The CANONICAL CONTROLLER and its REGULARITY

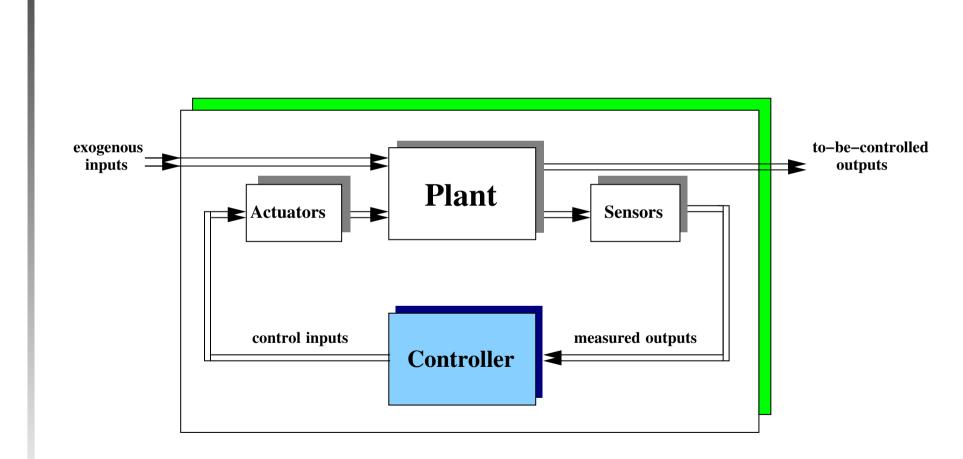


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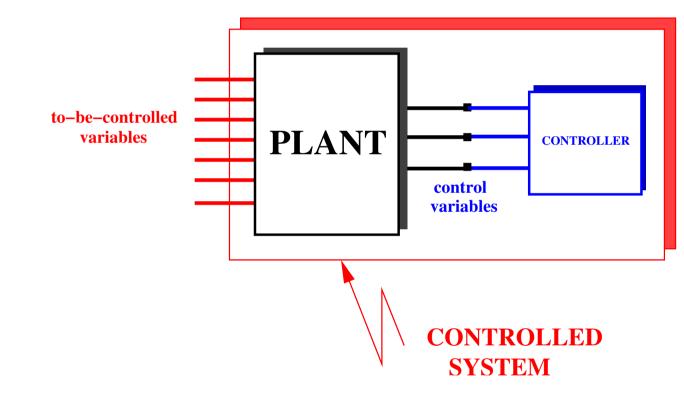
CDC2003, Maui

December 10, 2003

Feedback control

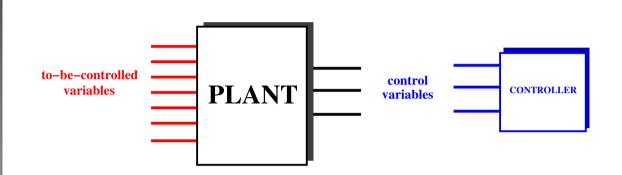


Behavioral control



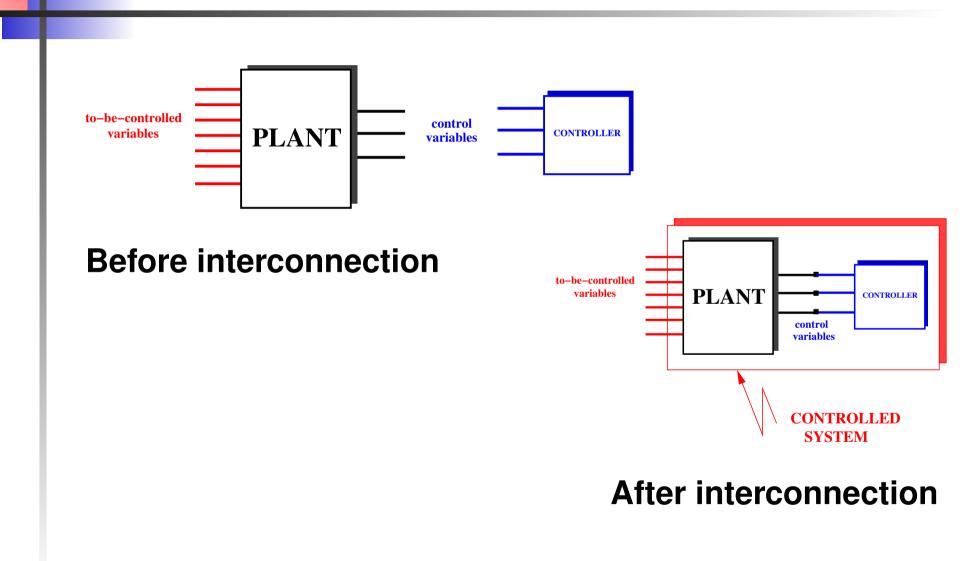
Control as interconnection

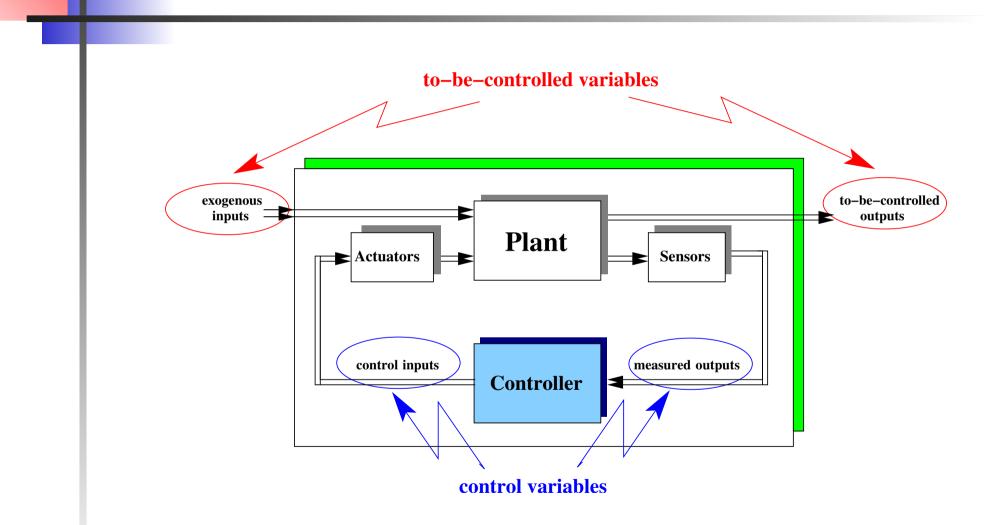
Behavioral control



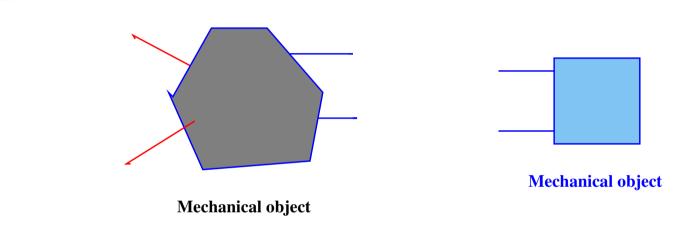
Before interconnection

Behavioral control

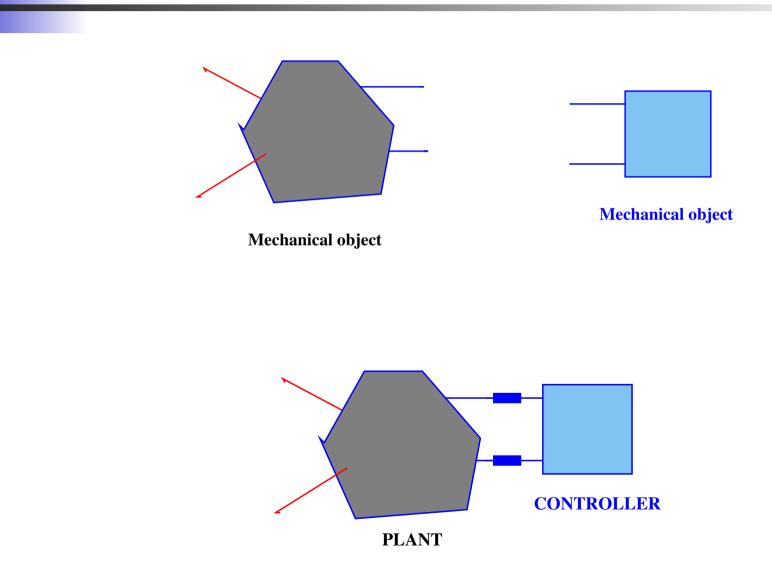












Mathematization

Domain of the to-be-controlled variables: ₩ Domain of the control variables: ℂ Typically: families of time-signals

Mathematization

Full plant behavior:

$$\mathcal{P}_{\mathrm{full}} = \{(oldsymbol{w}, oldsymbol{c}) \in \mathbb{W} imes \mathbb{C} \mid \mathsf{allowed} \ \mathsf{by} \ \mathsf{plant} \ \mathsf{laws} \}$$

 $\mathcal{C} = \{ \mathbf{c} \in \mathbb{C} \mid \text{allowed by controller laws} \}$



Full plant behavior:

$$\mathcal{P}_{\text{full}} = \{(w, c) \in \mathbb{W} \times \mathbb{C} \mid \text{allowed by plant laws} \}$$

 $\mathcal{C} = \{ \mathbf{c} \in \mathbb{C} \mid \text{allowed by controller laws} \}$

Controlled behavior:

 $\mathcal{K} := \{ oldsymbol{w} \in \mathbb{W} \mid \exists oldsymbol{c} \in \mathbb{C} \ ext{ such that } (oldsymbol{w}, oldsymbol{c}) \in \mathcal{P}_{ ext{full}} ext{ and } oldsymbol{c} \in \mathcal{C} \}.$



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and that \mathcal{K} is implementable



Controlled behavior:

We say that \mathcal{C} implements \mathcal{K} ,

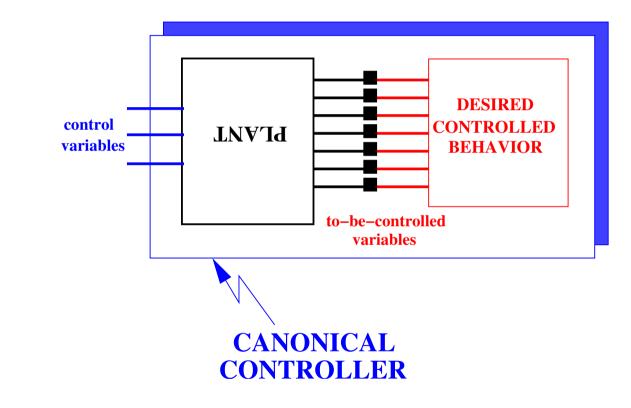
and that \mathcal{K} is implementable

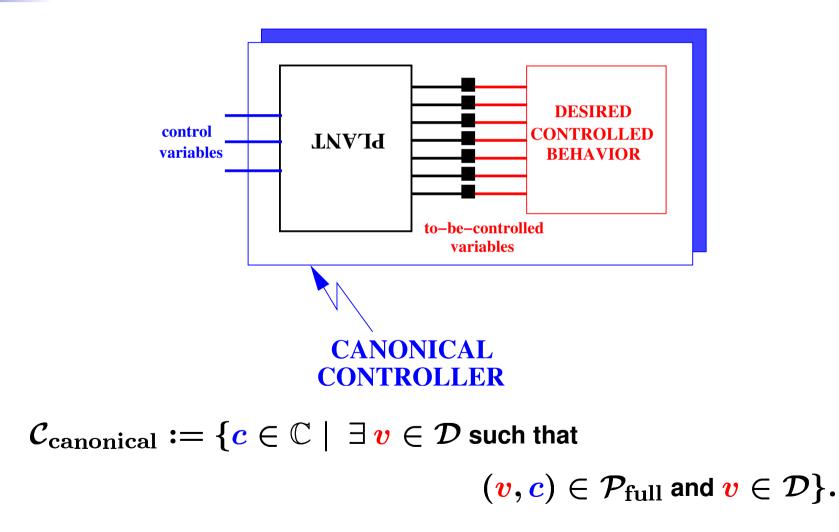
Questions:

Which C implements the

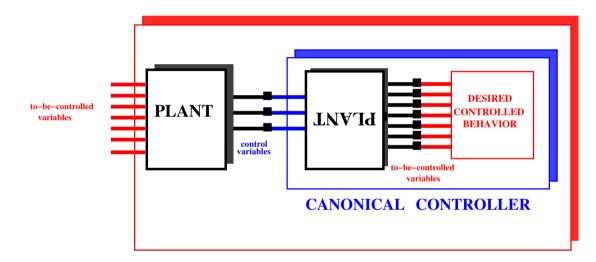
desired controlled behavior \mathcal{D} ?

Given \mathcal{P}_{full} , which $\mathcal{K} \subseteq \mathbb{W}$ are implementable?

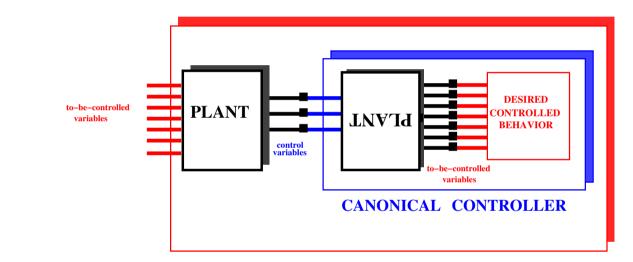




The canonically controlled system:

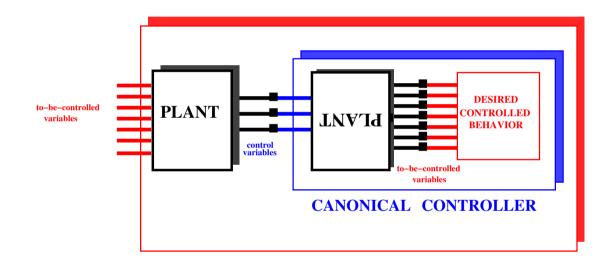


The canonically controlled system:



This is the **internal model principle** at work!

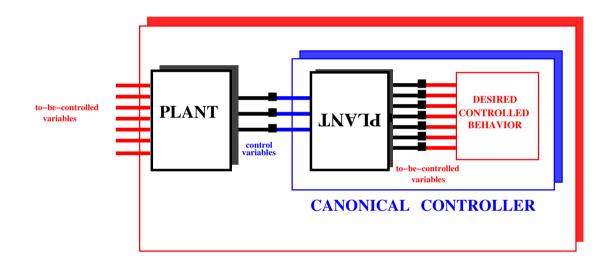
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Theorem:

 \mathcal{D} is implementable $\Leftrightarrow \mathcal{C}_{\mathrm{canonical}}$ implements it.

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Does $C_{canonical}$ have good properties?

LTIS

We henceforth restrict attention to

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The *behavior* \mathfrak{B} belongs to $\mathfrak{L}^{ imes}$: \Leftrightarrow \exists a polynomial matrix $R \in \mathbb{R}^{ullet imes imes}[\xi]$ such that

 $\mathfrak{B} = \{w \in \mathfrak{C}^\infty(\mathbb{R},\mathbb{R}^{\scriptscriptstyle {\mathbb W}}) \mid R(rac{d}{dt})w = 0\}$.

Control of LTIS

Plant:

$$\mathcal{P}_{\mathrm{full}} \in \mathfrak{L}^{w+c}.$$

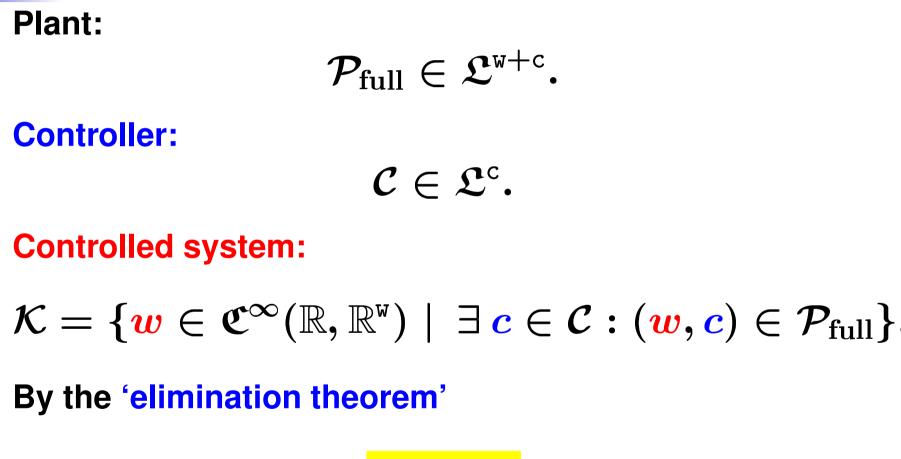
Controller:

 $\mathcal{C} \in \mathfrak{L}^{c}$.

Controlled system:

 $\mathcal{K} = \{ oldsymbol{w} \in \mathfrak{C}^\infty(\mathbb{R},\mathbb{R}^{w}) \mid \exists oldsymbol{c} \in \mathcal{C} : (oldsymbol{w},oldsymbol{c}) \in \mathcal{P}_{ ext{full}} \}$

Control of LTIS



$$\mathcal{K}\in\mathfrak{L}^{\scriptscriptstyle{W}}$$

Which behaviors $\mathcal{K} \in \mathfrak{L}^{w}$ can be implemented by attaching a controller $\mathcal{C} \in \mathfrak{L}^{c}$ to a given plant $\mathcal{P}_{full} \in \mathfrak{L}^{w+c}$?

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This question has a very concrete and intuitive answer. Theorem: Let $\mathcal{P}_{\mathrm{full}} \in \mathfrak{L}^{w+c}$ be given.

The behavior $\mathcal{K}\in\mathfrak{L}^{\scriptscriptstyle W}$ is implementable if and only if

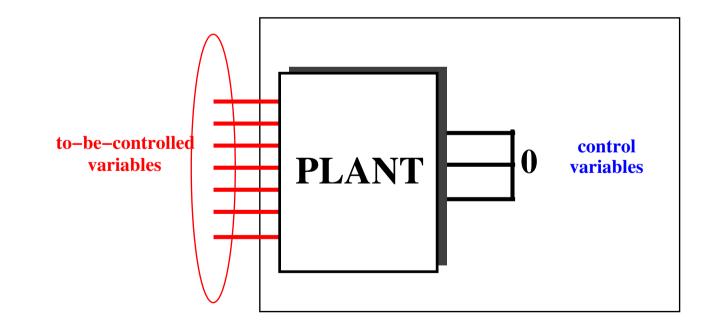
$$\mathcal{N} \subseteq \mathcal{K} \subseteq \mathcal{P}$$

The behavior $\mathcal{K}\in\mathfrak{L}^{w}$ is implementable if and only if $\mathcal{N}\subseteq\mathcal{K}\subseteq\mathcal{P}$

where $\mathcal{N} \in \mathfrak{L}^{\mathbb{W}}$ is the *hidden behavior* defined by $\mathcal{N} := \{ w \in \mathfrak{C}^{\infty}(\mathbb{R}, \mathbb{R}^{\mathbb{W}}) \mid (w, 0) \in \mathcal{P}_{\text{full}} \},\$ and \mathcal{P} is the *manifest plant behavior* defined by

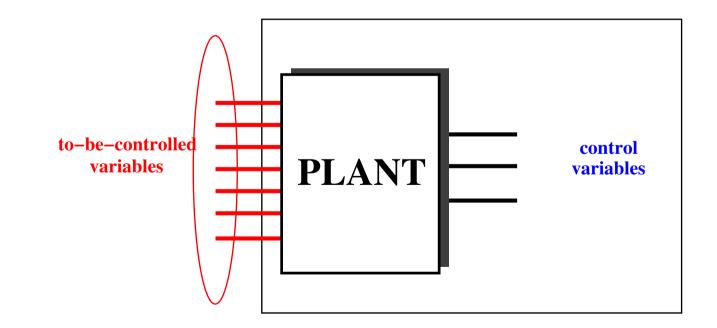
 $\mathcal{P} := \{ w \in \mathfrak{C}^{\infty}(\mathbb{R}, \mathbb{R}^{w}) \mid \exists \ c : (w, c) \in \mathcal{P}_{\mathrm{full}} \}.$

$\mathcal{N} \in \mathfrak{L}^{W}$, the *hidden behavior*





$\mathcal{P} \in \mathfrak{L}^{W}$, the *manifest plant behavior*



The effect of the canonical controller

Theorem: Consider $\mathcal{P}_{full} \in \mathfrak{L}^{w+c}$ and $\mathcal{D} \in \mathfrak{L}^{w}$. The controlled behavior implemented by the associated canonical controller $\mathcal{C}_{canonical} \in \mathfrak{L}^{c}$ is

$$\mathcal{K} = \mathcal{N} + \mathcal{D} \cap \mathcal{P}$$

with \mathcal{N} the hidden and \mathcal{P} the manifest plant behavior.

The effect of the canonical controller

Theorem: Consider $\mathcal{P}_{full} \in \mathfrak{L}^{w+c}$ and $\mathcal{D} \in \mathfrak{L}^{w}$. The controlled behavior implemented by the associated canonical controller $\mathcal{C}_{canonical} \in \mathfrak{L}^{c}$ is

$$\mathcal{K} = \mathcal{N} + \mathcal{D} \cap \mathcal{P}$$

with \mathcal{N} the hidden and \mathcal{P} the manifest plant behavior.

Corollary: The canonical controller implements $\mathcal{D} \in \mathfrak{L}^{\mathbb{W}}$ if and only if $\mathcal{N} \subseteq \mathcal{D} \subseteq \mathcal{P}$ i.e. if and only if \mathcal{D} is implementable.

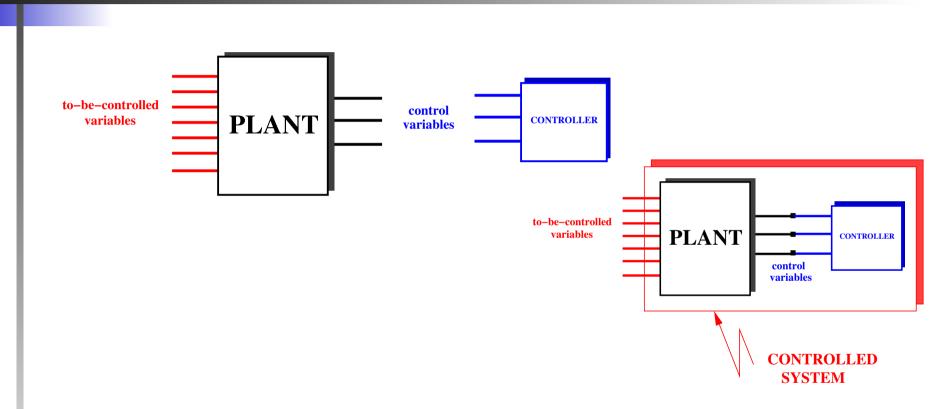
The full controlled behavior $\mathcal{K}_{\mathrm{full}} \subseteq \mathcal{P}_{\mathrm{full}}$ is defined by $\mathcal{K}_{\mathrm{full}} := \{(w, c) \in \mathcal{P}_{\mathrm{full}} \mid c \in \mathcal{C}\}.$

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The controller $\mathcal{C}\in\mathfrak{L}^{c}$ is said to be regular if

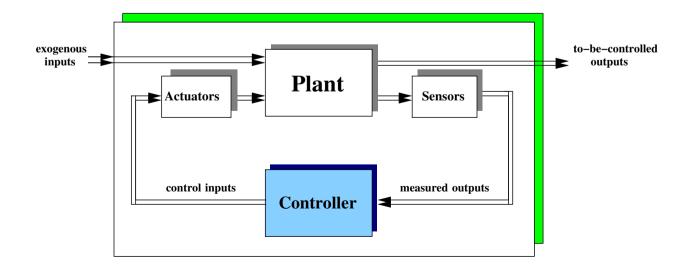
 $p(\mathcal{K}_{full}) = p(\mathcal{P}_{full}) + p(\mathcal{C})$

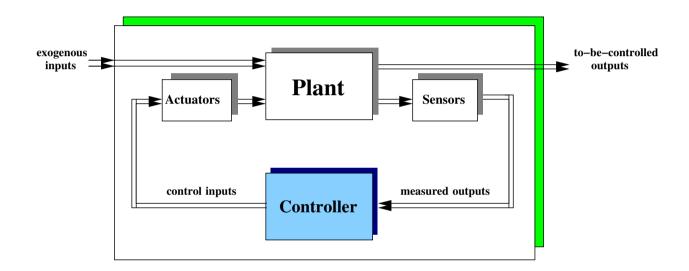


Regularity :=

if the controller has p bound (i.e. output) variables, then the plant looses p free variables after interconnection.

A controller is regular if and only if it can be realized as a feedback controller with a (possibly non-proper) transfer function from an output to an input in \mathcal{P}_{full} for an input/output partition of c.





 \Rightarrow A controller is regular if and only if it can be viewed as an 'intelligent controller' that processes sensor inputs outputs into actuator inputs.

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If \mathcal{P} is controllable, then every implementable \mathcal{K} is regularly implementable.

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The canonical controller is regular.

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The canonical controller is regular.

 \Rightarrow The canonical controller is regular if and only if every controller is regular.

Hence the canonical controller is maximally irregular.

