



PASSIVE SYNTHESIS

of the

TERMINAL BEHAVIOR

of

CIRCUITS

JAN C. WILLEMS

K.U. Leuven

Synthesis problem

Informal formulation

Given a system, a *behavior*, and a set of *building blocks*, find an *architecture* and an *embedding* of building blocks such that the interconnected system realizes the given behavior.

We take a look at the following classical case:

- ▶ **behavior**: a linear time-invariant differential (LTID) current/voltage behavior,
- ▶ **building blocks**: linear *passive* resistors, inductors, capacitors, and **transformers**
↪ RLCT synthesis.

Pedigree

Ronald Foster

Wilhelm Cauer

Otto Brune

Raoul Bott & Richard Duffin

Bernard Tellegen

Brockway McMillan

Vitold Belevitch

Sidney Darlington

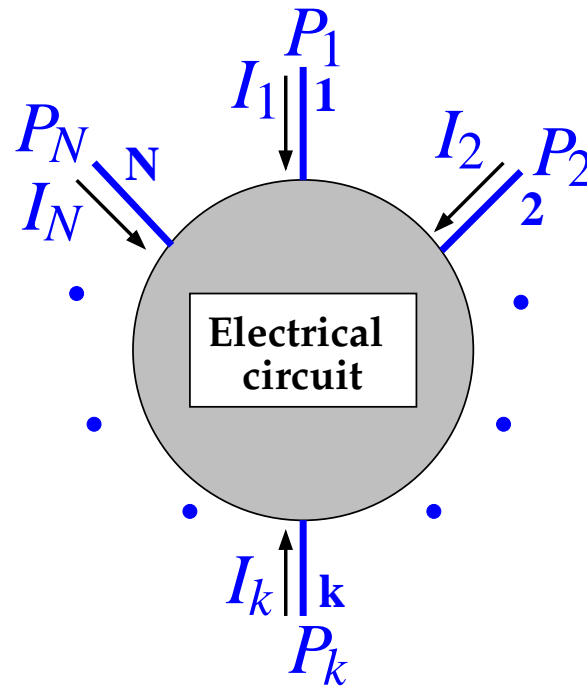
Dante Youla

and many others...

We add some footnotes to the work of these EE pioneers...

N -terminal circuits

Currents and potentials



At each terminal: a **current** and a **potential**

$\rightsquigarrow \Sigma = (\mathbb{R}, \mathbb{R}^N \times \mathbb{R}^N, \mathcal{B})$ **behavior** $\mathcal{B} \subseteq (\mathbb{R}^N \times \mathbb{R}^N)^{\mathbb{R}}$

$(I_1, I_2, \dots, I_N, P_1, P_2, \dots, P_N) : \mathbb{R} \rightarrow \mathbb{R}^N \times \mathbb{R}^N \in \mathcal{B}$ **means:**

this current/potential trajectory is compatible with the circuit architecture and its element values.

Elimination thm. \rightsquigarrow

RLCT circuit \Rightarrow LTID behavior

Synthesis

For which polynomial matrices $F \in \mathbb{R}[\xi]^{\bullet \times 2N}$ is

$$F \left(\frac{d}{dt} \right) \begin{bmatrix} I \\ P \end{bmatrix} = 0$$

the terminal behavior of an RLCT circuit?

!! *Given such an $F \in \mathbb{R}[\xi]^{\bullet \times 2N}$,
specify an RLCT circuit that has this terminal behavior !!*

Further cases of interest:

allow only: **RLC**, R, RC, RL, LC, RT, etc.

Our two footnotes

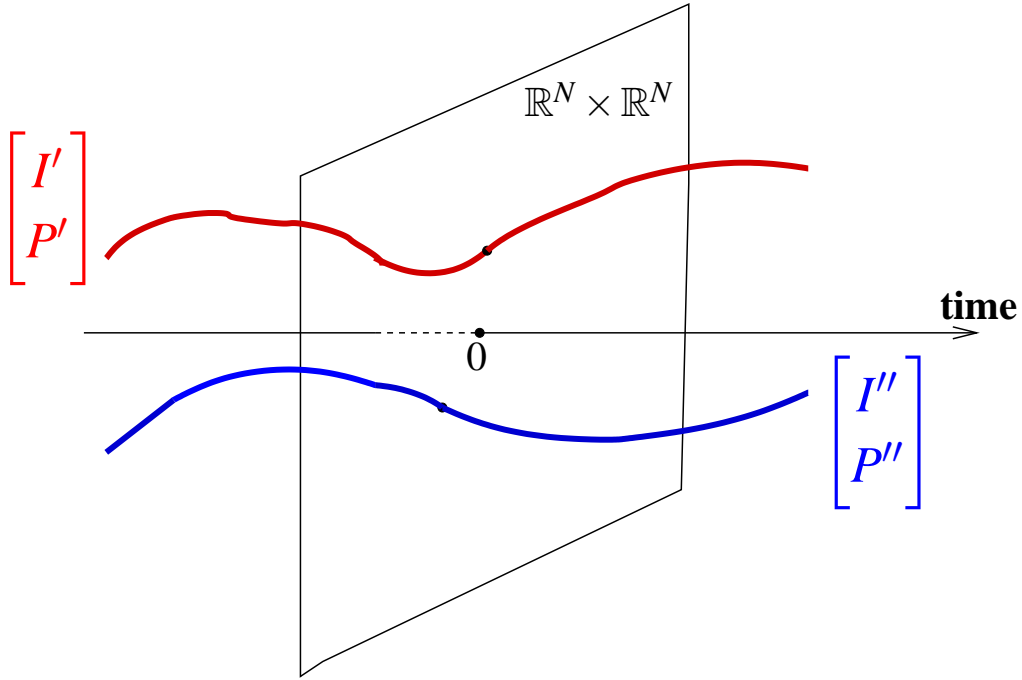
- ▶ Do we want to realize the *correct behavior*
or only the *correct controllable part* ?

Our two footnotes

- ▶ Do we want to realize the *correct behavior*
or only the *correct controllable part* ?
- ▶ Do we want to realize an *N-terminal* circuit,
or an *N-port* circuit?

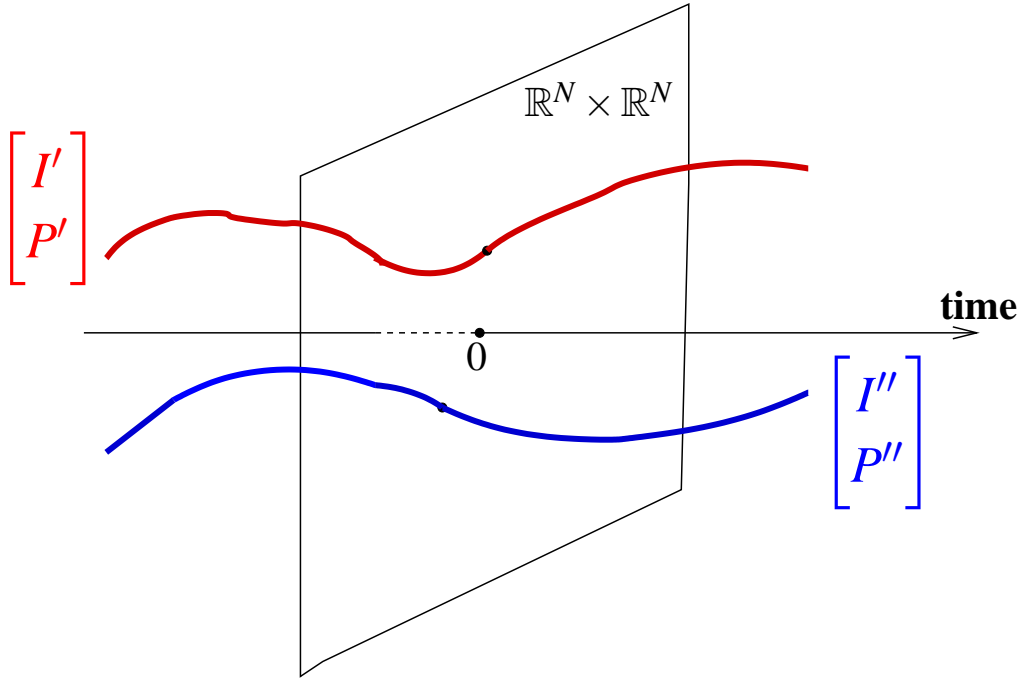
Controllability

Definition of controllability

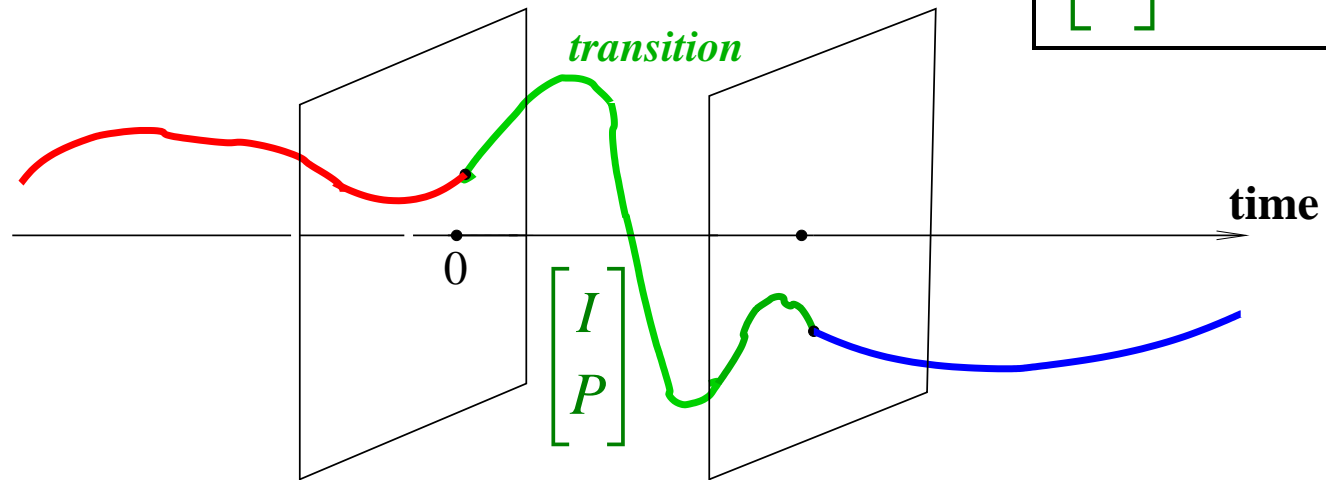


$$\begin{bmatrix} I' \\ P' \end{bmatrix}, \begin{bmatrix} I'' \\ P'' \end{bmatrix} \in \mathcal{B}$$

Definition of controllability



$$\begin{bmatrix} I \\ P \end{bmatrix} \in \mathcal{B}$$



controllability : \Leftrightarrow concatenability of trajectories after a delay .

Controllability of LTIDSs

The following are equivalent for $F \left(\frac{d}{dt} \right) \begin{bmatrix} I \\ P \end{bmatrix} = 0$.

- ▶ \mathcal{B} is **controllable**.
- ▶ F (WLOG full row rank) is **left prime**.
- ▶ ...

If $F = LF'$, with F' left prime, then

$$F' \left(\frac{d}{dt} \right) \begin{bmatrix} I \\ P \end{bmatrix} = 0$$

defines the '*controllable part*' of \mathcal{B} .

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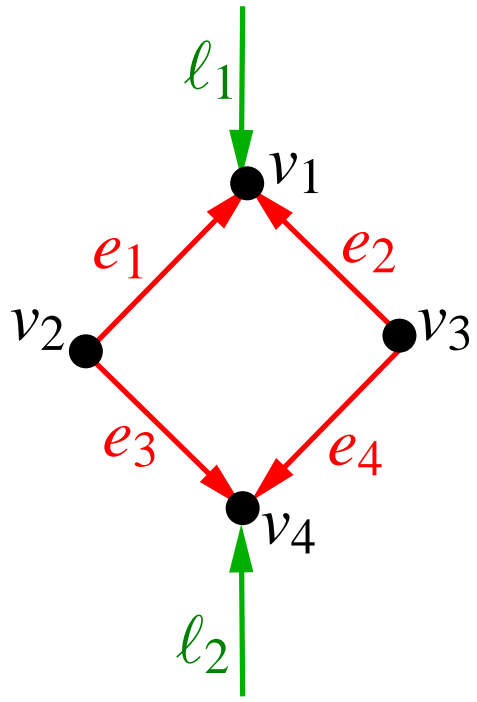
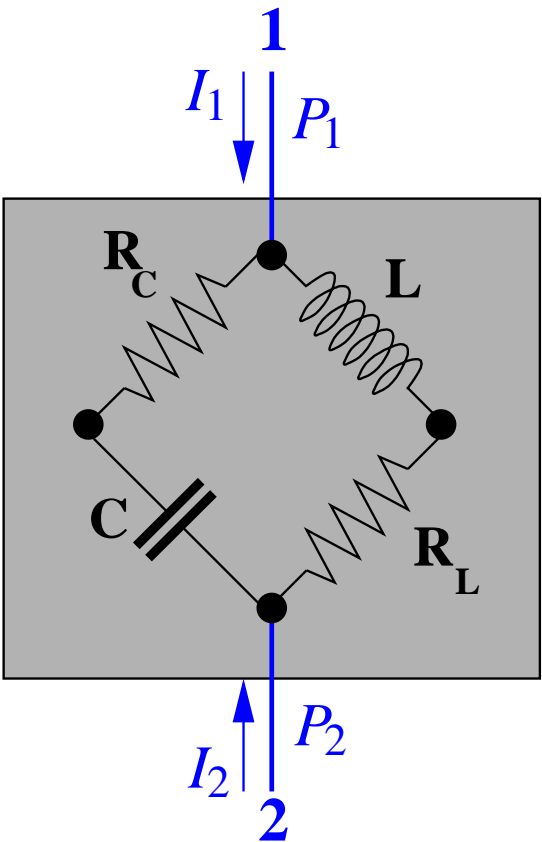
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Are uncontrollable circuits degenerate?

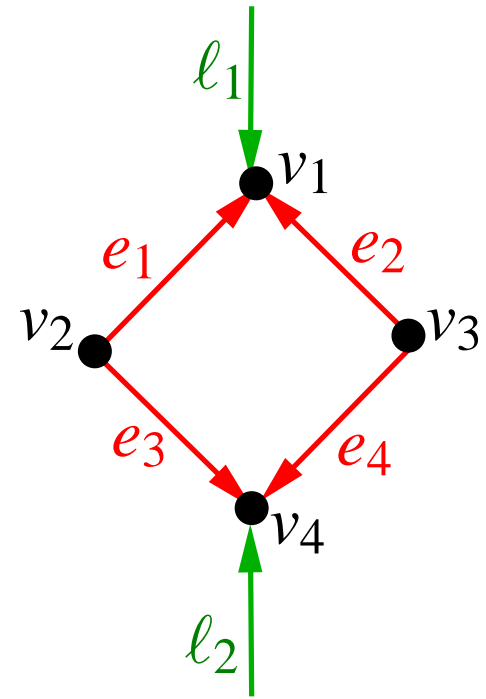
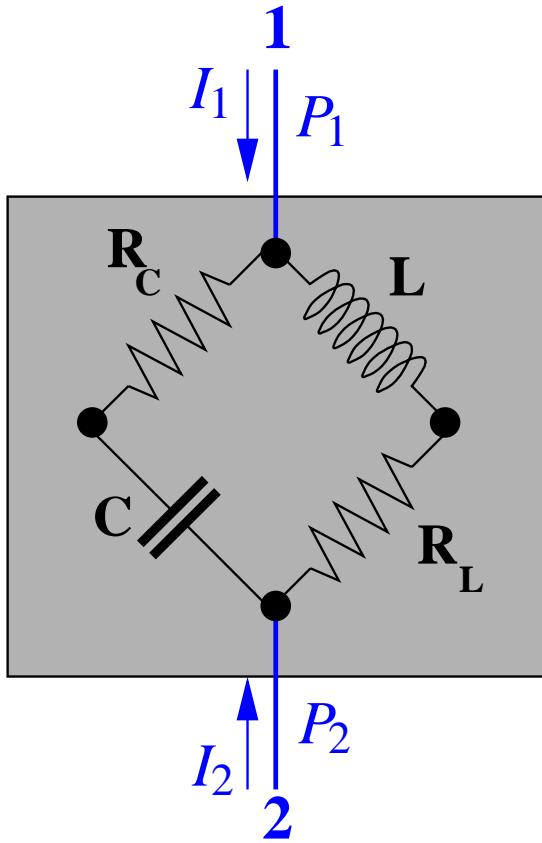
Example



$$A_V = \begin{bmatrix} +1 & +1 & 0 & 0 \\ -1 & 0 & -1 & 0 \\ 0 & -1 & 0 & -1 \\ 0 & 0 & +1 & +1 \end{bmatrix},$$

$$A_L = \begin{bmatrix} +1 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & +1 \end{bmatrix}.$$

Example



$$I = \begin{bmatrix} I_1 \\ I_2 \end{bmatrix}, P = \begin{bmatrix} P_1 \\ P_2 \end{bmatrix}, \quad I_{\mathbb{E}} = \begin{bmatrix} I_{e_1} \\ I_{e_2} \\ I_{e_3} \\ I_{e_4} \end{bmatrix}, P_{\mathbb{V}} = \begin{bmatrix} P_{v_1} \\ P_{v_2} \\ P_{v_3} \\ P_{v_4} \end{bmatrix}.$$

Behavioral equations

$$\begin{bmatrix} R_C & 0 & 0 & 0 \\ 0 & L \frac{d}{dt} & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & R_L \end{bmatrix} \begin{bmatrix} I_{e_1} \\ I_{e_2} \\ I_{e_3} \\ I_{e_4} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & C \frac{d}{dt} & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} -P_{v_1} + P_{v_2} \\ -P_{v_1} + P_{v_3} \\ P_{v_2} - P_{v_4} \\ P_{v_3} - P_{v_4} \end{bmatrix},$$

$$\begin{bmatrix} I_{e_1} + I_{e_2} + I_1 = 0 \\ I_{e_1} + I_{e_3} = 0 \\ I_{e_2} + I_{e_4} = 0 \\ I_{e_3} + I_{e_4} + I_2 = 0 \end{bmatrix}, \quad \begin{bmatrix} P_1 = P_{v_1} \\ P_2 = P_{v_4} \end{bmatrix}.$$

Elimination of $I_{\mathbb{E}}$ and $P_{\mathbb{V}} \rightsquigarrow$:

The circuit behavior

~> the following ODE defines \mathcal{B} .

Case 1: $CR_C \neq \frac{L}{R_L}$.

$$\begin{aligned} \left(\frac{R_C}{R_L} + \left(1 + \frac{R_C}{R_L} \right) CR_C \frac{d}{dt} + CR_C \frac{L}{R_L} \frac{d^2}{dt^2} \right) (P_1 - P_2) \\ = \left(1 + CR_C \frac{d}{dt} \right) \left(1 + \frac{L}{R_L} \frac{d}{dt} \right) R_C I_1, \end{aligned}$$

$$I_1 + I_2 = 0.$$

The circuit behavior

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Case 2:

$$CR_C = \frac{L}{R_L}.$$

$$\left(\frac{R_C}{R_L} + CR_C \frac{d}{dt} \right) (P_1 - P_2) = \left(1 + CR_C \frac{d}{dt} \right) R_C I_1,$$

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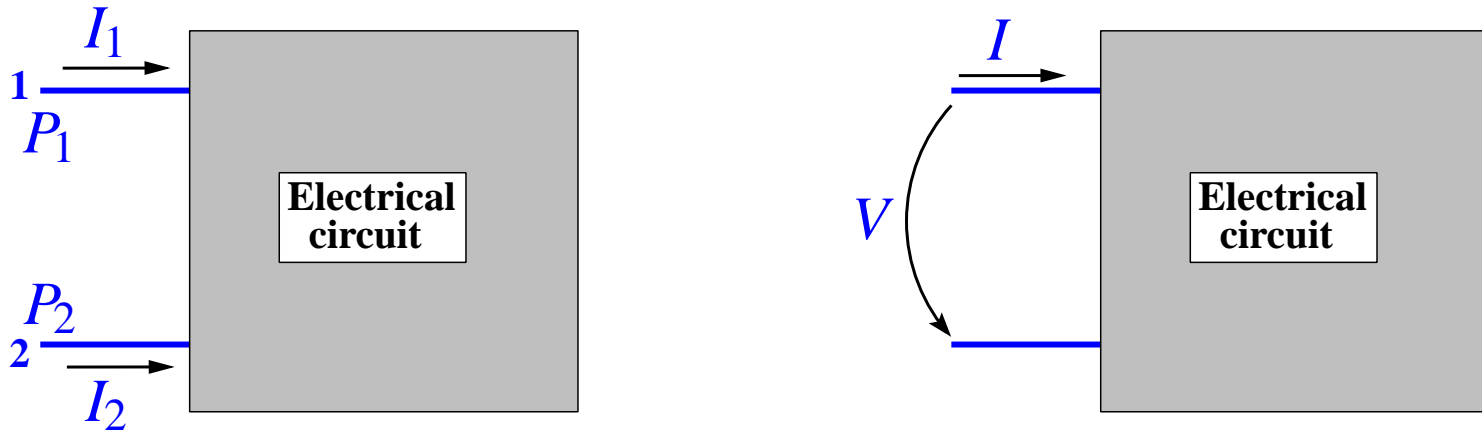
$$I_1 + I_2 = 0.$$

$$CR_C = \frac{L}{R_L} \text{ and } R_C = R_L \Leftrightarrow \text{uncontrollable.}$$

Hence: Linear passive circuits can become uncontrollable.

Realization of 2-terminal circuits

2-terminal circuits



KCL $\Rightarrow I_1 + I_2 = 0$, **KVL** \Rightarrow only $P_1 - P_2$ matters.

with $I := I_1 = -I_2$ and $V := P_1 - P_2$, this leads to

$$P \left(\frac{d}{dt} \right) V = Q \left(\frac{d}{dt} \right) I.$$

Define $Z := \frac{Q}{P}$ ‘*impedance*’.

2-terminal circuits

$$P \left(\frac{d}{dt} \right) V = Q \left(\frac{d}{dt} \right) I, \quad Z = \frac{Q}{P}.$$

**Which polynomial pairs (P, Q) are realizable
using RLCT? Using RLC?**

2-terminal circuits

$$P \left(\frac{d}{dt} \right) V = Q \left(\frac{d}{dt} \right) I, \quad Z = \frac{Q}{P}.$$

**Which polynomial pairs (P, Q) are realizable
using RLCT? Using RLC?**

Assume P and Q coprime (controllability).

Then RLCT realizable iff Z is positive real (Brune).

Iff Z is positive real,

then the controllable part is RLCT realizable (Brune).

Iff Z is positive real, then there exists RLC realization

with the ‘correct’ controllable part (Bott-Duffin).

Open problem

Which polynomial pairs (P, Q) are realizable using RLCT?

Necessary condition 1: $Z = \frac{Q}{P}$ is positive real.

Necessary condition 2: Uncontrollable part ‘stable’.

1 + 2 are not sufficient .

Sufficient condition: P and Q coprime, and $Z = \frac{Q}{P}$ p.r.

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Sufficient condition: P and Q coprime, and $Z = \frac{Q}{P}$ p.r.

Conclusions:

The set of RLCT realizable LTID behaviors is unknown .

Bott-Duffin realizes the impedance, but not the behavior .

Example

There is no present theory that guarantees that

$$\left(1 + \frac{d}{dt}\right) (P_1 - P_2) = \left(1 + \frac{d}{dt}\right) I_1, \quad I_1 + I_2 = 0$$

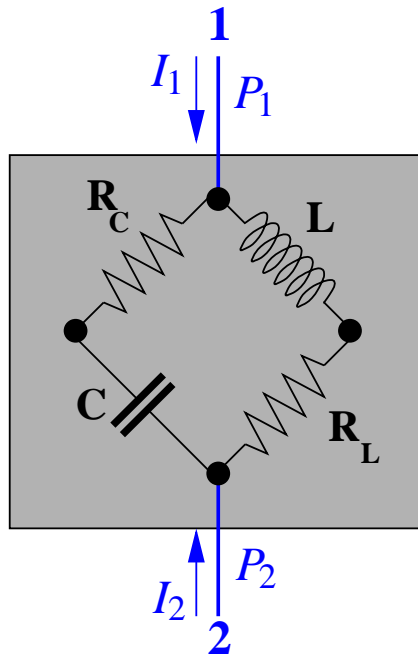
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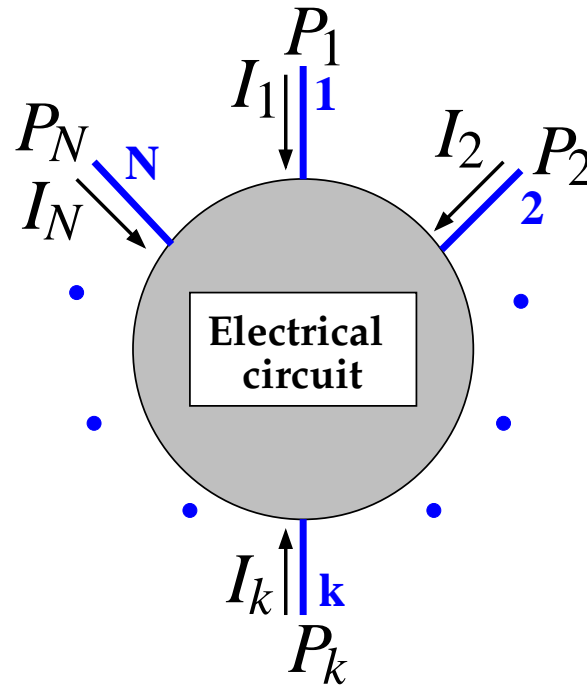
$$\left(1 + \frac{d}{dt}\right) (P_1 - P_2) = \left(1 + \frac{d}{dt}\right) I_1, \quad I_1 + I_2 = 0$$

is realizable. It is, using $R_C = R_L = 1, C = 1, L = 1$.



N -port versus N -terminal circuits

N -terminal circuit



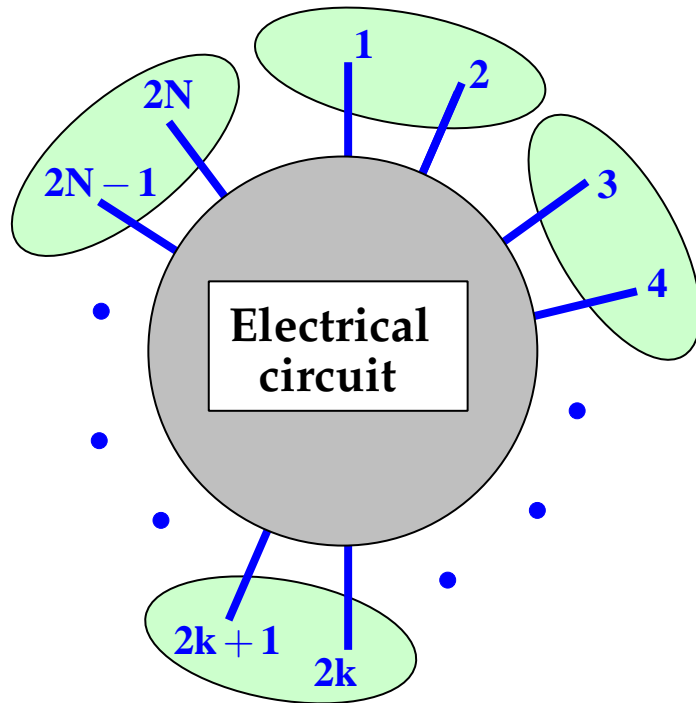
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this current/potential trajectory is compatible with the circuit architecture and its element values.

N -port



$2N$ -terminal circuit

behavior $\mathcal{B} \subseteq (\mathbb{R}^{2N} \times \mathbb{R}^{2N})^{\mathbb{R}}$

Pair the terminals, set

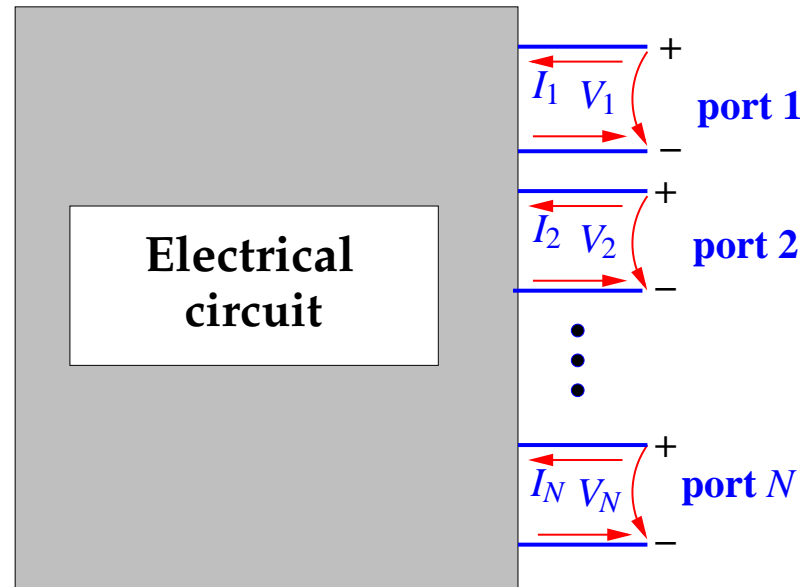
$$I_1 + I_2 = 0, \quad I_3 + I_4 = 0, \dots, \quad I_{2N-1} + I_{2N} = 0,$$

and take as variables the ‘port’ currents and ‘port’ voltages

$$I'_1 = I_1, \quad I'_2 = I_3, \dots, \quad I'_N = I_{2N-1},$$

$$V_1 = P_1 - P_2, \quad V_2 = P_3 - P_4, \dots, \quad V_N = P_{2N-1} - P_{2N}.$$

Currents and voltages



$\rightsquigarrow \Sigma = (\mathbb{R}, \mathbb{R}^N \times \mathbb{R}^N, \mathcal{B})$ port behavior $\mathcal{B} \subseteq (\mathbb{R}^N \times \mathbb{R}^N)^{\mathbb{R}}$

$(I_1, I_2, \dots, I_N, V_1, V_2, \dots, V_N) : \mathbb{R} \rightarrow \mathbb{R}^N \times \mathbb{R}^N \in \mathcal{B}$ means:

this current/voltage trajectory is compatible with the circuit *and the port current constraints*.

Classical synthesis problem

Given a LTID behavior $\mathcal{B} \subseteq (\mathbb{R}^N \times \mathbb{R}^N)^{\mathbb{R}}$,

find a $2N$ -terminal RLCT circuit with N -port behavior \mathcal{B} .

Classical synthesis problem

Given a LTID behavior $\mathcal{B} \subseteq (\mathbb{R}^N \times \mathbb{R}^N)^{\mathbb{R}}$,

find a $2N$ -terminal RLCT circuit with **N -port behavior** \mathcal{B} .

- ▶ For the 2-terminal case, KCL and KVL imply that 1-port synthesis is equivalent to 2-terminal synthesis.
- ▶ *If transformers are allowed* in the synthesis, then the results of the N -port case and the N -terminal case are transferrable.
Modulo controllability, a RLCT synthesis exists iff, **roughly**, the multivariable impedance is symmetric and positive real.
- ▶ Without transformers, the N -port and the N -terminal cases are distinct.

Resistive terminal synthesis

Transformerless resistive synthesis

The synthesis of resistive N -ports **without transformers** is one of the open problems of classical N -port synthesis.

For N -terminal synthesis, it can be solved completely.

Realizability

$$\begin{bmatrix} I_1 \\ I_2 \\ \vdots \\ I_N \end{bmatrix} = \begin{bmatrix} Y_{1,1} & Y_{1,2} & \cdots & Y_{1,N} \\ Y_{2,1} & Y_{2,2} & \cdots & Y_{2,N} \\ \vdots & \vdots & \vdots & \vdots \\ Y_{N,1} & Y_{N,2} & \cdots & Y_{N,N} \end{bmatrix} \begin{bmatrix} P_1 \\ P_2 \\ \vdots \\ P_N \end{bmatrix}$$

can be realized as an N -terminal circuit using only resistors
if and only if the matrix $Y \in \mathbb{R}^{N \times N}$

Realizability

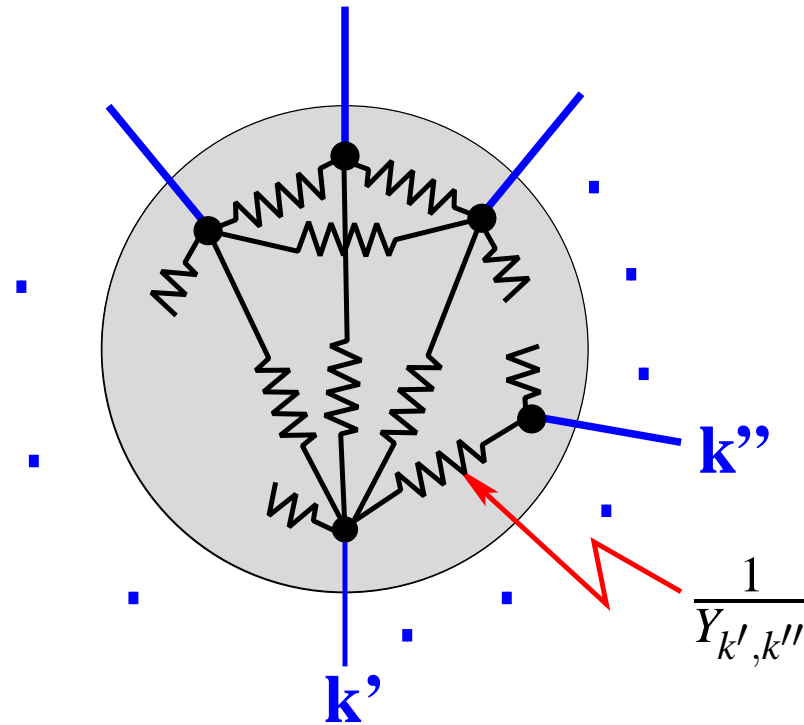
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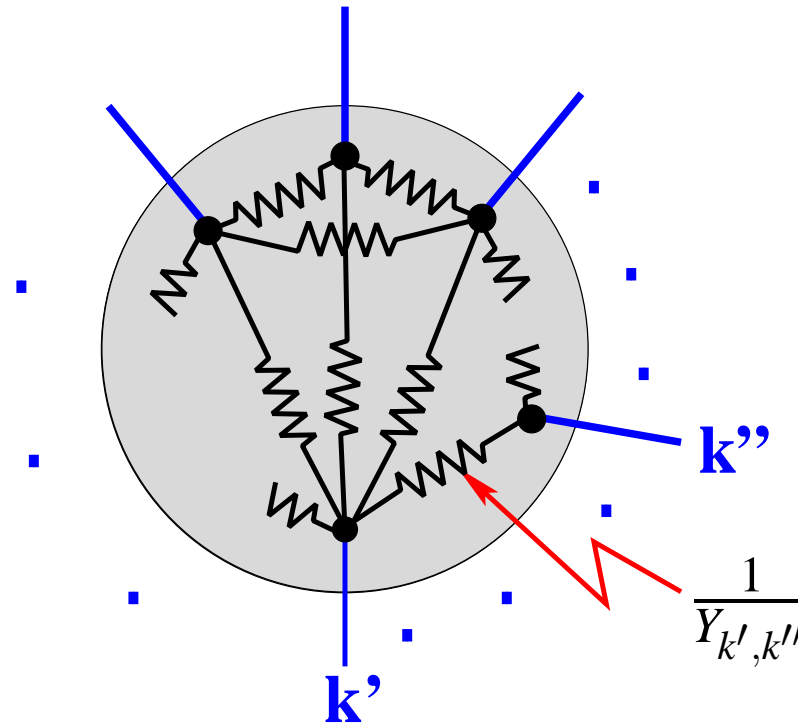
- ▶ is symmetric,
- ▶ has diagonal elements ≥ 0 ,
- ▶ off-diagonal elements ≤ 0 ,
- ▶ and row sums (and hence column sums) $= 0$.

‘hyperdominant with zero excess’.

Terminal synthesis of resistive circuits



Terminal synthesis of resistive circuits



- ▶ Generalizes to not-voltage-controlled case.

\rightsquigarrow realizability using only R 's of $F \begin{bmatrix} I \\ P \end{bmatrix} = 0$.

- ▶ Generalizes to inductive and capacitive circuits.

RLC terminal synthesis

Terminal synthesis of RLC circuits

$$\begin{bmatrix} I_1 \\ I_2 \\ \vdots \\ I_N \end{bmatrix} = \begin{bmatrix} Y_{1,1} \left(\frac{d}{dt} \right) & Y_{1,2} \left(\frac{d}{dt} \right) & \cdots & Y_{1,N} \left(\frac{d}{dt} \right) \\ Y_{2,1} \left(\frac{d}{dt} \right) & Y_{2,2} \left(\frac{d}{dt} \right) & \cdots & Y_{2,N} \left(\frac{d}{dt} \right) \\ \vdots & \vdots & \vdots & \vdots \\ Y_{N,1} \left(\frac{d}{dt} \right) & Y_{N,2} \left(\frac{d}{dt} \right) & \cdots & Y_{N,N} \left(\frac{d}{dt} \right) \end{bmatrix} \begin{bmatrix} P_1 \\ P_2 \\ \vdots \\ P_N \end{bmatrix}$$

can be realized as an N -terminal RLC circuit **if** $Y \in \mathbb{R}[\xi]^{N \times N}$

- ▶ is symmetric,
- ▶ has diagonal elements positive real,
- ▶ – off-diagonal elements positive real,
- ▶ and row sums (and hence column sums) = 0.

Not necessary!

Terminal synthesis of RLC circuits

$$\begin{bmatrix} I_1 \\ I_2 \\ \vdots \\ I_N \end{bmatrix} = \begin{bmatrix} Y_{1,1} \left(\frac{d}{dt} \right) & Y_{1,2} \left(\frac{d}{dt} \right) & \cdots & Y_{1,N} \left(\frac{d}{dt} \right) \\ Y_{2,1} \left(\frac{d}{dt} \right) & Y_{2,2} \left(\frac{d}{dt} \right) & \cdots & Y_{2,N} \left(\frac{d}{dt} \right) \\ \vdots & \vdots & \vdots & \vdots \\ Y_{N,1} \left(\frac{d}{dt} \right) & Y_{N,2} \left(\frac{d}{dt} \right) & \cdots & Y_{N,N} \left(\frac{d}{dt} \right) \end{bmatrix} \begin{bmatrix} P_1 \\ P_2 \\ \vdots \\ P_N \end{bmatrix}$$

can be realized as an N -terminal RLC circuit **only if**

$$Y \in \mathbb{R}[\xi]^{N \times N}$$

- ▶ is symmetric,
- ▶ has diagonal elements $Y_{k',k'}(\lambda) \geq 0$ for $\lambda \in \mathbb{R}, \lambda \geq 0$,
- ▶ off-diagonal elements $Y_{k',k''}(\lambda) \leq 0$ for $\lambda \in \mathbb{R}, \lambda \geq 0$,
- ▶ and row sums (and hence column sums) = 0.

Not sufficient!

Holy grail

Bott-Duffin synthesis from terminal point of view!

Conclusions

- ▶ **Classical impedance synthesis considers only the controllable part.**

RLCT synthesis of uncontrollable LTID behaviors is an open problem.

- ▶ **Classical impedance synthesis considers only the controllable part.**

RLCT synthesis of uncontrollable LTID behaviors is an open problem.

- ▶ **N -terminal synthesis is more natural than the classical N -port synthesis question.**

The resistive N -terminal synthesis problem is completely solvable.

For RLC N -terminal synthesis we presented a new necessary condition.

Copies of the lecture frames are available from/at

Jan.Willems@esat.kuleuven.be

<http://www.esat.kuleuven.be/~jwillems>

Thank you

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