



# **RATIONAL SYMBOLS**

Jan Willems, K.U. Leuven, Flanders, Belgium

Seminar, Kyoto University

June 26, 2008

– p. 1/5

#### Joint research with



#### Yutaka Yamamoto, Kyoto University, Japan

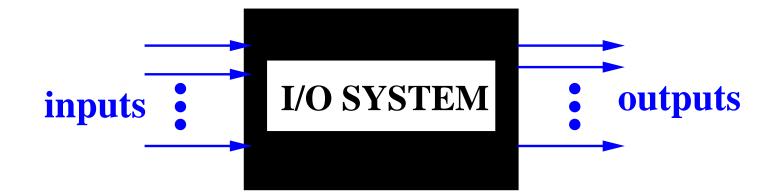


- I. Behaviors defined by rational symbols
- **II. Model reduction**
- (III. Parametrization of the stabilizing controllers)

# Introduction

**Motivation** 

In system theory, it is customary to think of dynamical models in terms of inputs and outputs.



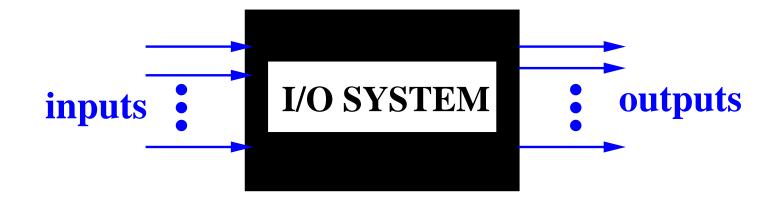
 $\rightsquigarrow$  say,

$$p_0y + p_1\frac{d}{dt}y + \dots + p_n\frac{d^n}{dt^n}y = q_0u + q_1\frac{d}{dt}u + \dots + q_n\frac{d^n}{dt^n}u$$

**i.e.,**  $p(\frac{d}{dt})y = q(\frac{d}{dt})u$ ,

**Motivation** 

In system theory, it is customary to think of dynamical models in terms of inputs and outputs.



$$p(\frac{d}{dt})y = q(\frac{d}{dt})u,$$
 or  $y = F(s)u$ 

with p, q polynomials, or F a rational transfer function.

#### **Motivation**

In system theory, it is customary to think of dynamical models in terms of inputs and outputs.

$$p(\frac{d}{dt})y = q(\frac{d}{dt})u,$$
 or  $y = F(s)u$ 

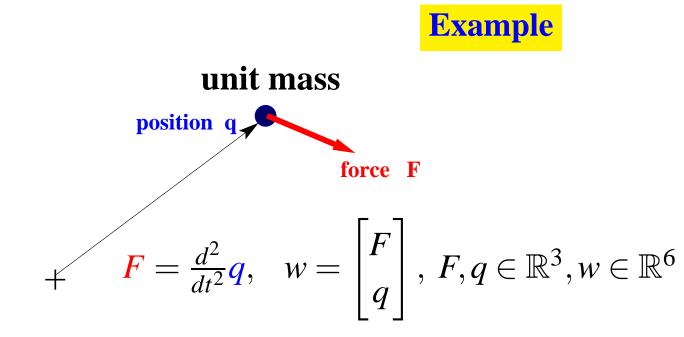
with p, q polynomials, or F a rational transfer function.

#### In the present talk, we will

**(for good reasons) make no distinction between** *u* **and** *y* 

$$w = \begin{bmatrix} u \\ y \end{bmatrix}$$

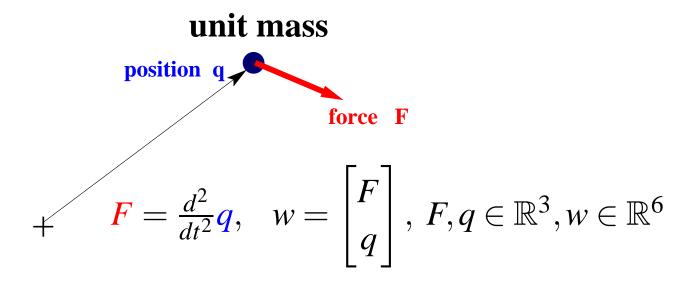
interpret F, not in terms of Laplace transforms, but in terms of differential equations. Important for, among other things, pedagogical reasons.





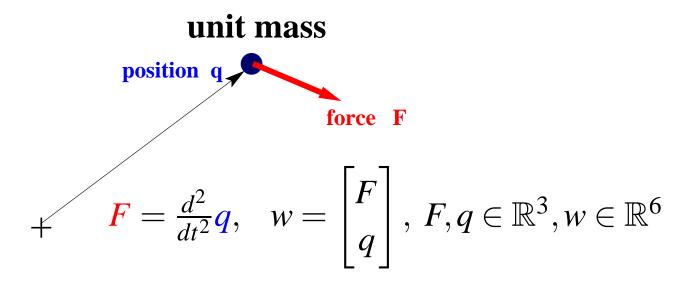
Isaac Newton by William Blake





$$\left[I_{3\times 3} : -\left(\frac{d}{dt}\right)^2 I_{3\times 3}\right] w = 0 \quad \rightsquigarrow q = \frac{1}{\left(\frac{d}{dt}\right)^2} F \quad \rightsquigarrow \quad \left[-\frac{1}{\left(\frac{d}{dt}\right)^2} I_{3\times 3} : I_{3\times 3}\right] w = 0$$





$$\left[I_{3\times3} : -\left(\frac{d}{dt}\right)^2 I_{3\times3}\right] w = 0 \quad \rightsquigarrow q = \frac{1}{\left(\frac{d}{dt}\right)^2} F \quad \rightsquigarrow \quad \left[-\frac{1}{\left(\frac{d}{dt}\right)^2} I_{3\times3} : I_{3\times3}\right] w = 0$$

In the scalar case with simple polynomials, it is easy to see how to proceed, but with general multivariable rational functions, less obvious. Today's pbm: What do we mean by

$$y = \frac{q(\frac{d}{dt})}{p(\frac{d}{dt})}u$$
, or  $G(\frac{d}{dt})w = 0$  with *G* rational?



# Linear time-invariant differential systems



# defined by rational symbols

– p. 7/5



A system 
$$\rightarrow$$
  $(\mathbb{T}, \mathbb{W}, \mathscr{B})$  where

■ 
$$T = set of independent variables
 $T = time \rightsquigarrow dynamical systems$   
 $T = time & space \rightsquigarrow distributed systems$$$

 $w : \mathbb{T} \to \mathbb{W}$  belongs to  $\mathscr{B}$  means: the model 'accepts' the trajectory w



# A dynamical system $\rightarrow$ $(\mathbb{R}, \mathbb{R}^{w}, \mathscr{B})$ where

- **•**  $\mathbb{T} = \text{set of independent variables } \longrightarrow \mathbb{T} = \mathbb{R}$  'time'
- $W = set of dependent variables; ~ W = \mathbb{R}^{W}$
- $\mathfrak{B} \quad \mathfrak{B} \quad \mathsf{the} \ \underline{behavior} \ \to \ \mathscr{B} \subseteq \mathbb{W}^{\mathbb{T}}, \\ \mathsf{time-trajectories} \ w : \mathbb{T} \to \mathbb{W}$ 
  - $\mathcal{B}$  = the solutions of a set of

linear constant coefficient ODEs



# A dynamical system $\rightarrow$ $(\mathbb{R}, \mathbb{R}^{w}, \mathscr{B})$ where

- **•**  $\mathbb{T} = \text{set of independent variables } \longrightarrow \mathbb{T} = \mathbb{R}$  'time'
- **•**  $\mathbb{W} = \text{set of dependent variables; } \rightarrow \mathbb{W} = \mathbb{R}^{\mathbb{W}}$
- $\mathscr{B}$  the *behavior*  $\to \mathscr{B} \subseteq \mathbb{W}^{\mathbb{T}}$ ,
- $\mathscr{B} =$ the solutions of

$$R_0 w + R_1 \frac{d}{dt} w + \dots + R_L \frac{d^L}{dt^L} w = 0, \ R_0, R_1, \dots$$
 matrices

**Polynomial matrix notation**  $\rightsquigarrow \frac{R\left(\frac{d}{dt}\right)w=0}{R \in \mathbb{R}\left[\xi\right]^{\bullet \times w}}$ 



A dynamical system 
$$\rightarrow$$
  $(\mathbb{R}, \mathbb{R}^{w}, \mathscr{B})$  where

- **•**  $\mathbb{T} =$ set of independent variables  $\rightarrow \mathbb{T} = \mathbb{R}$  'time'
- $W = set of dependent variables; ~ W = \mathbb{R}^{W}$
- $\mathscr{B}$  the *behavior*  $\to \mathscr{B} \subseteq \mathbb{W}^{\mathbb{T}}$ ,

 $\mathscr{B} = \operatorname{the} \mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{w})$ -solutions of

$$R_0 w + R_1 \frac{d}{dt} w + \dots + R_{\rm L} \frac{d^{\rm L}}{dt^{\rm L}} w = 0$$

**Polynomial matrix notation**  $\rightarrow R\left(\frac{d}{dt}\right)w = 0$ 

#### **Behaviors of LTIDSs allow many useful representations**

• As the set of solutions of  $R\left(\frac{d}{dt}\right)w = 0$   $R \in \mathbb{R}[\xi]^{\bullet \times w}$ 

**Behaviors of LTIDSs allow many useful representations** 

- As the set of solutions of  $R\left(\frac{d}{dt}\right)w = 0$   $R \in \mathbb{R}[\xi]^{\bullet \times w}$
- With input/output partition

$$P\left(\frac{d}{dt}\right)\mathbf{y} = Q\left(\frac{d}{dt}\right)\mathbf{u} \quad \mathbf{w} \cong \begin{bmatrix} \mathbf{u} \\ \mathbf{y} \end{bmatrix} \quad \det(P) \neq 0, P^{-1}Q \text{ proper}$$

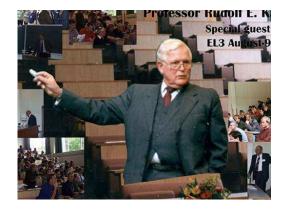
**Behaviors of LTIDSs allow many useful representations** 

- As the set of solutions of  $R\left(\frac{d}{dt}\right)w = 0$   $R \in \mathbb{R}[\xi]^{\bullet \times w}$
- With input/output partition

$$P\left(\frac{d}{dt}\right)y = Q\left(\frac{d}{dt}\right)u$$
  $w \cong$ 

Input/state/output representation
 ∃ matrices A, B, C, D such that
 ℬ consists of all w's generated by

$$\frac{d}{dt}x = Ax + Bu, \ y = Cx + Du \quad w \cong \begin{bmatrix} u \\ y \end{bmatrix}$$



Rudolf E. Kalman

# **Rational Symbols**

**Rational representations** 

In signal processing, control, etc., we often meet models that involve rational functions, instead of ODEs. Cfr. transfer functions,

$$y = F(\mathbf{s}')u$$

etc.  $\rightsquigarrow$ 

Let  $G \in \mathbb{R}(\xi)^{\bullet \times w}$ , and consider the 'differential equation'

 $G\left(\frac{d}{dt}\right)w = 0$  *G* is called the 'symbol'

What do we mean by its solutions, i.e. by the behavior?

**Rational representations** 

Let  $G \in \mathbb{R}(\xi)^{\bullet \times w}$ , and consider the 'differential equation'

 $G\left(\frac{d}{dt}\right)w = 0$  *G* is called the 'symbol'

What do we mean by its solutions, i.e. by the behavior?

$$\begin{bmatrix} M & \text{left prime} \end{bmatrix} :\Leftrightarrow \begin{bmatrix} M = FM' \end{bmatrix} \Rightarrow \begin{bmatrix} F \text{ unimodular} \end{bmatrix} \\ \Leftrightarrow & \exists H \text{ such that } MH = I. \\ \text{In scalar case, } M = \begin{bmatrix} m_1 & m_2 & \cdots & m_n \end{bmatrix}, \text{ this means:} \\ & m_1, m_2, \cdots, m_n \text{ have no common root.} \end{bmatrix}$$

**Rational representations** 

Let  $G \in \mathbb{R}(\xi)^{\bullet \times w}$ , and consider the 'differential equation'

 $G\left(\frac{d}{dt}\right)w = 0$  *G* is called the 'symbol'

What do we mean by its solutions, i.e. by the behavior? Let (P,Q) be a left coprime polynomial factorization of Gi.e.  $P,Q \in \mathbb{R}[\xi]^{\bullet \times \bullet}$ ,  $\det(P) \neq 0, G = P^{-1}Q, [P \vdots Q]$  left-prime.

E.g., in scalar case, means *P* and *Q* have no common roots.

$$\llbracket G(\frac{d}{dt})w = 0 \rrbracket \Leftrightarrow \llbracket P^{-1}Q(\frac{d}{dt})w = 0 \rrbracket :\Leftrightarrow \llbracket Q(\frac{d}{dt})w = 0 \rrbracket$$

**By definition**, therefore, the behavior of  $G(\frac{d}{dt})w = 0$  is equal to the behavior of  $Q(\frac{d}{dt})w = 0$ .

$$\llbracket G(\frac{d}{dt})w = 0 \rrbracket \Leftrightarrow \llbracket P^{-1}Q(\frac{d}{dt})w = 0 \rrbracket :\Leftrightarrow \llbracket Q(\frac{d}{dt})w = 0 \rrbracket$$

#### **Justification:**

**1.** *G* proper.  $G(\xi) = C(I\xi - A)^{-1}B + D$  controllable realization. Consider output nulling inputs:

$$\frac{d}{dt}x = Ax + Bw, \quad 0 = Cx + Dw$$

This set of *w*'s are exactly those that satisfy  $G\left(\frac{d}{dt}\right)w = 0$ . Analogous for  $\frac{d}{dt}x = Ax + Bw, 0 = Cx + D\left(\frac{d}{dt}\right)w, D \in \mathbb{R}[\xi]^{\bullet \times \bullet}$ .

- p. 11/5

$$\llbracket G(\frac{d}{dt})w = 0 \rrbracket \Leftrightarrow \llbracket P^{-1}Q(\frac{d}{dt})w = 0 \rrbracket :\Leftrightarrow \llbracket Q(\frac{d}{dt})w = 0 \rrbracket$$

#### **Justification:**

**2.** Consider y = G(s)w. View G(s) as a transfer f'n. Take your favorite definition of input/output pairs.

Output nulling inputs exactly those that satisfy  $G\left(\frac{d}{dt}\right)w = 0$ .

$$\llbracket G(\frac{d}{dt})w = 0 \rrbracket \Leftrightarrow \llbracket P^{-1}Q(\frac{d}{dt})w = 0 \rrbracket :\Leftrightarrow \llbracket Q(\frac{d}{dt})w = 0 \rrbracket$$

**<u>Note</u>!** With this def., we can deal with transfer functions,

$$y = F(\frac{d}{dt})u$$
, i.e.  $\left[F(\frac{d}{dt}) : -I\right] \begin{bmatrix} u \\ y \end{bmatrix} = 0$ 

**JBLIQUE FRANCA** 

- p. 11/5

with *F* a matrix of rational functions, while completely avoiding Laplace transforms, domains of convergence, and such mathematical traps.

# Caveats



#### Consider

$$y = F\left(\frac{d}{dt}\right)u$$

We now know what it means that  $(u, y) \in \mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{\bullet})$  satisfies this 'ODE'.

Is there a unique *y* for a given *u*?



#### Consider

$$y = F\left(\frac{d}{dt}\right)u$$

We now know what it means that  $(u, y) \in \mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{\bullet})$  satisfies this 'ODE'.

**Is there a unique** *y* **for a given** *u***?** 

$$F = P^{-1}Q$$
 coprime fact.  $\Leftrightarrow P^{-1}\begin{bmatrix} P & -Q \end{bmatrix}$  coprime fact.

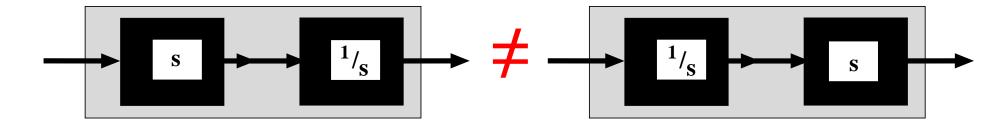
$$F = P^{-1}Q \quad \rightsquigarrow \quad y = F\left(\frac{d}{dt}\right)u \Leftrightarrow P\left(\frac{d}{dt}\right)y = Q\left(\frac{d}{dt}\right)u$$

If  $P \neq I$  (better, not unimodular), there are many sol'ns *y* of this ODE for a given *u*.

 $y = y_{\text{particular}} + y_{\text{homogeneous}}$ 

$$P(\frac{d}{dt})y_{\text{homogeneous}} = 0$$

 $G_1\left(\frac{d}{dt}\right)$  and  $G_2\left(\frac{d}{dt}\right)$  do not commute



$$G_1(s) = \frac{1}{s}$$
 and  $G_2(s) = s$   
 $y = \frac{1}{\frac{d}{dt}}v, \quad v = \frac{d}{dt}u \quad \Rightarrow \quad y(t) = u(t) + \text{ constant}$ 

 $y = \frac{d}{dt}v, v = \frac{1}{\frac{d}{dt}}u \Rightarrow y(t) = u(t)$ 

#### **LTIDSs** are **defined** in terms of **polynomial** symbols

$$\frac{R\left(\frac{d}{dt}\right)w=0}{R\in\mathbb{R}\left[\xi\right]^{\bullet\times\mathtt{w}}}$$

(behavior  $\mathscr{B}$ := the  $\mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{w})$  solutions)

#### **LTIDSs** are **defined** in terms of **polynomial** symbols

$$\frac{R\left(\frac{d}{dt}\right)w=0}{R\in\mathbb{R}\left[\xi\right]^{\bullet\times\mathbb{W}}}$$

(behavior  $\mathscr{B}$ := the  $\mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{W})$  solutions) but can also be represented by **rational** symbols

$$G\left(\frac{d}{dt}\right)w=0$$
  $G\in\mathbb{R}(\xi)^{\bullet imes w}$ 

#### **LTIDSs** are **defined** in terms of **polynomial** symbols

$$\frac{R\left(\frac{d}{dt}\right)w=0}{R\in\mathbb{R}\left[\xi\right]^{\bullet\times w}}$$

(behavior  $\mathscr{B}$ := the  $\mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{W})$  solutions) but can also be represented by **rational** symbols

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

**Behavior := the set of solutions of** 

$$Q\left(\frac{d}{dt}\right)w=0$$
  $Q\in\mathbb{R}\left[\xi\right]^{\bullet imes w}$ 

where  $G = P^{-1}Q$ ,  $P, Q \in \mathbb{R}[\xi]^{\bullet \times \bullet}$ , P and Q left coprime

#### **LTIDSs** are **defined** in terms of **polynomial** symbols

$$\frac{R\left(\frac{d}{dt}\right)w=0}{R\in\mathbb{R}\left[\xi\right]^{\bullet\times w}}$$

(behavior  $\mathscr{B}$ := the  $\mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{W})$  solutions) but can also be represented by **rational** symbols

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

This added flexibility  $\rightsquigarrow$  better adapted to certain applications,

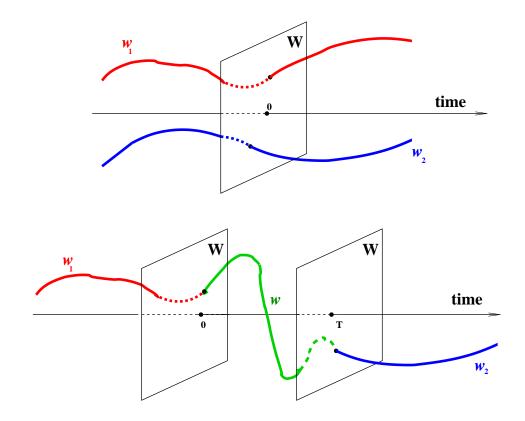
- e.g. distance between systems
- e.g. behavioral model reduction
- e.g. parametrization of the set of stabilizing controllers

# **Controllability c.s.**

**Controllability and stabilizability** 

 $\mathscr{B}$  is said to be **controllable** : $\Leftrightarrow$ 

 $\forall w_1, w_2 \in \mathscr{B}, \exists T \ge 0 \text{ and } w \in \mathscr{B} \text{ such that } \dots$ 

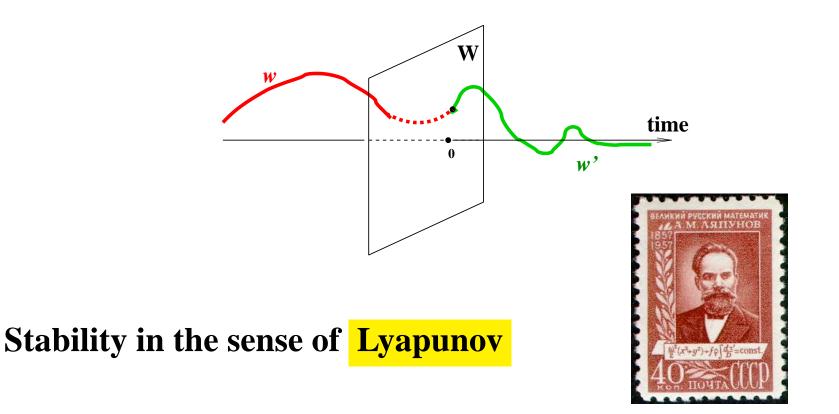


**Controllability and stabilizability** 

 $\mathscr{B}$  is said to be **controllable** : $\Leftrightarrow$ 

 $\mathscr{B}$  is said to be stabilizable : $\Leftrightarrow$ 

 $\forall w \in \mathscr{B}, \exists w' \in \mathscr{B}$  such that ...



What properties on *G* imply that the system with rational representation

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

has any of these properties?

What properties on *G* imply that the system with rational representation

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

has any of these properties?

Under what conditions on *G* does  $G\left(\frac{d}{dt}\right)w = 0$  define a controllable or a stabilizable system?

Can a rational representation be used to put one of these properties in evidence?



#### **Theorem:** The LTIDS

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

# is controllable if and only if

 $G(\lambda)$  has the same rank  $\forall \lambda \in \mathbb{C}$ 

## **Interpret carefully in cases like**

$$G(s) = \begin{bmatrix} s & 0 \\ 0 & \frac{1}{s} \end{bmatrix}, G(s) = \begin{bmatrix} s \\ \frac{1}{s} \end{bmatrix}, G(s) = \begin{bmatrix} s & \frac{1}{s} \end{bmatrix}$$



### **Theorem:** The LTIDS

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

# is controllable if and only if

 $G(\lambda)$  has the same rank  $\forall \lambda \in \mathbb{C}$ 

**Theorem:** The LTIDS

$$G\left(\frac{d}{dt}\right)w = 0 \qquad \qquad G \in \mathbb{R}(\xi)^{\bullet \times w}$$

is stabilizable if and only if

 $G(\lambda)$  has the same rank  $\forall \lambda \in \mathbb{C}$  with realpart $(\lambda) \geq 0$ 

For example,

# **Theorem:** A LTIDS is **controllable** if and only if its behavior allows an image representation

$$w = M(\frac{d}{dt})\ell$$
  $M \in \mathbb{R}(\xi)^{w \times \bullet}$ 

Module & vector spaces

### Take a LTIDS *B*.

# $n \in \mathbb{R}(\xi)^{1 \times w}$ is an annihilator $:\Leftrightarrow n(\frac{d}{dt})\mathscr{B} = 0$ , i.e.,

$$n(\frac{d}{dt})w = 0 \ \forall w \in \mathscr{B}$$

What structure does the set of annihilators of a given *B* have?

# Take a LTID behavior $\mathscr{B}$ .

 $n \in \mathbb{R}[\xi]^{1 \times w}$  is a polynomial annihilator : $\Leftrightarrow n(\frac{d}{dt})\mathscr{B} = 0$ 

# The polynomial annihilators form a $\mathbb{R}[\xi]$ -module: $n_1, n_2$ polynomial annihilators, $p \in \mathbb{R}[\xi]$ $\Rightarrow n_1 + pn_2$ polynomial annihilator.

# Take a LTID behavior $\mathscr{B}$ .

 $n \in \mathbb{R}[\xi]^{1 \times w}$  is a polynomial annihilator  $\Rightarrow n(\frac{d}{dt}) \mathscr{B} = 0$ 

The polynomial annihilators form a  $\mathbb{R}[\xi]$ -module:  $n_1, n_2$  polynomial annihilators,  $p \in \mathbb{R}[\xi]$  $\Rightarrow n_1 + pn_2$  polynomial annihilator.

 $n \in \mathbb{R}(\xi)^{1 \times w}$  is a rational annihilator : $\Leftrightarrow n(\frac{d}{dt})\mathscr{B} = 0$ 

The rational annihilators of a controllable  $\mathscr{B}$  form a  $\mathbb{R}(\xi)$ -vector space:  $n_1, n_2$  rational annihilators,  $p \in \mathbb{R}(\xi)$  $\Rightarrow n_1 + pn_2$  rational annihilator. **Module & vector spaces** 

By identifying a system with its polynomial annihilators, we obtain the one-to-one relation between LTIDSs with w variables and the

 $\mathbb{R}[\xi]$ - submodules of  $\mathbb{R}[\xi]^{\vee}$ 

By identifying a system with its rational annihilators, we obtain the one-to-one relation between the controllable LTIDSs with w variables and the

 $\mathbb{R}(\xi)$ -subspaces of  $\mathbb{R}(\xi)^{\mathtt{w}}$ 

**LTIDS**  $\cong$  finite dimensional  $\mathbb{R}[\xi]$ -modules

**Controllable LTIDS**  $\cong$  **finite dimensional**  $\mathbb{R}(\xi)$ **-subspaces.** 



# **Model reduction**

What is a good, computable, definition for the distance between two LTIDS?

**Basic issue underlying model reduction, robustness, etc.** 

- Approximate a system by a simpler one.
- If a system has a particular property (e.g., stabilized by a controller), will this also hold for close by systems?

What is meant by 'approximate', by 'close by'?

There is an elegant theory for reducing the state space dimension of stable LTI input/output systems. Let *B* be described by

$$\frac{d}{dt}x = Ax + Bu, \ y = Cx + Du \quad w \cong \begin{bmatrix} u \\ y \end{bmatrix}$$

# with A Hurwitz(: $\Leftrightarrow$ eigenvalues in left half plane).

There are effective methods (balancing, AAK) with good error bounds (in terms of the  $\mathscr{H}_{\infty}$  norm) for approximating  $\mathscr{B}$ by a (stable) system with a lower dimensional state space.



**Keith Glover** 

There is an elegant theory for reducing the state space dimension of stable LTI input/output systems. Let *B* be described by

$$\frac{d}{dt}x = Ax + Bu, \ y = Cx + Du$$

with A Hurwitz. T'f f'n  $F(s) = C(Is - A)^{-1}B + D$ 

proper stable rational. Reduced system

 $\frac{d}{dt}x_{\text{reduced}} = A_{\text{reduced}}x_{\text{reduced}} + B_{\text{reduced}}u, \ y = C_{\text{reduced}}x_{\text{reduced}} + Du$ 

**T'f f'n**  $F_{\text{reduced}}(s) = C_{\text{reduced}}(Is - A_{\text{reduced}})^{-1}B_{\text{reduced}} + D$ proper stable rational. Balanced model reduction  $\Rightarrow$ 

 $||F(i\omega) - F_{\text{reduced}}(i\omega)|| \le 2 \left(\sum_{\text{neglected Hankel SVs}} \sigma_{k}\right) \quad \forall \omega \in \mathbb{R}$ 

There is an elegant theory for reducing the state space dimension of stable LTI input/output systems. Let *B* be described by

$$\frac{d}{dt}x = Ax + Bu, \ y = Cx + Du$$

with A Hurwitz.

# F(s) proper stable rational $\Rightarrow$ reducible !

Extend this to situations where we do not make a distinction between inputs and outputs, and to unstable systems!

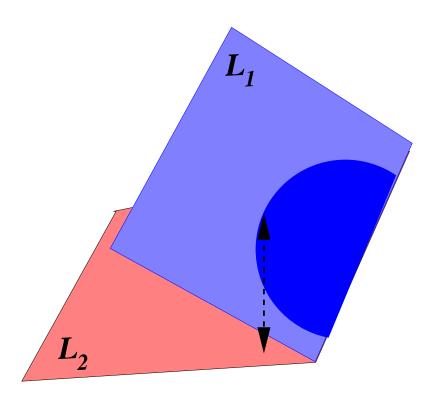
# **Distance between systems**

#### **Distance between linear subspaces**

In the behavioral theory, we identify a dynamical system with its behavior, a subspace  $\mathscr{B} \subseteq \mathfrak{C}^{\infty}(\mathbb{R}, \mathbb{R}^{w})$ . We are hence led to study the distance between linear subspaces of a vector space.

# $\mathscr{L}_1, \mathscr{L}_2 \subseteq \mathbb{R}^n$ , linear subspaces

$$\overrightarrow{d}\left(\mathscr{L}_{1},\mathscr{L}_{2}\right) \cong \max_{\substack{x_{1}\in\mathscr{L}_{1},||x_{1}||=1}} \min_{x_{2}\in\mathscr{L}_{2}}\left|\left|x_{1}-x_{2}\right|\right|$$



# $\mathscr{L}_1, \mathscr{L}_2 \subseteq \mathbb{R}^n$ , linear subspaces

$$d(\mathscr{L}_{1},\mathscr{L}_{2}) := \max\{\max_{x_{1}\in\mathscr{L}_{1},||x_{1}||=1} \min_{x_{2}\in\mathscr{L}_{2}} ||x_{1}-x_{2}||, \max_{x_{2}\in\mathscr{L}_{1},||x_{2}||=1} \min_{x_{1}\in\mathscr{L}_{1}} ||x_{1}-x_{2}||\}$$

$$0 \le d(\mathscr{L}_1, \mathscr{L}_2) \le 1$$

= 1 if dimension $(\mathscr{L}_1) \neq \text{dimension}(\mathscr{L}_2)$ 

 $\mathscr{L}_1, \mathscr{L}_2 \subseteq \mathbb{R}^n$ , linear subspaces  $P_{\mathscr{L}} \perp$  projection onto  $\mathscr{L}$  $S_1, S_2$  matrices, columns orthonormal basis for  $\mathscr{L}_1, \mathscr{L}_2$  $S_1 S_1^\top, S_2 S_2^\top$  orthogonal projectors

$$d(\mathscr{L}_{1},\mathscr{L}_{2}) = ||P_{\mathscr{L}_{1}} - P_{\mathscr{L}_{2}}|| \quad `gap', `aperture'$$
$$= ||S_{1}S_{1}^{\top} - S_{2}S_{2}^{\top}||$$
$$= \min_{\substack{\text{matrices } U}} ||S_{1} - S_{2}U||$$
$$= \min_{\substack{U \text{ such that } U\mathscr{L}_{1} = \mathscr{L}_{2}}} ||I - U||$$

 $\mathscr{L}_1, \mathscr{L}_2 \subseteq \mathbb{R}^n$ , linear subspaces  $P_{\mathscr{L}} \perp$  projection onto  $\mathscr{L}$   $S_1, S_2$  matrices, columns orthonormal basis for  $\mathscr{L}_1, \mathscr{L}_2$  $S_1 S_1^\top, S_2 S_2^\top$  orthogonal projectors

$$d(\mathscr{L}_{1},\mathscr{L}_{2}) = \frac{||P_{\mathscr{L}_{1}} - P_{\mathscr{L}_{2}}||}{||S_{1}S_{1}^{\top} - S_{2}S_{2}^{\top}||}$$
  
$$= \min_{\substack{\text{matrices } U}} ||S_{1} - S_{2}U||$$
  
$$= \min_{\substack{U \text{ such that } U\mathscr{L}_{1} = \mathscr{L}_{2}}} ||I - U||$$

Note

$$d(\mathscr{L}_1, \mathscr{L}_2) = ||S_1 S_1^\top - S_2 S_2^\top|| \le ||S_1 - S_2||$$

**Distance between controllable behaviors** 

 $\min \rightarrow \inf, \max \rightarrow \sup$ , etc., readily generalized to closed subspaces of Hilbert space.

For LTIDS, behaviors  $\mathscr{B} \mapsto \mathscr{B} \cap \mathscr{L}_2(\mathbb{R}, \mathbb{R}^w)$ . Keep notation. So, we consider only  $\mathscr{L}_2$ -behavior for measuring distance.

$$d(\mathscr{B}_1,\mathscr{B}_2) := gap(\mathscr{B}_1,\mathscr{B}_2)$$

 $\forall w_1 \in \mathscr{B}_1, \exists w_2 \in \mathscr{B}_2 \text{ such that } ||w_1 - w_2|| \leq \operatorname{gap}(\mathscr{B}_1, \mathscr{B}_2) ||w_1||$ 

and vice-versa. Small gap  $\Rightarrow$  the models are 'close'.

**Distance between controllable behaviors** 

 $\min \rightarrow \inf, \max \rightarrow \sup$ , etc., readily generalized to closed subspaces of Hilbert space.

For LTIDS, behaviors  $\mathscr{B} \mapsto \mathscr{B} \cap \mathscr{L}_2(\mathbb{R}, \mathbb{R}^w)$ . Keep notation. So, we consider only  $\mathscr{L}_2$ -behavior for measuring distance.

$$d(\mathscr{B}_1,\mathscr{B}_2) := gap(\mathscr{B}_1,\mathscr{B}_2)$$

 $\forall w_1 \in \mathscr{B}_1, \exists w_2 \in \mathscr{B}_2 \text{ such that } ||w_1 - w_2|| \leq \operatorname{gap}(\mathscr{B}_1, \mathscr{B}_2) ||w_1||$ 

and vice-versa. Small gap  $\Rightarrow$  the models are 'close'.

- Mow to compute the gap?
- Model reduce according to the gap!

# Let $\mathscr{B}$ be the behavior of a controllable LTIDS. Then it allows a rational symbol based image representation

$$w = M(\frac{d}{dt})\ell \quad \text{with} \quad M \in \mathbb{R} \left(\xi\right)^{w \times \bullet} \quad \& \quad M(-\xi)^{\top}M(\xi) = I$$
  
i.e.,  $||\ell||^{2}_{\mathscr{L}_{2}(\mathbb{R},\mathbb{R}^{\bullet})} = ||w||^{2}_{\mathscr{L}_{2}(\mathbb{R},\mathbb{R}^{w})} \quad \text{`norm preserving image repr.'}$   
$$\int_{-\infty}^{+\infty} ||w(t)||^{2} dt = \frac{1}{2\pi} \int_{-\infty}^{+\infty} ||\hat{w}(i\omega)||^{2} d\omega =$$
  
$$\frac{1}{2\pi} \int_{-\infty}^{+\infty} ||M(i\omega)\hat{\ell}(i\omega)||^{2} d\omega = \frac{1}{2\pi} \int_{-\infty}^{+\infty} ||\hat{\ell}(i\omega)||^{2} d\omega = \int_{-\infty}^{+\infty} ||\ell(t)||^{2} dt$$

**<u>Note</u>**: *M* cannot be polynomial, it must be rational. **Obviously** *M* must be proper. Can also make it stable.

Let  $\mathscr{B}$  be the behavior of a controllable LTIDS. Then it allows a rational symbol based image representation

$$w = M(\frac{d}{dt})\ell$$
 with  $M \in \mathbb{R}(\xi)^{w \times \bullet}$  &  $M(-\xi)^{\top}M(\xi) = I$ 

i.e.,  $||\ell||^2_{\mathscr{L}_2(\mathbb{R},\mathbb{R}^{\bullet})} = ||w||^2_{\mathscr{L}_2(\mathbb{R},\mathbb{R}^{w})}$  'norm preserving image repr.'

<u>Note</u>: *M* cannot be polynomial, it must be rational. Obviously *M* must be proper. Can also make it stable. **Proof:** Start with an observable polynomial image representation  $w = N(\frac{d}{dt})\ell, N \in \mathbb{R}[\xi]^{w \times m(\mathscr{B})}$ . Factor

 $N(-\xi)^{\top}N(\xi) = F(-\xi)^{\top}F(\xi), F \in \mathbb{R}\left[\xi\right]^{\mathrm{m}(\mathscr{B}) \times \mathrm{m}(\mathscr{B})}$ 

Can make determinant(F) Hurwitz. Take  $M = NF^{-1}$ .

Let  $\mathscr{B}$  be the behavior of a controllable LTIDS. Then it allows a rational symbol based image representation

$$w = M(\frac{d}{dt})\ell$$
 with  $M \in \mathbb{R}(\xi)^{w \times \bullet}$  &  $M(-\xi)^{\top}M(\xi) = I$ 

i.e.,  $||\ell||^2_{\mathscr{L}_2(\mathbb{R},\mathbb{R}^{\bullet})} = ||w||^2_{\mathscr{L}_2(\mathbb{R},\mathbb{R}^{w})}$  'norm preserving image repr.'

**Note:** *M* cannot be polynomial, it must be rational. Obviously *M* must be proper. Can also make it stable. Note that

$$f \in \mathscr{L}_2(\mathbb{R}, \mathbb{R}^{\mathsf{w}}) \mapsto M(i\omega)M(-i\omega)^{\top}\hat{f}(i\omega)$$

is the orthogonal projection onto  $\mathscr{B} \cap \mathscr{L}_2(\mathbb{R}, \mathbb{R}^w)$ .

# Let $\mathscr{B}$ be the behavior of a controllable LTIDS. Then it allows a rational symbol based image representation

$$w = M(\frac{d}{dt})\ell$$
 with  $M \in \mathbb{R}(\xi)^{w \times \bullet}$  &  $M(-\xi)^{\top}M(\xi) = I$ 

i.e.,  $||\ell||^2_{\mathscr{L}_2(\mathbb{R},\mathbb{R}^{\bullet})} = ||w||^2_{\mathscr{L}_2(\mathbb{R},\mathbb{R}^{w})}$  'norm preserving image repr.'

<u>Note</u>: *M* cannot be polynomial, it must be rational. **Obviously** *M* must be proper. Can also make it stable.  $\mathscr{B}_1 \mapsto M_1, \mathscr{B}_2 \mapsto M_2$  norm preserving, then

$$\begin{split} \mathsf{gap}(\mathscr{B}_1,\mathscr{B}_2) &= ||M_1(i\omega)M_1(-i\omega)^\top - M_2(i\omega)M_2(-i\omega)^\top||_{\mathscr{L}_{\infty}} \\ &\leq ||M_1(i\omega) - M_2(i\omega)||_{\mathscr{H}_{\infty}} \end{split}$$

**Model reduction by balancing** 

## Start with *B*. Take representatation

$$w = M(\frac{d}{dt})\ell$$
 with  $M \in \mathbb{R}(\xi)^{w \times \bullet}$  norm preserving, stable

Now model reduce  $w = M(\frac{d}{dt})\ell$  (viewed as a stable input/output system) using, for example, balancing

$$\rightsquigarrow w = M_{\texttt{reduced}}(\frac{d}{dt})\ell$$

and an error bound

$$||M - M_{\text{reduced}}||_{\mathscr{H}_{\infty}} \leq 2 \left(\sum_{\text{neglected SVs of } M} \sigma_{k}\right)$$

Start with stable norm preserving representation of  $\mathscr{B}$ 

$$w = M(\frac{d}{dt})\ell$$
 with  $M \in \mathbb{R}(\xi)^{w \times \bullet}$ 

Model reduce using balancing  $\rightsquigarrow w = M_{\text{reduced}}(\frac{d}{dt})\ell$ . Call behavior  $\mathscr{B}_{\text{reduced}}$ . Error bound

$$gap(\mathscr{B}, \mathscr{B}_{reduced}) = ||MM^{\top} - M_{reduced}M^{\top}_{reduced}||_{\mathscr{L}_{\infty}}$$

$$\leq ||M - M_{reduced}||_{\mathscr{H}_{\infty}}$$

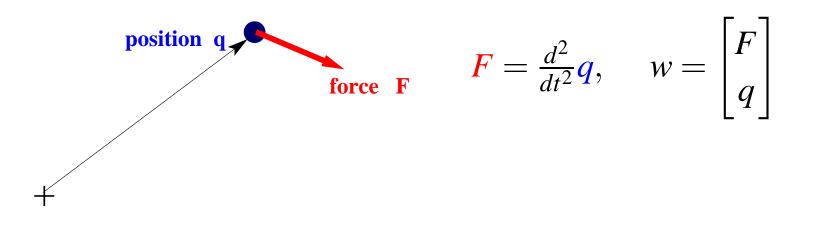
$$\leq 2 \left( \sum_{neglected SVs of M} \sigma_{k} \right)$$

 $\forall w \in \mathscr{B} \exists w' \in \mathscr{B}_{red}$  such that  $||w - w'|| \leq 2(\sum_{neglected SVs} \sigma_k)||w||$ 

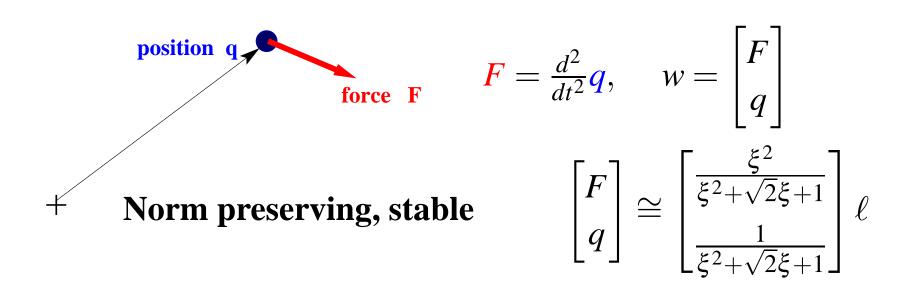
and vice-versa.

 $\sum_{neglected \ SVs} \sigma_k$  small  $\Rightarrow$  good approximation in the gap.

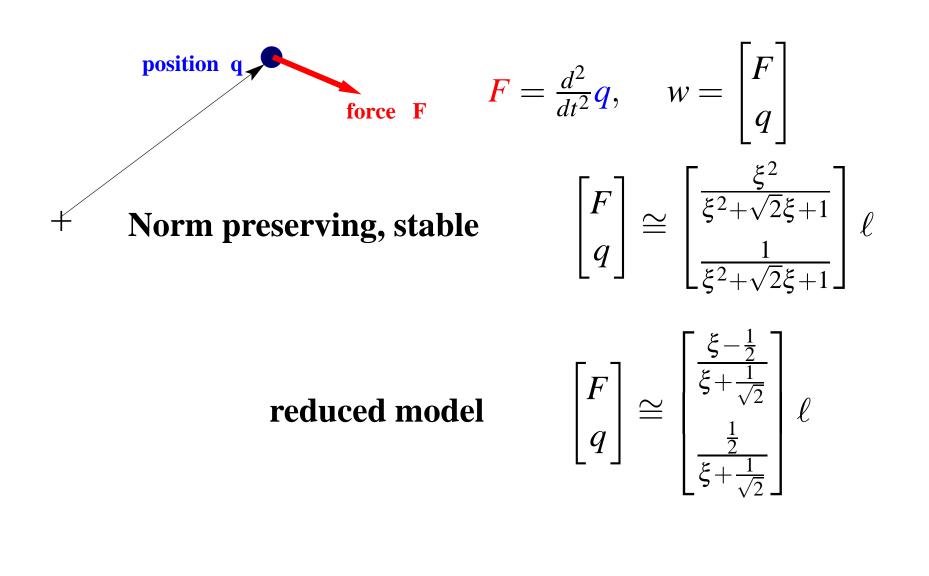




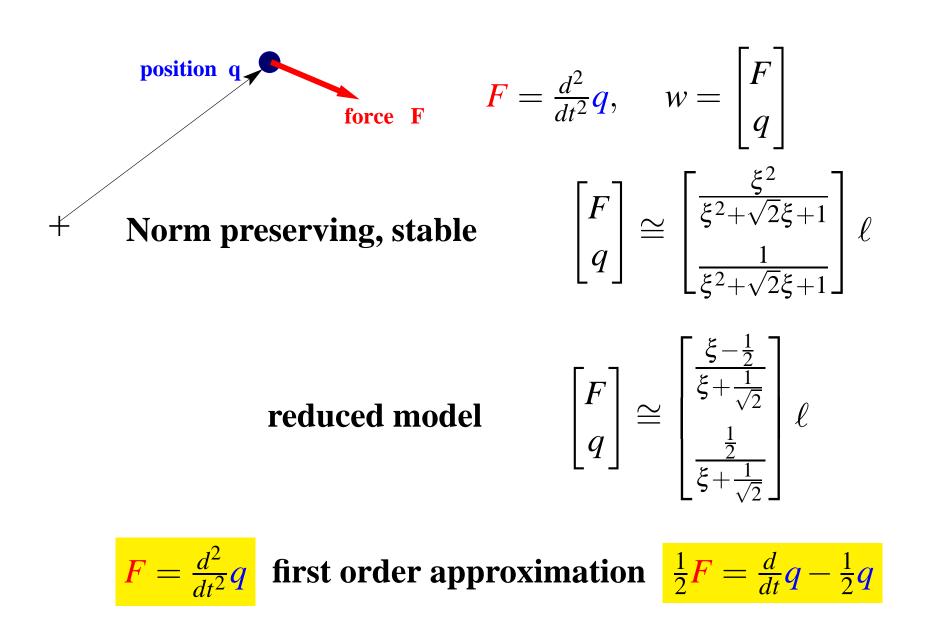
Example



Example



Example







# ■ $G(\frac{d}{dt})w = 0$ defined in terms left-coprime factorization of rational *G*.



- $G(\frac{d}{dt})w = 0$  defined in terms left-coprime factorization of rational *G*.
- ▶  $y = G(\frac{d}{dt})u$  does not require Laplace transform.

- $G(\frac{d}{dt})w = 0$  defined in terms left-coprime factorization of rational *G*.
- $y = G(\frac{d}{dt})u$  does not require Laplace transform.
- Controllability, stabilizability, etc. of  $G(\frac{d}{dt})w = 0$ decidable from *G*.

- $G(\frac{d}{dt})w = 0$  defined in terms left-coprime factorization of rational *G*.
- $y = G(\frac{d}{dt})u$  does not require Laplace transform.
- Controllability, stabilizability, etc. of  $G(\frac{d}{dt})w = 0$ decidable from *G*.
- Annihilators: finite dimensional R [ξ]-module.
   In controllable case, finite dimensional R (ξ)-vector space.

- $G(\frac{d}{dt})w = 0$  defined in terms left-coprime factorization of rational *G*.
- $y = G(\frac{d}{dt})u$  does not require Laplace transform.
- Controllability, stabilizability, etc. of  $G(\frac{d}{dt})w = 0$ decidable from *G*.
- Annihilators: finite dimensional R [ξ]-module.
   In controllable case, finite dimensional R (ξ)-vector space.
- Norm preserving representation  $w = M(\frac{d}{dt})\ell$  achievable with rational *M*.

- $G(\frac{d}{dt})w = 0$  defined in terms left-coprime factorization of rational *G*.
- $y = G(\frac{d}{dt})u$  does not require Laplace transform.
- **Controllability, stabilizability, etc. of**  $G(\frac{d}{dt})w = 0$ decidable from *G*.
- Annihilators: finite dimensional R [ξ]-module.
   In controllable case, finite dimensional R (ξ)-vector space.
- Norm preserving representation  $w = M(\frac{d}{dt})\ell$  achievable with rational *M*.
- Stable norm preserving representation  $w = M(\frac{d}{dt})\ell$  leads to model reduction of unstable systems and systems without input/output partition.



# **Parametrization of stabilizing controllers**

# $\mathbb{R}(\xi)$ and some of its subrings

Field of (real) rationals

**Subrings of interest** 

polynomials proper rationals stable rationals

proper stable rationals

## **unimodularity** $:\Leftrightarrow$ invertibility in the ring

Field of (real) rationalsnonzero

**Subrings of interest** 

polynomialsnonzero constantproper rationalsbiproperstable rationalsminiphaseproper stable rationalsbiproper & miniphase

## **unimodularity** $:\Leftrightarrow$ invertibility in the ring

Field of (real) rationalsnonzero

**Subrings of interest** 

polynomialsnonzero constantproper rationalsbiproperstable rationalsminiphaseproper stable rationalsbiproper & miniphase

unimodularity of square matrices over rings ⇔ determinant unimodular

left primeness of matrices over rings

$$:\Leftrightarrow \ \left[\!\left[ \left[\!\left[M = FM'\right]\!\right] \Rightarrow \left[\!\left[F \text{ unimodular}\right]\!\right]\!\right]\!\right]$$

The LTIDS *B* admits a representation that is **left prime** over

- rationals: always
- proper rationals: always
- **stable rationals: iff** *B* is **stabilizable**
- **proper stable rationals: iff** *B* is **stabilizable**
- **polynomials: iff**  $\mathscr{B}$  is controllable

Left prime representations over subrings allow to express certain system properties... The LTIDS *B* admits a representation that is left prime over

- **stable rationals: iff** *B* is **stabilizable**
- **proper stable rationals: iff** *B* is **stabilizable**

 $\mathscr{B}$  stabilizable  $\Leftrightarrow \exists G$ , matrix of rational functions, such that

- (i)  $\mathscr{B} = \operatorname{kernel}\left(G\left(\frac{d}{dt}\right)\right)$
- (ii) G is proper (no poles at  $\infty$ )
- (iii)  $G^{\infty} := \text{limit}_{\lambda \to \infty} G(\lambda)$  has full row rank (no zeros at  $\infty$ )
- (iv) G has no poles in  $\mathbb{C}_+ := \{\lambda \in \mathbb{C} \mid \texttt{real}(\lambda \ge 0\}$
- (v)  $G(\lambda)$  has full row rank  $\forall \lambda \in \mathbb{C}_+$  (no zeros in  $\mathbb{C}_+$ )

# **Unimodular completion**

**Unimodular completion lemma** 

Let G be a matrix over one of our rings (polynomial, proper rat., stable rat., proper stable rat.).

**;** Does there exist a unimodular completion G'i.e. a matrix G' over that same ring such that

$$\begin{bmatrix} G \\ G' \end{bmatrix}$$

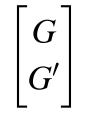
is unimodular (determinant is invertible in the ring) ?

**Unimodular completion lemma** 

Let G be a matrix over one of our rings (polynomial, proper rat., stable rat., proper stable rat.).

; There exists a unimodular completion G'

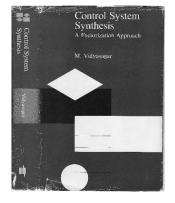
i.e. a matrix G' over that same ring such that



is unimodular

if and only if

*G* is **left prime** over the ring !





M. Vidyasagar

**Unimodular completion lemma** 

#### G: 1 row, 2 columns

$$G = \begin{bmatrix} p & q \end{bmatrix} \qquad G' = \begin{bmatrix} -y & x \end{bmatrix} \qquad \begin{bmatrix} G \\ G' \end{bmatrix} = \begin{bmatrix} p & q \\ -y & x \end{bmatrix}$$

#### G: 1 row, 2 columns

$$G = \begin{bmatrix} p & q \end{bmatrix} \qquad G' = \begin{bmatrix} -y & x \end{bmatrix} \qquad \begin{bmatrix} G \\ G' \end{bmatrix} = \begin{bmatrix} p & q \\ -y & x \end{bmatrix}$$

determinant = px + qy, unimodularity  $\Leftrightarrow px + qy = 1$ 

solvable for  $x, y \Leftrightarrow p$  & q coprime  $\Leftrightarrow G = \begin{bmatrix} p & q \end{bmatrix}$  left prime

#### G: 1 row, 2 columns

$$G = \begin{bmatrix} p & q \end{bmatrix} \qquad G' = \begin{bmatrix} -y & x \end{bmatrix} \qquad \begin{bmatrix} G \\ G' \end{bmatrix} = \begin{bmatrix} p & q \\ -y & x \end{bmatrix}$$

determinant = px + qy, unimodularity  $\Leftrightarrow px + qy = 1$ 

solvable for  $x, y \Leftrightarrow p$  & q coprime  $\Leftrightarrow G = \begin{bmatrix} p & q \end{bmatrix}$  left prime

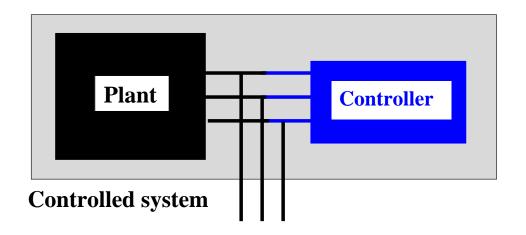
Our rings are **Hermite rings** 



*G* left prime  $\Leftrightarrow$  unimodularly completable  $\Leftrightarrow \exists H : GH = I \Leftrightarrow \cdots$ 

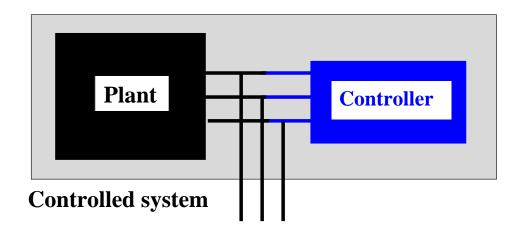
# Control





## Plant $\mathscr{P}$ , controller $\mathscr{C}$ , controlled system $\mathscr{P} \cap \mathscr{C}$



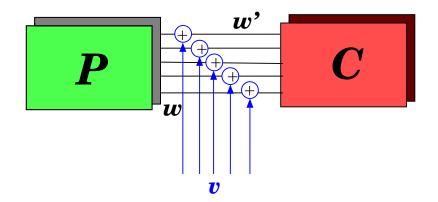


# Plant $\mathscr{P}$ , controller $\mathscr{C}$ , controlled system $\mathscr{P} \cap \mathscr{C}$ $\llbracket \mathscr{C} \text{ is stabilizing} \rrbracket :\Leftrightarrow \llbracket \mathscr{P} \cap \mathscr{C} \text{ is stable} \rrbracket$ $\Leftrightarrow \llbracket \llbracket w \in \mathscr{P} \cap \mathscr{C} \rrbracket \Rightarrow \llbracket w(t) \to 0 \text{ for } t \to \infty \rrbracket \rrbracket$



# $\llbracket \mathscr{C} \text{ is a } \operatorname{\mathbf{regular}} \operatorname{\mathbf{controller}} \rrbracket :\Leftrightarrow \llbracket \mathscr{P} + \mathscr{C} = \mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{\mathtt{W}}) \rrbracket$

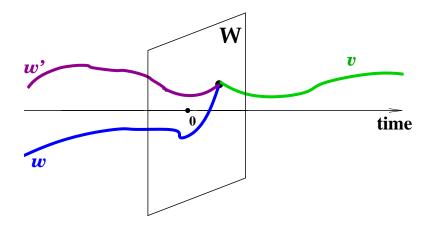
 $\forall v \in \mathscr{C}^{\infty}(\mathbb{R}, \mathbb{R}^{w}) \ \exists w \in \mathscr{P} \text{ and } w' \in \mathscr{C} \text{ such that } v = w + w'$ 



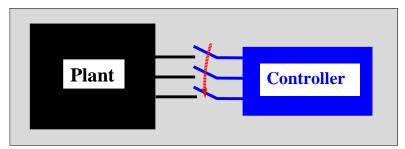


## $\llbracket \mathscr{C}$ is a superregular controller $\rrbracket :\Leftrightarrow$ in addition,

#### $\llbracket \forall w \in \mathscr{P}, \forall w' \in \mathscr{C} \ \exists v \text{ such that } w \wedge_0 v, w' \wedge_0 v \in \mathscr{P} \cap \mathscr{C} \rrbracket$



#### A superregular controller can be engaged at any time



**Controlled** system

#### superregular $\Rightarrow$ controller can be engaged at any time

(Super)regular controllers

# Usual feedback controllers are superregular PID controllers are regular, but not superregular

Controllers that are not superregular are relevant: control is interconnection, not just signal processing



**Harry Trentelman** 

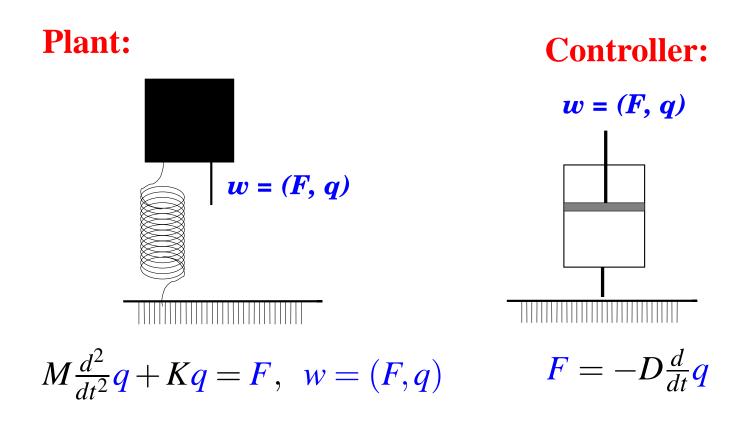


Madhu Belur

#### A regular, but not superregular, controller

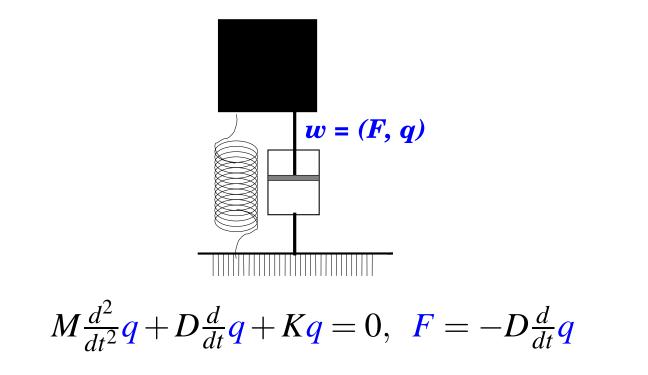
# Plant: w = (F, q) $M\frac{d^2}{dt^2}q + Kq = F, w = (F, q)$

#### A regular, but not superregular, controller



A regular, but not superregular, controller

#### **Controlled system:**



# **Existence of stabilizing controllers**



#### Proposition

## $\mathscr{P}$ is stabilizable $\Leftrightarrow \exists$ a regular stabilizing controller

 $\Leftrightarrow \exists$  a superregular stabilizing controller



#### Proposition

# $\mathscr{P}$ is stabilizable $\Leftrightarrow \exists$ a regular stabilizing controller

 $\Leftrightarrow \exists \ a \ superregular \ stabilizing \ controller$ 

## $\mathscr{P}$ is controllable $\Leftrightarrow \exists$ pole placement for $\mathscr{P} \cap \mathscr{C}$

**\nexists** a controller that is superregular &  $\mathscr{P} \cap \mathscr{C}$  has a low order characterisitic polynomial.

# **Parametrization of stabilizing controllers**

**Parametrization of superregular stabilizing controllers** 

Start with  $G\left(\frac{d}{dt}\right)w = 0$  a (rational symbol based) representation of the plant

Assume *G* left prime over proper stable rational functions. Iff the plant is stabilizable, such a *G* exists.

**Parametrization of superregular stabilizing controllers** 

Start with  $G\left(\frac{d}{dt}\right) w = 0$  a (rational symbol based) representation of the plant

Assume *G* left prime over proper stable rational functions. Iff the plant is stabilizable, such a G exists.



**Parametrization of superregular stabilizing controllers** 

Start with  $G\left(\frac{d}{dt}\right)w = 0$  a (rational symbol based) representation of the plant

Assume *G* left prime over proper stable rational functions. Iff the plant is stabilizable, such a G exists.

 $\Rightarrow \exists G' \text{ such that } \begin{vmatrix} G \\ G' \end{vmatrix} \text{ is unimodular over proper stable rat.}$ 

**Par'ion of superregular stabilizing controllers**  $C\left(\frac{d}{dt}\right)w = 0$ 

 $C = F_1 G + F_2 G'$ 

 $F_1$  free over ring of proper stable rational  $F_2$  unimodular over proper stable rational



Using **rational symbol** based representations  $G\left(\frac{d}{dt}\right)w = 0$ that are left prime over suitable rings, we obtain parametrizations of regular and superregular stabilizing controllers

 $\cong$  Kučera-Youla parametrization, with proper attention for the uncontrollable part



Vladimir Kučera



**Dante Youla** 



Margreta Kuijper

#### **Details & copies of the lecture frames are available from/at**

Jan.Willems@esat.kuleuven.be

http://www.esat.kuleuven.be/~jwillems

#### **Details & copies of the lecture frames are available from/at**

Jan.Willems@esat.kuleuven.be

http://www.esat.kuleuven.be/~jwillems

