

The Behavioral Approach to Systems Theory

**Paolo Rapisarda, Un. of Southampton, U.K.
&
Jan C. Willems, K.U.Leuven, Belgium**

**MTNS 2006
Kyoto, Japan, July 24–28, 2006**

Lecture 3: State and state construction

Lecturer: Paolo Rapisarda

Outline

The axiom of state

Discrete-time systems

First-order representations

State maps

The shift-and-cut map

Algebraic characterization

Continuous-time systems

Computation of state-space representations

Questions

- **Are state representations “natural”?**

Questions

- **Are state representations “natural”?**
 - **First principles and “tearing and zooming” modelling**
~> **systems of high-order differential equations**

Questions

- **Are state representations “natural”?**
 - **First principles and “tearing and zooming” modelling**
 \rightsquigarrow **systems of high-order differential equations**
 - **Algebraic constraints among variables**

Questions

- **Are state representations “natural”?**
 - **First principles and “tearing and zooming” modelling**
 \rightsquigarrow **systems of high-order differential equations**
 - **Algebraic constraints among variables**
- **What makes a latent variable a “state”?**

Questions

- **Are state representations “natural”?**
 - **First principles and “tearing and zooming” modelling**
 \rightsquigarrow **systems of high-order differential equations**
 - **Algebraic constraints among variables**
- **What makes a latent variable a “state”?**
- **What does that imply for the equations?**

Questions

- **Are state representations “natural”?**
 - **First principles and “tearing and zooming” modelling**
 \rightsquigarrow **systems of high-order differential equations**
 - **Algebraic constraints among variables**
- **What makes a latent variable a “state”?**
- **What does that imply for the equations?**
- **How to construct a state from the equations?**

Questions

- **Are state representations “natural”?**
 - **First principles and “tearing and zooming” modelling**
 \rightsquigarrow **systems of high-order differential equations**
 - **Algebraic constraints among variables**
- **What makes a latent variable a “state”?**
- **What does that imply for the equations?**
- **How to construct a state from the equations?**
- **How to construct a state representation from the equations?**

The basic idea

It's the Mariners' final game in the World Series. You're late...



The **current score is what matters...**

The basic idea

- The state contains all the **relevant information** about the **future** behavior of the system
- The state is the **memory** of the system
- **Independence of past and future** given the state

The axiom of state

$\Sigma = (\mathbb{T}, \mathbb{W}, \mathbb{X}, \mathfrak{B}_{\text{full}})$ is a **state system** if

$$(w_1, x_1), (w_2, x_2) \in \mathfrak{B}_{\text{full}} \text{ and } x_1(T) = x_2(T)$$

\Downarrow

$$(w_1, x_1) \underset{T}{\wedge} (w_2, x_2) \in \mathfrak{B}_{\text{full}}$$

$\underset{T}{\wedge}$ is **concatenation at T** :

$$(f_1 \underset{T}{\wedge} f_2)(t) := \begin{cases} f_1(t) & \text{for } t < T \\ f_2(t) & \text{for } t \geq T \end{cases}$$

Graphically...

$$(w_1, x_1), (w_2, x_2) \in \mathfrak{B}_{\text{full}} \text{ and } x_1(T) = x_2(T)$$

\Downarrow

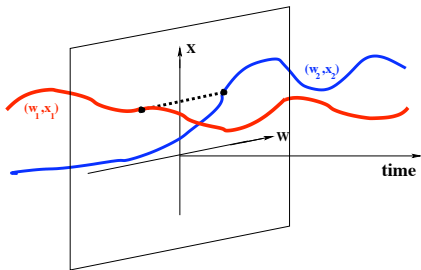
$$(w_1, x_1) \underset{T}{\wedge} (w_2, x_2) \in \mathfrak{B}_{\text{full}}$$

Graphically...

$$(\mathbf{w}_1, \mathbf{x}_1), (\mathbf{w}_2, \mathbf{x}_2) \in \mathcal{B}_{\text{full}} \text{ and } \mathbf{x}_1(T) = \mathbf{x}_2(T)$$

\Downarrow

$$(\mathbf{w}_1, \mathbf{x}_1) \underset{T}{\wedge} (\mathbf{w}_2, \mathbf{x}_2) \in \mathcal{B}_{\text{full}}$$

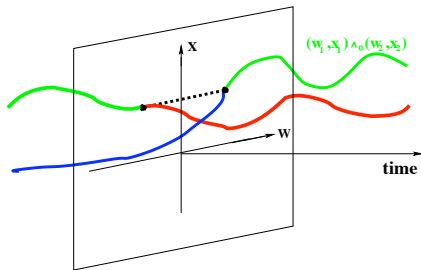
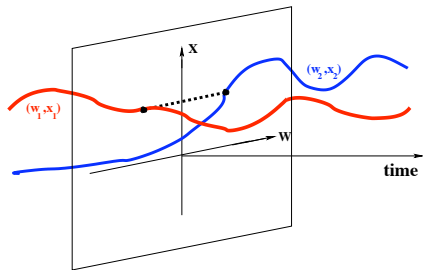


Graphically...

$$(\mathbf{w}_1, \mathbf{x}_1), (\mathbf{w}_2, \mathbf{x}_2) \in \mathfrak{B}_{\text{full}} \text{ and } \mathbf{x}_1(T) = \mathbf{x}_2(T)$$

\Downarrow

$$(\mathbf{w}_1, \mathbf{x}_1) \underset{T}{\wedge} (\mathbf{w}_2, \mathbf{x}_2) \in \mathfrak{B}_{\text{full}}$$



Example 1: discrete-time system

$\Sigma = (\mathbb{Z}, \mathbb{R}^w, \mathbb{R}^1, \mathcal{B}_{\text{full}})$, with

$$\mathcal{B}_{\text{full}} := \{(\mathbf{w}, \ell) \mid \mathbf{F} \circ (\sigma \ell, \ell, \mathbf{w}) = \mathbf{0}\}$$

where

$$\begin{aligned}\sigma &: (\mathbb{R}^1)^{\mathbb{Z}} \rightarrow (\mathbb{R}^1)^{\mathbb{Z}} \\ (\sigma(\ell))(k) &:= \ell(k+1)\end{aligned}$$

Example 1: discrete-time system

$\Sigma = (\mathbb{Z}, \mathbb{R}^w, \mathbb{R}^1, \mathcal{B}_{\text{full}})$, with

$$\mathcal{B}_{\text{full}} := \{(w, \ell) \mid F \circ (\sigma \ell, \ell, w) = \mathbf{0}\}$$

Special case: input-state-output equations

$$\sigma x = f(x, u)$$

$$y = h(x, u)$$

$$w = (u, y)$$

Example 2: continuous-time system

$\Sigma = (\mathbb{R}, \mathbb{R}^w, \mathbb{R}^1, \mathcal{B}_{\text{full}})$, with

$$\mathcal{B}_{\text{full}} := \left\{ (w, \ell) \mid F \circ \left(\frac{d}{dt} \ell, \ell, w \right) = \mathbf{0} \right\}$$

Special case: input-state-output equations

$$\begin{aligned} \frac{d}{dt} x &= f(x, u) \\ y &= h(x, u) \\ w &= (u, y) \end{aligned}$$

Outline

The axiom of state

Discrete-time systems

First-order representations

State maps

The shift-and-cut map

Algebraic characterization

Continuous-time systems

Computation of state-space representations

First-order discrete-time representations

Theorem: A 'complete' latent variable system

$$\Sigma = (\mathbb{Z}, \mathbb{R}^w, \mathbb{R}^x, \mathfrak{B}_{\text{full}})$$

is a state system iff $\mathfrak{B}_{\text{full}}$ can be described by

$$F \circ (\sigma x, x, w) = 0$$

First-order discrete-time representations

Theorem: A ‘complete’ latent variable system

$$\Sigma = (\mathbb{Z}, \mathbb{R}^w, \mathbb{R}^x, \mathfrak{B}_{\text{full}})$$

is a state system iff $\mathfrak{B}_{\text{full}}$ can be described by

$$F \circ (\sigma x, x, w) = 0$$

0-th order in w , 1st order in x

First-order discrete-time representations

Theorem: A ‘complete’ latent variable system

$$\Sigma = (\mathbb{Z}, \mathbb{R}^w, \mathbb{R}^x, \mathcal{B}_{\text{full}})$$

is a state system iff $\mathcal{B}_{\text{full}}$ can be described by

$$F \circ (\sigma x, x, w) = 0$$

Linear case:

$$E\sigma x + Fx + Gw = 0$$

First-order discrete-time representations

Theorem: A ‘complete’ latent variable system

$$\Sigma = (\mathbb{Z}, \mathbb{R}^w, \mathbb{R}^x, \mathcal{B}_{\text{full}})$$

is a state system iff $\mathcal{B}_{\text{full}}$ can be described by

$$F \circ (\sigma x, x, w) = 0$$

Linear case:

$$E\sigma x + Fx + Gw = 0$$

1st order in x is **consequence** of state property!

Proof (linear case)

$$\mathcal{V} := \left\{ \begin{bmatrix} a \\ b \\ c \end{bmatrix} \mid \exists (\mathbf{x}, \mathbf{w}) \in \mathfrak{B}_{\text{full}} \text{ s. t. } \begin{bmatrix} \mathbf{x}(1) \\ \mathbf{x}(0) \\ \mathbf{w}(0) \end{bmatrix} = \begin{bmatrix} a \\ b \\ c \end{bmatrix} \right\}$$

$$\mathcal{V} \text{ linear} \Rightarrow \exists \mathbf{E}, \mathbf{F}, \mathbf{G} \text{ s.t. } \mathcal{V} = \ker(\begin{bmatrix} \mathbf{E} & \mathbf{F} & \mathbf{G} \end{bmatrix})$$

Proof (linear case)

$$\mathcal{V} := \left\{ \begin{bmatrix} a \\ b \\ c \end{bmatrix} \mid \exists (x, w) \in \mathfrak{B}_{\text{full}} \text{ s. t. } \begin{bmatrix} x(1) \\ x(0) \\ w(0) \end{bmatrix} = \begin{bmatrix} a \\ b \\ c \end{bmatrix} \right\}$$

$$\mathcal{V} \text{ linear} \Rightarrow \exists E, F, G \text{ s.t. } \mathcal{V} = \ker(\begin{bmatrix} E & F & G \end{bmatrix})$$

↓

$$[(x, w) \in \mathfrak{B}_{\text{full}} \implies E\sigma x + Fx + Gw = 0]$$

Proof (linear case)

$$\mathcal{V} := \left\{ \begin{bmatrix} a \\ b \\ c \end{bmatrix} \mid \exists (x, w) \in \mathfrak{B}_{\text{full}} \text{ s. t. } \begin{bmatrix} x(1) \\ x(0) \\ w(0) \end{bmatrix} = \begin{bmatrix} a \\ b \\ c \end{bmatrix} \right\}$$

$$\mathcal{V} \text{ linear} \Rightarrow \exists E, F, G \text{ s.t. } \mathcal{V} = \ker(\begin{bmatrix} E & F & G \end{bmatrix})$$

Converse by induction, using axiom of state:

$$E\sigma x + Fx + Gw = 0 \text{ on } [0, k] \implies (w, x)_{|[0, k]} \in \mathfrak{B}_{\text{full}|[0, k]}$$

Then apply completeness of \mathfrak{B}

State construction: basic idea

Problem: Given kernel or hybrid description, find a state representation

$$E\sigma x + Fx + Gw = 0$$

State construction: basic idea

Problem: Given kernel or hybrid description, find a state representation

$$E\sigma x + Fx + Gw = 0$$

First compute polynomial operator in the shift acting on system variables, inducing a state variable:

$$X(\sigma)w = x$$

$$X(\sigma) \begin{bmatrix} w \\ \ell \end{bmatrix} = x$$

State construction: basic idea

Problem: Given kernel or hybrid description, find a state representation

$$E\sigma x + Fx + Gw = 0$$

First compute polynomial operator in the shift acting on system variables, inducing a state variable:

$$X(\sigma)w = x$$

$$X(\sigma) \begin{bmatrix} w \\ \ell \end{bmatrix} = x$$

Then use original eqs. and X to obtain 1st order representation.

State maps for kernel representations

$X \in \mathbb{R}^{\bullet \times w}[\xi]$ induces a **state map $X(\sigma)$** for $\ker(R(\sigma))$ if the behavior $\mathfrak{B}_{\text{full}}$ with latent variable x , described by

$$\begin{aligned} R(\sigma)w &= 0 \\ X(\sigma)w &= x \end{aligned}$$

satisfies the axiom of state.

Example

$$\mathfrak{B} = \{w \mid r(\sigma)w = 0\}$$

where $r \in \mathbb{R}[\xi]$, $\deg(r) = n$.

(Minimal) state map induced by

$$\begin{bmatrix} 1 \\ \xi \\ \vdots \\ \xi^{n-1} \end{bmatrix} \rightsquigarrow \begin{bmatrix} W \\ \sigma W \\ \vdots \\ \sigma^{n-1} W \end{bmatrix}$$

The axiom of state revisited

A **linear** system $\Sigma = (\mathbb{T}, \mathbb{W}, \mathbb{X}, \mathfrak{B}_{\text{full}})$ with latent variable x is a state system if

$$(w, x) \in \mathfrak{B}_{\text{full}} \text{ and } x(T) = 0$$

↓

$$(0, 0) \underset{T}{\wedge} (w, x) \in \mathfrak{B}_{\text{full}}$$

The axiom of state revisited

A **linear** system $\Sigma = (\mathbb{T}, \mathbb{W}, \mathbb{X}, \mathfrak{B}_{\text{full}})$ with latent variable x is a state system if

$$(w, x) \in \mathfrak{B}_{\text{full}} \text{ and } x(T) = 0$$

\Downarrow

$$(0, 0) \underset{T}{\wedge} (w, x) \in \mathfrak{B}_{\text{full}}$$

- Time-invariance \implies can choose $T = 0$;

The axiom of state revisited

A **linear** system $\Sigma = (\mathbb{T}, \mathbb{W}, \mathbb{X}, \mathcal{B}_{\text{full}})$ with latent variable x is a state system if

$$(w, x) \in \mathcal{B}_{\text{full}} \text{ and } x(T) = 0$$

↓

$$(0, 0) \underset{T}{\wedge} (w, x) \in \mathcal{B}_{\text{full}}$$

- Time-invariance \implies can choose $T = 0$;
- **Concatenability with zero trajectory** is key.

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

... 0 0 R_0 R_1 R_2 R_3 ...

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

$$R_0 w(0) + R_1 w(1) + \dots + R_L w(L) = 0$$

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

$$\dots \quad 0 \quad R_0 \quad R_1 \quad R_2 \quad R_3 \quad R_4 \quad \dots$$

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

$$R_0 w(0) + R_1 w(1) + \dots + R_L w(L) = 0$$

$$R_1 w(0) + R_2 w(1) + \dots + R_L w(L-1) = 0$$

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

... R_0 R_1 R_2 R_3 R_4 R_5 ...

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

$$R_0 w(0) + R_1 w(1) + \dots + R_L w(L) = 0$$

$$R_1 w(0) + R_2 w(1) + \dots + R_L w(L - 1) = 0$$

$$R_2 w(0) + R_3 w(1) + \dots + R_L w(L - 2) = 0$$

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

$$R_0 w(0) + R_1 w(1) + \dots + R_L w(L) = 0$$

$$R_1 w(0) + R_2 w(1) + \dots + R_L w(L - 1) = 0$$

$$R_2 w(0) + R_3 w(1) + \dots + R_L w(L - 2) = 0$$

$$\vdots \qquad = \vdots$$

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

$$\dots \quad R_{L-3} \quad R_{L-2} \quad R_{L-1} \quad R_L \quad 0 \quad 0 \quad \dots$$

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

$$R_0 w(0) + R_1 w(1) + \dots + R_L w(L) = 0$$

$$R_1 w(0) + R_2 w(1) + \dots + R_L w(L-1) = 0$$

$$R_2 w(0) + R_3 w(1) + \dots + R_L w(L-2) = 0$$

$$\vdots \quad = \quad \vdots$$

$$R_{L-1} w(0) + R_L w(1) = 0$$

When is $w \in \mathfrak{B}$ concatenable with zero?

$$R_0 w + R_1 \sigma w + \dots + R_L \sigma^L w = 0$$

$$\dots \quad R_{L-2} \quad R_{L-1} \quad R_L \quad 0 \quad 0 \quad 0 \quad \dots$$

...	0	0	$w(0)$	$w(1)$	$w(2)$	$w(3)$...
...	$k = -2$	$k = -1$	$k = 0$	$k = 1$	$k = 2$	$k = 3$...

$$R_0 w(0) + R_1 w(1) + \dots + R_L w(L) = 0$$

$$R_1 w(0) + R_2 w(1) + \dots + R_L w(L-1) = 0$$

$$R_2 w(0) + R_3 w(1) + \dots + R_L w(L-2) = 0$$

$$\vdots \quad = \quad \vdots$$

$$R_{L-1} w(0) + R_L w(1) = 0$$

$$R_L w(0) = 0$$

The shift-and-cut map

$$\begin{aligned}\sigma_+ : \mathbb{R}[\xi] &\rightarrow \mathbb{R}[\xi] \\ \sigma_+(\sum_{i=0}^n p_i \xi^i) &:= \sum_{i=0}^{n-1} p_{i+1} \xi^i\end{aligned}$$

“Divide by ξ and take polynomial part”

Extended componentwise to vectors and matrices

Example

$$R(\xi) = R_0 + R_1\xi + \dots + R_{L-1}\xi^{L-1} + R_L\xi^L$$

Example

$$R(\xi) = R_0 + R_1\xi + \dots + R_{L-1}\xi^{L-1} + R_L\xi^L$$

$$\sigma_+(R(\xi)) = R_1 + \dots + R_{L-1}\xi^{L-2} + R_L\xi^{L-1}$$

Example

$$R(\xi) = R_0 + R_1\xi + \dots + R_{L-1}\xi^{L-1} + R_L\xi^L$$

$$\sigma_+(R(\xi)) = R_1 + \dots + R_{L-1}\xi^{L-2} + R_L\xi^{L-1}$$

$$\sigma_+^2(R(\xi)) = R_2 + \dots + R_{L-1}\xi^{L-3} + R_L\xi^{L-2}$$

Example

$$R(\xi) = R_0 + R_1\xi + \dots + R_{L-1}\xi^{L-1} + R_L\xi^L$$

$$\sigma_+(R(\xi)) = R_1 + \dots + R_{L-1}\xi^{L-2} + R_L\xi^{L-1}$$

$$\sigma_+^2(R(\xi)) = R_2 + \dots + R_{L-1}\xi^{L-3} + R_L\xi^{L-2}$$

$\vdots = \vdots$

Example

$$R(\xi) = R_0 + R_1\xi + \dots + R_{L-1}\xi^{L-1} + R_L\xi^L$$

$$\sigma_+(R(\xi)) = R_1 + \dots + R_{L-1}\xi^{L-2} + R_L\xi^{L-1}$$

$$\sigma_+^2(R(\xi)) = R_2 + \dots + R_{L-1}\xi^{L-3} + R_L\xi^{L-2}$$

$$\vdots = \vdots$$

$$\sigma_+^L(R(\xi)) = R_L$$

Shift-and-cut and concatenability with zero

**w is
concatenable
with zero**

\Leftrightarrow

$$(\sigma_+(R)(\sigma)w)(0) = 0$$

$$(\sigma_+^2(R)(\sigma)w)(0) = 0$$

$$\vdots = \vdots$$

$$(\sigma_+^L(R)(\sigma)w)(0) = 0$$

$\text{col}((\sigma_+^i(R))_{i=1,\dots,L}(\sigma))$ is a state map!

Shift-and-cut and concatenability with zero

$$\begin{array}{l} \mathbf{w \text{ is}} \\ \mathbf{concatenable} \\ \mathbf{with zero} \end{array} \Leftrightarrow \begin{array}{l} (\sigma_+(R)(\sigma)w)(0) = 0 \\ (\sigma_+^2(R)(\sigma)w)(0) = 0 \\ \vdots = \vdots \\ (\sigma_+^L(R)(\sigma)w)(0) = 0 \end{array}$$

$\text{col}((\sigma_+^i(R))_{i=1,\dots,L}(\sigma))$ is a state map!

**Other equations equivalent to shift-and-cut ones
 \implies different state maps are possible!**

Shift-and-cut and concatenability with zero

**w is
concatenable
with zero**

\Leftrightarrow

$$\begin{aligned}(\sigma_+(R)(\sigma)w)(0) &= 0 \\(\sigma_+^2(R)(\sigma)w)(0) &= 0 \\&\vdots \\(\sigma_+^L(R)(\sigma)w)(0) &= 0\end{aligned}$$

$\text{col}((\sigma_+^i(R))_{i=1,\dots,L}(\sigma))$ is a state map!

Other equations equivalent to **shift-and-cut ones**
 \implies different state maps are possible!

Example: scalar systems

$$r_0 W + r_1 \sigma W + \dots + \sigma^n W = 0$$

Example: scalar systems

$$r_0 w + r_1 \sigma w + \dots + \sigma^n w = 0$$

Observe w concatenable with zero iff $w = 0$. Indeed,

$$\begin{aligned}\sigma_+^n(r)(\sigma)w &= w \\ \sigma_+^{n-1}(r)(\sigma)w &= r_{n-1}w + \sigma w \\ &\vdots = \vdots \\ \sigma_+(r)(\sigma)w &= r_1 w + \dots + \sigma^{n-1}w\end{aligned}$$

Example: scalar systems

$$r_0 w + r_1 \sigma w + \dots + \sigma^n w = 0$$

Observe w concatenable with zero iff $w = 0$. Indeed,

$$\begin{aligned}\sigma_+^n(r)(\sigma)w &= w \\ \sigma_+^{n-1}(r)(\sigma)w &= r_{n-1}w + \sigma w \\ &\vdots = \vdots \\ \sigma_+(r)(\sigma)w &= r_1 w + \dots + \sigma^{n-1}w\end{aligned}$$

Zero at $t = 0$ iff $(\sigma^k w)(0) = 0$ for $k = 0, \dots, n - 1$.

From kernel representation to state map

Denote $\text{col}((\sigma_+^i(R)))_{i=1,\dots,L} =: \Sigma_R$.

Theorem: Let $\mathfrak{B} = \ker(R(\sigma))$. Then

$$\begin{aligned} R(\sigma)w &= 0 \\ \Sigma_R(\sigma)w &= x \end{aligned}$$

is a **state representation** of \mathfrak{B} with **state variable** x .

Algebraic characterization

Theorem: Let $\mathfrak{B} = \ker(R(\sigma))$, and define Σ_R as above. Then

$$\Xi_R := \{f \in \mathbb{R}^{1 \times w}[\xi] \mid \exists g \in \mathbb{R}^{1 \times \bullet}[\xi], \alpha \in \mathbb{R}^{1 \times \bullet} \\ \text{s.t. } f = \alpha \Sigma_R + gR\}$$

is a vector space over \mathbb{R} .

Algebraic characterization

Theorem: Let $\mathfrak{B} = \ker(R(\sigma))$, and define Σ_R as above. Then

$$\Xi_R := \{f \in \mathbb{R}^{1 \times w}[\xi] \mid \exists g \in \mathbb{R}^{1 \times \bullet}[\xi], \alpha \in \mathbb{R}^{1 \times \bullet} \\ \text{s.t. } f = \alpha \Sigma_R + gR\}$$

is a vector space over \mathbb{R} .

$X \in \mathbb{R}^{\bullet \times w}[\xi]$ is state map for \mathfrak{B} iff $\text{row span}(X) = \Xi_R$

Algebraic characterization

Theorem: Let $\mathfrak{B} = \ker(R(\sigma))$, and define Σ_R as above. Then

$$\Xi_R := \{f \in \mathbb{R}^{1 \times w}[\xi] \mid \exists g \in \mathbb{R}^{1 \times \bullet}[\xi], \alpha \in \mathbb{R}^{1 \times \bullet} \text{ s.t. } f = \alpha \Sigma_R + gR\}$$

is a vector space over \mathbb{R} .

$X \in \mathbb{R}^{\bullet \times w}[\xi]$ is state map for \mathfrak{B} iff $\text{row span}(X) = \Xi_R$

X is **minimal** if and only if its rows are a **basis** for Ξ_R .

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

If $(y, u) \in \mathfrak{B}$, then for all $g \in \mathbb{R}[\xi]$

$$\begin{aligned} [\sigma + 2 \quad -1] \begin{bmatrix} y \\ u \end{bmatrix} &= [\sigma + 2 \quad -1] \begin{bmatrix} y \\ u \end{bmatrix} \\ &+ \underbrace{g(\sigma) [\sigma^2 + 2\sigma + 3 \quad -\sigma - 3]}_{=0 \text{ on } \mathfrak{B}} \begin{bmatrix} y \\ u \end{bmatrix} \end{aligned}$$

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

If $(y, u) \in \mathfrak{B}$, then for all $g \in \mathbb{R}[\xi]$

$$\begin{aligned} [\sigma + 2 \quad -1] \begin{bmatrix} y \\ u \end{bmatrix} &= [\sigma + 2 \quad -1] \begin{bmatrix} y \\ u \end{bmatrix} \\ &+ \underbrace{g(\sigma) [\sigma^2 + 2\sigma + 3 \quad -\sigma - 3]}_{=0 \text{ on } \mathfrak{B}} \begin{bmatrix} y \\ u \end{bmatrix} \end{aligned}$$

‘equivalence modulo R ’

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

$$\sigma_+^2 \rightsquigarrow [1 \quad 0] \rightsquigarrow [1 \quad 0]$$

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

$$\sigma_+^2 \rightsquigarrow [1 \quad 0] \rightsquigarrow [1 \quad 0]$$

If $(y, u) \in \mathfrak{B}$, then for all $g \in \mathbb{R}[\xi]$

$$\begin{aligned} [1 \quad 0] \begin{bmatrix} y \\ u \end{bmatrix} &= [1 \quad 0] \begin{bmatrix} y \\ u \end{bmatrix} \\ &+ \underbrace{g(\sigma) [\sigma^2 + 2\sigma + 3 \quad -\sigma - 3]}_{=0 \text{ on } \mathfrak{B}} \begin{bmatrix} y \\ u \end{bmatrix} \end{aligned}$$

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

$$\sigma_+^2 \rightsquigarrow [1 \quad 0] \rightsquigarrow [1 \quad 0]$$

$$\Xi_R = \{ \alpha [\xi + 2 \quad -1] + g(\xi) [\xi^2 + 2\xi + 3 \quad -\xi - 3], \\ \beta [1 \quad 0] + f(\xi) [\xi^2 + 2\xi + 3 \quad -\xi - 3] \\ \alpha, \beta \in \mathbb{R}, f, g \in \mathbb{R}[\xi] \}$$

Example

$$(\sigma^2 + 2\sigma + 3)y = (\sigma + 3)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

$$\sigma_+ \rightsquigarrow [\xi + 2 \quad -1] \rightsquigarrow [\sigma + 2 \quad -1]$$

$$\sigma_+^2 \rightsquigarrow [1 \quad 0] \rightsquigarrow [1 \quad 0]$$

$$\Xi_R = \{ \alpha [\xi + 2 \quad -1] + g(\xi) [\xi^2 + 2\xi + 3 \quad -\xi - 3], \\ \beta [1 \quad 0] + f(\xi) [\xi^2 + 2\xi + 3 \quad -\xi - 3] \\ \alpha, \beta \in \mathbb{R}, f, g \in \mathbb{R}[\xi] \}$$

Any set of generators of $\Xi_R \rightsquigarrow$ a state map

Outline

The axiom of state

Discrete-time systems

First-order representations

State maps

The shift-and-cut map

Algebraic characterization

Continuous-time systems

Computation of state-space representations

On the space of solutions

\mathcal{C}^∞ -solutions to $R(\frac{d}{dt})w = 0$ too small $\rightsquigarrow \mathcal{L}_1^{\text{loc}}$

Equality in the sense of distributions:

$$R\left(\frac{d}{dt}\right)w = 0 \quad \Leftrightarrow \quad \int_{-\infty}^{+\infty} w(t)^\top \left(R\left(-\frac{d}{dt}\right)^\top f\right)(t) dt = 0$$

for all testing functions f .

The axiom of state revisited

$\Sigma = (\mathbb{T}, \mathbb{W}, \mathbb{X}, \mathfrak{B}_{\text{full}})$ is a state system if

$(w_1, x_1), (w_2, x_2) \in \mathfrak{B}_{\text{full}}$ and $x_1(T) = x_2(T)$
and x_1, x_2 continuous at T

\Downarrow

$(w_1, x_1) \underset{T}{\wedge} (w_2, x_2) \in \mathfrak{B}_{\text{full}}$

‘State map’ $\rightsquigarrow X\left(\frac{d}{dt}\right)$

From kernel representation to state map

Denote $\text{col}((\sigma_+^i(R)))_{i=1,\dots,L} =: \Sigma_R$.

Theorem: Let $\mathfrak{B} = \ker(R(\frac{d}{dt}))$. Then

$$\begin{aligned} R\left(\frac{d}{dt}\right)w &= 0 \\ \Sigma_R\left(\frac{d}{dt}\right)w &= x \end{aligned}$$

is a state representation of \mathfrak{B} with state variable x .

¿How to prove it?

When is $w \in \mathfrak{B}$ concatenable with zero?

$$\begin{aligned} \mathbf{0} \underset{\mathbf{0}}{\wedge} w \in \mathfrak{B} &\iff \int_{-\infty}^{+\infty} (\mathbf{0} \underset{\mathbf{0}}{\wedge} w)(t)^\top (R(-\frac{d}{dt})^\top f)(t) dt = 0 \\ &\iff \int_0^{+\infty} w(t)^\top (R(-\frac{d}{dt})^\top f)(t) dt = 0 \end{aligned}$$

for all testing functions f

When is $w \in \mathfrak{B}$ concatenable with zero?

$$\begin{aligned} \mathbf{0} \underset{0}{\wedge} w \in \mathfrak{B} &\iff \int_{-\infty}^{+\infty} (\mathbf{0} \underset{0}{\wedge} w)(t)^\top (R(-\frac{d}{dt})^\top f)(t) dt = 0 \\ &\iff \int_0^{+\infty} w(t)^\top (R(-\frac{d}{dt})^\top f)(t) dt = 0 \end{aligned}$$

for all testing functions f

Integrating repeatedly by parts on f yields:

$$\begin{aligned} \sum_{k=1}^{\deg(R)} \sum_{j=k}^{\deg(R)} (-1)^{k-1} \left(\frac{d^{j-k}}{dt^{j-k}} w\right)(0)^\top R_j^\top \left(\frac{d^{k-1}}{dt^{k-1}} f\right)(0) \\ + \int_0^{+\infty} \left(R\left(\frac{d}{dt}\right)w\right)(t)^\top f(t) dt = 0 \end{aligned}$$

When is $w \in \mathfrak{B}$ concatenable with zero?

$$\begin{aligned} \mathbf{0} \underset{0}{\wedge} w \in \mathfrak{B} &\iff \int_{-\infty}^{+\infty} (\mathbf{0} \underset{0}{\wedge} w)(t)^\top (R(-\frac{d}{dt})^\top f)(t) dt = 0 \\ &\iff \int_0^{+\infty} w(t)^\top (R(-\frac{d}{dt})^\top f)(t) dt = 0 \end{aligned}$$

for all testing functions f

Integrating repeatedly by parts on f yields:

$$\begin{aligned} \sum_{k=1}^{\deg(R)} \sum_{j=k}^{\deg(R)} (-1)^{k-1} \left(\frac{d^{j-k}}{dt^{j-k}} w \right)(0)^\top R_j^\top \left(\frac{d^{k-1}}{dt^{k-1}} f \right)(0) \\ + \int_0^{+\infty} \left(R \left(\frac{d}{dt} \right) w \right)(t)^\top f(t) dt = 0 \end{aligned}$$

$w \in \mathfrak{B}$ concatenable with zero if and only if...

$$\sum_{k=1}^{\deg(R)} \sum_{j=k}^{\deg(R)} (-1)^{k-1} \left(\frac{d^{j-k}}{dt^{j-k}} w \right)(0)^\top R_j^\top \left(\frac{d^{k-1}}{dt^{k-1}} f \right)(0) = 0$$

$$\Updownarrow$$

$$\left[\begin{array}{c} f(0) \\ \left(\frac{d}{dt} f \right)(0) \\ \vdots \\ (-1)^{\deg(R)-1} \left(\frac{d^{\deg(R)-1}}{dt^{\deg(R)-1}} f \right)(0) \end{array} \right]^\top (\Sigma_R \left(\frac{d}{dt} \right) w)(0) = 0$$

$$\Updownarrow$$

$$(\Sigma_R \left(\frac{d}{dt} \right) w)(0) = 0$$

The **shift-and-cut** state map!

Outline

The axiom of state

Discrete-time systems

First-order representations

State maps

The shift-and-cut map

Algebraic characterization

Continuous-time systems

Computation of state-space representations

From kernel representation to state representation

$$R \in \mathbb{R}^{g \times w}[\xi] \rightsquigarrow \text{state map } X \in \mathbb{R}^{n \times w}[\xi]$$

Find:

$$E, F \in \mathbb{R}^{(n+g) \times n}, G \in \mathbb{R}^{(n+g) \times w}$$

$$T \in \mathbb{R}^{(n+g) \times g}[\xi] \text{ with } \text{rank}(T(\lambda)) = g \quad \forall \lambda \in \mathbb{C}$$

satisfying

$$E\xi X(\xi) + FX(\xi) + G = T(\xi)R(\xi)$$

Linear equations, Gröbner bases computations!

From I/O representation to I/O/S representation

I/O representation

$$R = \begin{bmatrix} P & -Q \end{bmatrix}$$



state map

$$\begin{bmatrix} X_y & X_u \end{bmatrix}$$

Find:

$$A \in \mathbb{R}^{n \times n}, B \in \mathbb{R}^{n \times m}, C \in \mathbb{R}^{p \times p}, D \in \mathbb{R}^{p \times m}$$

$$T \in \mathbb{R}^{(n+p) \times p}[\xi] \text{ with } \text{rank}(T(\lambda)) = g \quad \forall \lambda \in \mathbb{C}$$

satisfying

$$\begin{bmatrix} \xi X_y(\xi) & \xi X_u(\xi) \\ I_p & 0 \end{bmatrix} = \begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} X_y(\xi) & X_u(\xi) \\ 0 & I_m \end{bmatrix} + T(\xi)R(\xi)$$

On the choice of state map

State map
+
system equations



state-space
equations

On the choice of state map

State map
+
system equations



**state-space
equations**

On the choice of state map

**State map
+
system equations**



**state-space
equations**

$$\left(\frac{d^2}{dt^2} + 2\frac{d}{dt} + 3\right)y = \left(\frac{d}{dt} + 3\right)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

On the choice of state map

State map
+
system equations

\rightsquigarrow

state-space
equations

$$\left(\frac{d^2}{dt^2} + 2\frac{d}{dt} + 3\right)y = \left(\frac{d}{dt} + 3\right)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

Take $X(\xi) = \begin{bmatrix} 1 & 0 \\ \xi + 2 & -1 \end{bmatrix}$ ('reverse shift-and-cut').

Then

$$A = \begin{bmatrix} -2 & 1 \\ -3 & 0 \end{bmatrix} \quad B = \begin{bmatrix} -1 \\ -3 \end{bmatrix}$$
$$C = [1 \quad 0] \quad D = [0]$$

'observer canonical form'

On the choice of state map

State map
+
system equations

\rightsquigarrow

state-space
equations

$$\left(\frac{d^2}{dt^2} + 2\frac{d}{dt} + 3\right)y = \left(\frac{d}{dt} + 3\right)u$$

$$[\xi^2 + 2\xi + 3 \quad -\xi - 3]$$

Take $X(\xi) = \begin{bmatrix} 1 & 0 \\ \xi & -1 \end{bmatrix}$. Then

$$\begin{aligned} A &= \begin{bmatrix} 0 & 1 \\ -3 & -2 \end{bmatrix} & B &= \begin{bmatrix} 1 \\ 1 \end{bmatrix} \\ C &= [1 \quad 0] & D &= [0] \end{aligned}$$

'observable canonical form'

Summary

Summary

- **The state is constructed!**

Summary

- **The state is constructed!**
- **Axiom of state**

Summary

- **The state is constructed!**
- **Axiom of state**
- **Concatenability with zero**

Summary

- **The state is constructed!**
- **Axiom of state**
- **Concatenability with zero**
- **State maps**

Summary

- **The state is constructed!**
- **Axiom of state**
- **Concatenability with zero**
- **State maps**
- **State maps \rightsquigarrow state-space equations**

Summary

- **The state is constructed!**
- **Axiom of state**
- **Concatenability with zero**
- **State maps**
- **State maps \rightsquigarrow state-space equations**
- **Algorithms!**