B}$ is controllable.If: Take $R$ for minimal kernel representation of $\mathscr{B}$. Apply constancy of rank to conclude Smith form of $R$ is

$$
R=U\left[\begin{array}{ll}
I_{\mathrm{p}} & 0_{\mathrm{p} \times \mathrm{m}}
\end{array}\right] V
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## Image representations and controllability

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$R=U\left[\begin{array}{ll}I_{\mathrm{p}} & 0_{\mathrm{p} \times \mathrm{m}}\end{array}\right] V$.
Now $U\left(\frac{d}{d t}\right)\left[\begin{array}{ll}I_{\mathrm{p}} & 0_{\mathrm{p} \times \mathrm{m}}\end{array}\right] \underbrace{V\left(\frac{d}{d t}\right) w}_{=: w^{\prime}}=0 \Leftrightarrow\left[\begin{array}{ll}I_{\mathrm{p}} & 0_{\mathrm{p} \times \mathrm{m}}\end{array}\right] w^{\prime}=0 \Leftrightarrow$

$$
w^{\prime}=\left[\begin{array}{c}
c_{\mathrm{p}} \\
I_{\mathrm{m}}
\end{array}\right] \ell
$$

with $\ell \in \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{m}\right)$ free.

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Theorem: There exists $M \in \mathbb{R}^{\mathrm{w} \times} \cdot[\xi]$ such that $\mathscr{B}=\operatorname{im} M\left(\frac{d}{d t}\right) \Leftrightarrow \mathscr{B}$ is controllable.

Consequently,

$$
w^{\prime}=V\left(\frac{d}{d t}\right) w=\left[\begin{array}{c}
0_{\mathrm{p}} \\
I_{\mathrm{m}}
\end{array}\right] \ell
$$

from which

$$
w=V\left(\frac{d}{d t}\right)^{-1}\left[\begin{array}{c}
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I_{\mathrm{m}}
\end{array}\right] \ell=: M\left(\frac{d}{d t}\right) \ell
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$$

Note also that $M$ can be chosen with $\mathrm{m}(B)$ columns.

# Complementability 

 anddecomposition of behaviors

## Complementability

Theorem: Let $\mathscr{B} \in \mathscr{L}^{\text {w }}$ be controllable. There exists $\mathscr{B}^{\prime} \in \mathscr{L}^{\text {w }}$ such that

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\mathscr{B} \oplus \mathscr{B}^{\prime}=\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{W}\right)
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Proof: Let $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$ be a minimal kernel representation. $\mathscr{B}$ controllable $\Leftrightarrow$ Smith form of $R$ is

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R=U\left[\begin{array}{ll}
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Define

$$
R^{\prime}:=U\left[\begin{array}{ll}
0 & I_{\mathrm{w}-\mathrm{p}}
\end{array}\right] V
$$

and $\mathscr{B}^{\prime}:=\operatorname{ker} R^{\prime}\left(\frac{d}{d t}\right) . \mathscr{B}^{\prime}$ is also controllable.

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$$
\mathscr{B} \oplus \mathscr{B}^{\prime}=\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{w}}\right)
$$

Proof: Observe that $\mathscr{B} \cap \mathscr{B}^{\prime}$ is represented in kernel form by

$$
U\left[\begin{array}{cc}
I_{\mathrm{p}} & 0 \\
0 & I_{\mathrm{w}-\mathrm{p}}
\end{array}\right] V
$$

a unimodular matrix. Consequently, $\mathscr{B} \cap \mathscr{B}^{\prime}=\{0\}$.

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\mathscr{B} \oplus \mathscr{B}^{\prime}=\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathbf{W}}\right)
$$

Proof: Easy to see image representations of $\mathscr{B}, \mathscr{B}^{\prime}$ given by

$$
\mathscr{B}=\operatorname{im} V^{-1}\left[\begin{array}{c}
0 \\
I_{\mathrm{w}-\mathrm{p}}
\end{array}\right] \quad \mathscr{B}^{\prime}=\operatorname{im} V^{-1}\left[\begin{array}{c}
I_{\mathrm{p}} \\
0
\end{array}\right]
$$

Consequently $\mathscr{B}+\mathscr{B}^{\prime}$ represented by

$$
V^{-1}\left[\begin{array}{cc}
0 & I_{\mathrm{p}} \\
I_{\mathrm{w}-\mathrm{p}} & 0
\end{array}\right]
$$

unimodular, consequently bijective.

## Decomposition of behaviors

Theorem: Let $\mathscr{B}=\operatorname{ker} R\left(\frac{d}{d t}\right)$, with $R \in \mathbb{R}^{\mathrm{p} \times \mathrm{w}}[\xi]$ full row rank. There exist $\mathscr{B}_{\text {aut }} \subseteq \mathscr{B}$ and $\mathscr{B}_{\text {contr }} \subseteq \mathscr{B}$ such that

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Proof: Write Smith form of $R=U\left[\begin{array}{ll}D & 0_{\mathrm{p} \times(\mathrm{w}-\mathrm{p})}\end{array}\right] V$, define $\mathscr{B}^{\prime}:=V\left(\frac{d}{d t}\right) \mathscr{B}$.

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Proof: Write Smith form of $R=U\left[\begin{array}{ll}D & 0_{\mathrm{p} \times(\mathrm{w}-\mathrm{p})}\end{array}\right] V$, define $\mathscr{B}^{\prime}:=V\left(\frac{d}{d t}\right) \mathscr{B}$.

$$
w^{\prime} \in \mathscr{B}^{\prime} \Longleftrightarrow w^{\prime}=\left[\begin{array}{l}
w_{1}^{\prime} \\
w_{2}^{\prime}
\end{array}\right]
$$

with $w_{1}^{\prime} \in \operatorname{ker} D\left(\frac{d}{d t}\right), w_{2}^{\prime}$ free.

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$$
\text { If } D=I_{\mathrm{p}} \Longrightarrow \text { take } \mathscr{B}_{\text {contr }}^{\prime}=\mathscr{B}^{\prime}, \mathscr{B}_{\text {aut }}^{\prime}=\{0\} .
$$

## Decomposition of behaviors

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$$

with $\mathscr{B}_{\text {contr }}$ controllable and $\mathscr{B}_{\text {aut }}$ autonomous.

$$
\begin{gathered}
\text { If } D \neq I_{\mathrm{p}} \text {, define } \\
\mathscr{B}_{\text {contr }}^{\prime}=\left\{\left[\begin{array}{c}
w_{1}^{\prime} \\
0
\end{array}\right] \left\lvert\, w_{1}^{\prime} \in \operatorname{ker} D\left(\frac{d}{d t}\right)\right.\right\} \\
\mathscr{B}_{\text {aut }}^{\prime}=\left\{\left.\left[\begin{array}{c}
0 \\
w_{2}^{\prime}
\end{array}\right] \right\rvert\, w_{2}^{\prime} \in \mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{\mathrm{w}-\mathrm{p}}\right)\right\} .
\end{gathered}
$$

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with $\mathscr{B}_{\text {contr }}$ controllable and $\mathscr{B}_{\text {aut }}$ autonomous. Then transform back to $w$ variables.

## Observability



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¿Can $w_{2}$ be determined knowing $w_{1}$ and the system dynamics?

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¿Can $w_{2}$ be determined knowing $w_{1}$ and the system dynamics?
$\mathscr{B} \in \mathscr{L}^{\mathrm{w}}, w=\left(w_{1}, w_{2}\right) . w_{2}$ is observable from $w_{1}$ if

$$
\left(w_{1}, w_{2}^{\prime}\right),\left(w_{1}, w_{2}^{\prime \prime}\right) \in \mathscr{B} \Longrightarrow w_{2}^{\prime}=w_{2}^{\prime \prime}
$$

## Algebraic characterization of observability

Assume $\mathscr{B}$ represented in kernel form as

$$
R_{1}\left(\frac{d}{d t}\right) w_{1}+R_{2}\left(\frac{d}{d t}\right) w_{2}=0
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have a unique solution $w_{2}$ ?
It has $\Leftrightarrow R_{2}\left(\frac{d}{d t}\right)$ injective $\Leftrightarrow R_{2}(\lambda)$ has full column rank for all $\lambda \in \mathbb{C}$


¿Is $w_{2}$ observable from $w_{1}$ ?

## Example


¿Is $w_{2}$ observable from $w_{1}$ ?
¿Can one determine $w_{2}$
from knowledge of $w_{1}$ and the system dynamics?


$$
\left[\begin{array}{c}
m_{1} \frac{d^{2}}{d t^{2}}+c_{1} \frac{d}{d t}+k_{1} \\
-c_{1} \frac{d}{d t}-k_{1}
\end{array}\right] w_{1}=\left[\begin{array}{c}
c_{1} \frac{d}{d t}+k_{1} \\
-m_{2} \frac{d^{2}}{d t^{2}}-\left(c_{2}+c_{1}\right) \frac{d}{d t}-\left(k_{1}+k_{2}\right)
\end{array}\right] w_{2}
$$

## Example

$$
\left[\begin{array}{c}
m_{1} \frac{d^{2}}{d t^{2}}+c_{1} \frac{d}{d t}+k_{1} \\
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Is polynomial differential operator on RHS injective?

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$$
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\end{array}\right] w_{2}
$$

Is polynomial differential operator on RHS injective?

$$
\left[\begin{array}{c}
c_{1} \lambda+k_{1} \\
-m_{2} \lambda^{2}-\left(c_{2}+c_{1}\right) \lambda-\left(k_{1}+k_{2}\right)
\end{array}\right]
$$

has full column rank $\forall \lambda \in \mathbb{C}(\Longleftrightarrow$ observability $) \Leftrightarrow$

$$
-m_{2} k_{1}^{2}+c_{1} c_{2} k_{1}-k_{2} c_{2}^{2} \neq 0
$$

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Remarks
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Rank constancy test generalization of 'Hautus test' for state-space systems.

## Remarks

Rank constancy test generalization of 'Hautus test' for state-space systems.

Trajectory-, not representation-based definition as in state-space framework.

## Summary of Lecture 3

Polynomial differential operators and their properties are key;

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Inputs: free variables;

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- Autonomous systems;


## The main points

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- Polynomial differential operators and their properties are key;

Inputs: free variables;
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Controllability and observability: system, not representation, properties;

- Algebraic characterizations;
- Image representations.

End of Lecture 3

