

DSP-CIS

Part-IV : Filter Banks & Subband Systems

Chapter-13 : Frequency Domain Filtering

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Part-IV : Filter Banks & Subband Systems

Chapter-11 Filter Bank Preliminaries

Chapter-12 Filter Bank Design

Chapter-13 Frequency Domain Filtering

- Frequency Domain FIR Filter Realization
- Frequency Domain Adaptive Filtering

Chapter-14 Time-Frequency Analysis & Scaling

FIR Filter Realization

FIR Filter Realization

=Construct (realize) LTI system (with delay elements, adders and multipliers), such that the I/O behavior is given by..

$$y[k] = b_0 \cdot u[k] + b_1 \cdot u[k-1] + \dots + b_{L-1} \cdot u[k-L+1]$$

Several possibilities exist...

1. Direct form
2. Transposed direct form
3. Lattice realization (LPC lattice)
4. Lossless lattice realization
5. Frequency domain realization: see Part IV ←

For convenience, number of filter coefficients is now L, instead of L+1 (...easier formulas)

Frequency Domain FIR Filter Realization

Have to know a theorem from linear algebra here:

- A 'circulant' matrix is a matrix where each row is obtained from the previous row using a right-shift (by 1 position), the rightmost element which spills over is circulated back to become the leftmost element

$$\begin{bmatrix} a & d & c & b \\ b & a & d & c \\ c & b & a & d \\ d & c & b & a \end{bmatrix}$$

- The eigenvalue decomposition of a circulant matrix is always given as... (4x4 example)

$$\begin{bmatrix} a & d & c & b \\ b & a & d & c \\ c & b & a & d \\ d & c & b & a \end{bmatrix} = F^{-1} \cdot \begin{bmatrix} A & 0 & 0 & 0 \\ 0 & B & 0 & 0 \\ 0 & 0 & C & 0 \\ 0 & 0 & 0 & D \end{bmatrix} \cdot F, \quad \text{with} \quad \begin{bmatrix} A \\ B \\ C \\ D \end{bmatrix} = F \cdot \begin{bmatrix} a \\ b \\ c \\ d \end{bmatrix}$$

with F the DFT-matrix. This means that the eigenvectors are equal to the column-vectors of the IDFT-matrix, and that then eigenvalues are obtained as the DFT of the first column of the circulant matrix (proof by Matlab)

Frequency Domain FIR Filter Realization

FIR Filter Realization (example L=4, similar for other L)

$$y[k] = b_0 \cdot u[k] + b_1 \cdot u[k-1] + b_2 \cdot u[k-2] + b_3 \cdot u[k-3]$$

$$B(z) = b_0 + b_1 z^{-1} + b_2 z^{-2} + b_3 z^{-3}$$

Consider a 'block processing' where a block of L_B output samples are computed at once, with 'block length' $L_B=L$:

$$\begin{bmatrix} y[k] \\ y[k-1] \\ y[k-2] \\ y[k-3] \end{bmatrix} = \begin{bmatrix} b_0 & b_1 & b_2 & b_3 & 0 & 0 & 0 & 0 \\ 0 & b_0 & b_1 & b_2 & b_3 & 0 & 0 & 0 \\ 0 & 0 & b_0 & b_1 & b_2 & b_3 & 0 & 0 \\ 0 & 0 & 0 & b_0 & b_1 & b_2 & b_3 & 0 \end{bmatrix} \cdot \begin{bmatrix} u[k] \\ u[k-1] \\ u[k-2] \\ u[k-3] \\ u[k-4] \\ u[k-5] \\ u[k-6] \\ u[k-7] \end{bmatrix}$$

Frequency Domain FIR Filter Realization

Now some matrix manipulation...

$$\begin{bmatrix} y[k] \\ y[k-1] \\ y[k-2] \\ y[k-3] \end{bmatrix} = \begin{bmatrix} 0_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot \begin{matrix} \text{=circulant matrix} \\ \begin{bmatrix} 0 & 0 & 0 & 0 & b_0 & b_1 & b_2 & b_3 \\ b_3 & 0 & 0 & 0 & 0 & b_0 & b_1 & b_2 \\ b_2 & b_3 & 0 & 0 & 0 & 0 & b_0 & b_1 \\ b_1 & b_2 & b_3 & 0 & 0 & 0 & 0 & b_0 \\ b_0 & b_1 & b_2 & b_3 & 0 & 0 & 0 & 0 \\ 0 & b_0 & b_1 & b_2 & b_3 & 0 & 0 & 0 \\ 0 & 0 & b_0 & b_1 & b_2 & b_3 & 0 & 0 \\ 0 & 0 & 0 & b_0 & b_1 & b_2 & b_3 & 0 \end{bmatrix} \cdot \begin{bmatrix} u[k] \\ u[k-1] \\ u[k-2] \\ u[k-3] \\ u[k-4] \\ u[k-5] \\ u[k-6] \\ u[k-7] \end{bmatrix} \end{matrix}$$

$$= \begin{bmatrix} 0_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot F^{-1} \cdot \begin{matrix} \begin{bmatrix} B_0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & B_1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & B_2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & B_3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & B_4 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & B_5 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & B_6 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & B_7 \end{bmatrix} \cdot F \cdot \begin{bmatrix} u[k] \\ u[k-1] \\ u[k-2] \\ u[k-3] \\ u[k-4] \\ u[k-5] \\ u[k-6] \\ u[k-7] \end{bmatrix} \end{matrix}$$

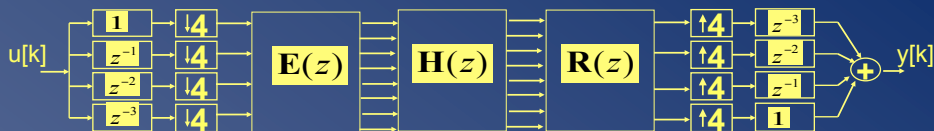
Frequency Domain FIR Filter Realization

- This means that a block of $L_B=L$ output samples can be computed as follows (read previous formula from right to left) :
 - Compute DFT of 2L input samples, i.e. last L samples combined ('overlapped') with previous L samples
 - Perform component-wise multiplication with... (=freq.domain representation of the FIR filter)
 - Compute IDFT
 - Throw away 1st half of result, select ('save') 2nd half
- This is referred to as an 'overlap-save' procedure (and 'frequency domain filter realization' because of the DFT/IDFT)

$$\begin{bmatrix} B_0 \\ B_1 \\ B_2 \\ B_3 \\ B_4 \\ B_5 \\ B_6 \\ B_7 \end{bmatrix} = F \cdot \begin{bmatrix} 0 \\ b_3 \\ b_2 \\ b_1 \\ b_0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

Frequency Domain FIR Filter Realization

- This corresponds to a filter bank-type realization as follows...



Analysis bank:

$$\mathbf{E}(z) = F \cdot \begin{bmatrix} I_{4 \times 4} \\ z^{-1} \cdot I_{4 \times 4} \end{bmatrix}$$

Subband processing:

$$\mathbf{H}(z) = \text{diag}\{B_0, B_1, \dots, B_7\}$$

Synthesis bank:

$$\mathbf{R}(z) = \begin{bmatrix} 0 & I_{4 \times 4} \end{bmatrix} F^{-1}$$

This is a 2L-channel filter bank, with L-fold downsampling

The analysis FB is a 2L-channel uniform DFT filter bank (see Chapter 11)

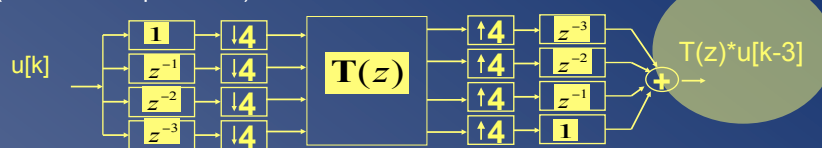
The synthesis FB is matched to the analysis bank, for PR: $\mathbf{R}(z) \cdot \mathbf{E}(z) = z^{-1} I_{4 \times 4}$

Frequency Domain FIR Filter Realization

- ☺ Overlap-save procedure is very efficient for large L :
 - Computational complexity (with FFT/IFFT i.o. DFT/IDFT) is $2 \cdot \lceil \alpha \cdot 2L \cdot \log(2L) \rceil + 2L$ multiplications for L output samples, i.e. $O(\log(L))$ per sample for large L
 - Compare to computational complexity for direct form realization: L multiplications per output sample, i.e. $O(L)$ per sample
- ☹ Overlap-save procedure introduces $O(L)$ processing delay/latency (e.g. $y[k-L+1]$ only available sometime after time k)
- Conclusion: For large L, complexity reduction is large, but latency is also large
- Will derive 'intermediate' realizations, with a smaller latency at the expense of a smaller complexity reduction. This will be based on an N^{th} order polyphase decomposition of $B(z)$...

Frequency Domain FIR Filter Realization

A compact derivation will rely on a result from filter bank theory (return to Chapter-10...)



(...and now let $B(z)$ take the place of 'distortion function' $T(z)$)

This means that a filter (specified with N^{th} order polyphase decomposition)

$$B(z) = p_0(z^4) + z^{-1}p_1(z^4) + z^{-2}p_2(z^4) + z^{-3}p_3(z^4)$$

can be realized in a multirate structure, based on a pseudo-circulant matrix

$$\mathbf{T}(z) = \begin{bmatrix} p_0(z) & p_1(z) & p_2(z) & p_3(z) \\ z^{-1} \cdot p_3(z) & p_0(z) & p_1(z) & p_2(z) \\ z^{-1} \cdot p_2(z) & z^{-1} \cdot p_3(z) & p_0(z) & p_1(z) \\ z^{-1} \cdot p_1(z) & z^{-1} \cdot p_2(z) & z^{-1} \cdot p_3(z) & p_0(z) \end{bmatrix}$$

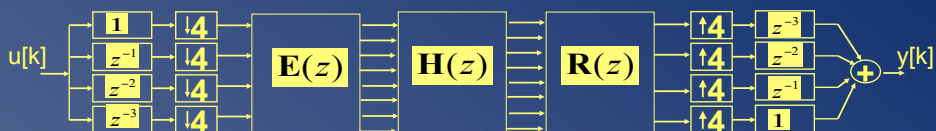
Frequency Domain FIR Filter Realization

Now some matrix manipulation... (compare to p.6)

$$\mathbf{T}(z) = \begin{bmatrix} 0 & I_{4 \times 4} \end{bmatrix} \cdot \begin{bmatrix} 0 & 0 & 0 & 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) \\ p_3(z) & 0 & 0 & 0 & 0 & p_0(z) & p_1(z) & p_2(z) \\ p_2(z) & p_3(z) & 0 & 0 & 0 & 0 & p_0(z) & p_1(z) \\ p_1(z) & p_2(z) & p_3(z) & 0 & 0 & 0 & 0 & p_0(z) \\ p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 & 0 & 0 & 0 \\ 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 & 0 & 0 \\ 0 & 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 & 0 \\ 0 & 0 & 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 \end{bmatrix} \cdot \begin{bmatrix} I_{4 \times 4} \\ z^{-1} I_{4 \times 4} \end{bmatrix} \\
 = \underbrace{\begin{bmatrix} 0 & I_{4 \times 4} \end{bmatrix}}_{\mathbf{R}(z)} \cdot F^{-1} \cdot \begin{bmatrix} P_0(z) & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & P_1(z) & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & P_2(z) & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & P_3(z) & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & P_4(z) & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & P_5(z) & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & P_6(z) & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & P_7(z) \end{bmatrix} \cdot F \cdot \begin{bmatrix} I_{4 \times 4} \\ z^{-1} I_{4 \times 4} \end{bmatrix} = \mathbf{E}(z)$$

Frequency Domain FIR Filter Realization

- An (8-channel) filter bank representation of this is...



Analysis bank:

$$\mathbf{E}(z) = F \cdot \begin{bmatrix} I_{4 \times 4} \\ z^{-1} I_{4 \times 4} \end{bmatrix}$$

Subband processing:

$$\mathbf{H}(z) = \text{diag}\{P_0(z), P_1(z), \dots, P_7(z)\}$$

Synthesis bank:

$$\mathbf{R}(z) = \begin{bmatrix} 0 & I_{4 \times 4} \end{bmatrix} F^{-1}$$

This is a 2N-channel filter bank, with N-fold downsampling

The analysis FB is a 2N-channel uniform DFT filter bank (see Chapter 11)

The synthesis FB is matched to the analysis bank, for PR: $\mathbf{R}(z) \cdot \mathbf{E}(z) = z^{-1} I_{4 \times 4}$

Frequency Domain FIR Filter Realization

- This is again known as an 'overlap-save' realization :

- Analysis bank: performs 2N-point DFT (FFT) of a block of (N=4) samples, together with the previous block of (N) samples (hence 'overlap')

$$\mathbf{E}(z) = F \cdot \begin{bmatrix} I_{4 \times 4} \\ z^{-1} \cdot I_{4 \times 4} \end{bmatrix}$$

← **block'**
← **'previous block'**

- Synthesis bank: performs 2N-point IDFT (IFFT), throws away the 1st half of the result, saves the 2nd half

(hence 'save')

$$\mathbf{R}(z) = \begin{bmatrix} 0 & I_{4 \times 4} \end{bmatrix} F^{-1}$$

← **'throw away'** ← **'save'**

- Subband processing corresponds to 'frequency domain' operation

- Complexity/latency? See p.16...

Frequency Domain FIR Filter Realization

Derivation on p.10 can also be modified as follows...

$$\mathbf{T}(z) = \begin{bmatrix} z^{-1} I_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot \begin{bmatrix} 0 & 0 & 0 & 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) \\ p_3(z) & 0 & 0 & 0 & 0 & p_0(z) & p_1(z) & p_2(z) \\ p_2(z) & p_3(z) & 0 & 0 & 0 & 0 & p_0(z) & p_1(z) \\ p_1(z) & p_2(z) & p_3(z) & 0 & 0 & 0 & 0 & p_0(z) \\ p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 & 0 & 0 & 0 \\ 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 & 0 & 0 \\ 0 & 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 & 0 \\ 0 & 0 & 0 & p_0(z) & p_1(z) & p_2(z) & p_3(z) & 0 \end{bmatrix} \begin{bmatrix} I_{4 \times 4} \\ 0_{4 \times 4} \end{bmatrix}$$

$$= \underbrace{\begin{bmatrix} z^{-1} I_{4 \times 4} & I_{4 \times 4} \end{bmatrix} F^{-1}}_{\mathbf{R}(z)} \cdot \begin{bmatrix} P_0(z) & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & P_1(z) & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & P_2(z) & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & P_3(z) & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & P_4(z) & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & P_5(z) & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & P_6(z) & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & P_7(z) \end{bmatrix} \cdot \underbrace{F \cdot \begin{bmatrix} I_{4 \times 4} \\ 0_{4 \times 4} \end{bmatrix}}_{\mathbf{E}(z)}$$

Frequency Domain FIR Filter Realization

- This is known as an **'overlap-add'** realization :
 - Analysis bank: performs 2N-point DFT (FFT) of a block of (N=4) samples, padded with N zero samples

$$\mathbf{E}(z) = F \cdot \begin{bmatrix} I_{4 \times 4} \\ 0_{4 \times 4} \end{bmatrix}$$

← 'block'
← 'zero padding'

- Synthesis bank: performs 2N-point IDFT (IFFT), adds 2nd half of the result to 1st half of previous IDFT (hence 'add')

$$\mathbf{R}(z) = \begin{bmatrix} z^{-1} \cdot I_{4 \times 4} & I_{4 \times 4} \end{bmatrix} F^{-1}$$

↑ 'overlap' ↑ 'add'

- Subband processing corresponds to 'frequency domain' operation

Frequency Domain FIR Filter Realization

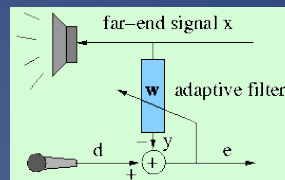
- Computational complexity** is (with FFT/IFFT i.o. DFT/IDFT, plus subband processing) $2 \cdot [\alpha \cdot 2N \cdot \log(2N)] + 2L$ multiplications for N output samples, i.e. $O(\log(N)) + O(L/N)$ per sample
 - ☺ For large $N \approx L$ this is $O(\log(L))$ i.e. dominated by FFT/IFFT (cheap!)
 - ☹ For $N \ll L$ this is $O(L)$, i.e. dominated by subband processing
- Processing delay/latency** is $O(N)$
- Standard 'overlap-add' and 'overlap-save' (=p.7) realizations are derived when **0th order** poly-phase components are used in the above derivation ($N=L$, i.e. each poly-phase component has only 1 filter coefficient). For large L, this leads to a large complexity reduction, but also a large latency ($=O(L)$)
- In the more general case, with **higher-order** polyphase components ($N < L$, i.e. each poly-phase component has >1 filter coefficients) a smaller complexity reduction is achieved, but the latency is also smaller ($=O(N)$).

Frequency Domain Adaptive Filtering

- A similar derivation can be made for LMS-based adaptive filtering with block processing ('Block-LMS'). The adaptive filter then consist in a filtering operation plus an adaptation operation, which corresponds to a correlation operation. Both operations can be performed cheaply in the frequency domain..
- Starting point is the LMS update equation

$$\mathbf{w}_{k+1} = \mathbf{w}_k + \mu \cdot \mathbf{x}_k \cdot (d[k] - \mathbf{x}_k^T \mathbf{w}_k)$$

$$\mathbf{x}_k = \begin{bmatrix} x[k] \\ \vdots \\ x[k-L+1] \end{bmatrix}, \quad \mathbf{w}_k = \begin{bmatrix} w[0] \\ \vdots \\ w[L-1] \end{bmatrix}$$



Frequency Domain Adaptive Filtering

Consider block processing with so-called '**Block-LMS**'

- Remember that LMS is a 'stochastic gradient' algorithm, where instantaneous estimates of the autocorrelation matrix and cross-correlation vector are used to compute a gradient (=steepest descent vector)
- Block-LMS uses averaged estimates, with averaging over a block of L_B ('block length') samples, and hence an averaged gradient.

The update formula is then..

where n is the block index

$$\mathbf{w}_{n+1} = \mathbf{w}_n + \mu \cdot \sum_{k=nL_B+1}^{nL_B+L_B} \mathbf{x}_k \cdot (d[k] - \mathbf{x}_k^T \mathbf{w}_n)$$

- Compared to LMS, Block-LMS does fewer updates (one per L_B samples), but with (presumably) better gradient estimates. Overall, convergence could be faster or slower (=unpredictable).
- The important thing is that Block-LMS can be realized cheaply...

Frequency Domain Adaptive Filtering

Will consider case where block length $L_B =$ filter length L
 The update formulas are then given as follows

1) Compute a priori residuals (example $L_B=L=4$, similar for other L)

$$\begin{bmatrix} e[4n+4] \\ e[4n+3] \\ e[4n+2] \\ e[4n+1] \end{bmatrix} = \begin{bmatrix} d[4n+4] \\ d[4n+3] \\ d[4n+2] \\ d[4n+1] \end{bmatrix} - \begin{bmatrix} w[0] & w[1] & w[2] & w[3] & 0 & 0 & 0 & 0 \\ 0 & w[0] & w[1] & w[2] & w[3] & 0 & 0 & 0 \\ 0 & 0 & w[0] & w[1] & w[2] & w[3] & 0 & 0 \\ 0 & 0 & 0 & w[0] & w[1] & w[2] & w[3] & 0 \end{bmatrix} \begin{bmatrix} x[4n+4] \\ x[4n+3] \\ x[4n+2] \\ x[4n+1] \\ x[4n] \\ x[4n-1] \\ x[4n-2] \\ x[4n-3] \end{bmatrix}$$

as on p.6

$$\begin{bmatrix} d[4n+4] \\ d[4n+3] \\ d[4n+2] \\ d[4n+1] \end{bmatrix} - \begin{bmatrix} 0_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot \mathcal{F}^{-1} \cdot \begin{bmatrix} W_0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & W_1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & W_2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & W_3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & W_4 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & W_5 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & W_6 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & W_7 \end{bmatrix} \cdot \mathcal{F} \cdot \begin{bmatrix} x[4n+4] \\ x[4n+3] \\ x[4n+2] \\ x[4n+1] \\ x[4n] \\ x[4n-1] \\ x[4n-2] \\ x[4n-3] \end{bmatrix}$$

with $W_i = \dots$

=frequency domain filtering

Frequency Domain Adaptive Filtering

Will consider case where block length $L_B =$ filter length L
 The update formulas are then given as follows

2) Filter update (example $L_B=L=4$, similar for other L)

$$\begin{bmatrix} w[0] \\ w[1] \\ w[2] \\ w[3] \end{bmatrix}_{n+1} = \begin{bmatrix} w[0] \\ w[1] \\ w[2] \\ w[3] \end{bmatrix}_n + \mu \cdot \begin{bmatrix} x[4n+4] & x[4n+3] & x[4n+2] & x[4n+1] \\ x[4n+3] & x[4n+2] & x[4n+1] & x[4n] \\ x[4n+2] & x[4n+1] & x[4n] & x[4n-1] \\ x[4n+1] & x[4n] & x[4n-1] & x[4n-2] \end{bmatrix} \begin{bmatrix} e[4n+4] \\ e[4n+3] \\ e[4n+2] \\ e[4n+1] \end{bmatrix}$$

$$= \begin{bmatrix} w[0] \\ w[1] \\ w[2] \\ w[3] \end{bmatrix}_n + \mu \cdot \begin{bmatrix} e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 & 0 & 0 & 0 \\ 0 & e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 & 0 & 0 \\ 0 & 0 & e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 & 0 \\ 0 & 0 & 0 & e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 \end{bmatrix} \begin{bmatrix} x[4n+4] \\ x[4n+3] \\ x[4n+2] \\ x[4n+1] \\ x[4n] \\ x[4n-1] \\ x[4n-2] \\ x[4n-3] \end{bmatrix}$$

as on p.6

$$\begin{bmatrix} w[0] \\ w[1] \\ w[2] \\ w[3] \end{bmatrix}_n + \mu \cdot \begin{bmatrix} 0_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot \mathcal{F}^{-1} \cdot \begin{bmatrix} E_0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & E_1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & E_2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & E_3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & E_4 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & E_5 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & E_6 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & E_7 \end{bmatrix} \cdot \mathcal{F} \cdot \begin{bmatrix} x[4n+4] \\ x[4n+3] \\ x[4n+2] \\ x[4n+1] \\ x[4n] \\ x[4n-1] \\ x[4n-2] \\ x[4n-3] \end{bmatrix}$$

with $E_i = \dots$

=frequency domain correlation

Frequency Domain Adaptive Filtering

This is referred to as **FDAF** ('Frequency Domain Adaptive Filtering')

- FDAF is functionally equivalent to Block-LMS (but cheaper, see below)
- **Convergence**: Instead of using one and the same stepsize μ for all 'frequency bins', frequency dependent stepsizes can be applied..
 - In the update formula, μ is removed and E_i is replaced by $\mu_i \cdot E_i$
 - Stepsize μ_i dependent on the energy in the i^{th} frequency bin
 - Leads to increased convergence speed at only a small extra cost

Frequency Domain Adaptive Filtering

This is referred to as **FDAF** ('Frequency Domain Adaptive Filtering')

- **Complexity** ≈ 5 (I)FFT's (size $2L$)
per block of L output samples (check!)
- Hence for large L , FDAF is very efficient/cheap, only $O(\log(L))$ multiplications per output sample (compared to $O(L)$ for (Block-)LMS)

Example: $L_B=L=1024$, then $\frac{\text{cost LMS}}{\text{cost FDAF}} \approx 20$

- **Processing delay/latency** is again $O(L)$.

Example: $L_B=L=1024$ and $f_s=8000\text{Hz}$, then delay is 256 ms !

In cases where this is objectionable (e.g. acoustic echo cancellation), need 'intermediate' algorithms with smaller latency and smaller complexity reduction, based on a polyphase decomposition...(read on)

Frequency Domain Adaptive Filtering

For large L , a block length of $L_B=L$ may lead to a too large latency

If an N^{th} order polyphase decomposition of the adaptive filter is considered (hence with $L_p=L/N$ coefficients per polyphase component), then a frequency domain adaptive filtering algorithm with block length $L_B=N$ can be derived as follows...

(where "N takes the place of L")

Example $L_B=N=4$, i.e. (as on p.9)

$$W(z) = p_0(z^4) + z^{-1}p_1(z^4) + z^{-2}p_2(z^4) + z^{-3}p_3(z^4)$$

with

$$p_0(z) = p_0^0 + p_0^1 z^{-1} + p_0^2 z^{-2} + \dots + p_0^{L_p+1} z^{-L_p+1}$$

$$p_1(z) = \text{etc.}$$

Frequency Domain Adaptive Filtering

(compare to p.19-20)

The update formulas are given as follows

1) Compute a priori residuals (example $L_B=N=4$, similar for other N)

$$\begin{bmatrix} e[4n+4] \\ e[4n+3] \\ e[4n+2] \\ e[4n+1] \end{bmatrix} = \begin{bmatrix} d[4n+4] \\ d[4n+3] \\ d[4n+2] \\ d[4n+1] \end{bmatrix} - \sum_{i=0}^{L_p} \begin{bmatrix} p_0^i & p_1^i & p_2^i & p_3^i & 0 & 0 & 0 & 0 \\ 0 & p_0^i & p_1^i & p_2^i & p_3^i & 0 & 0 & 0 \\ 0 & 0 & p_0^i & p_1^i & p_2^i & p_3^i & 0 & 0 \\ 0 & 0 & 0 & p_0^i & p_1^i & p_2^i & p_3^i & 0 \end{bmatrix} \begin{bmatrix} x[4(n-i)+4] \\ x[4(n-i)+3] \\ x[4(n-i)+2] \\ x[4(n-i)+1] \\ x[4(n-i)] \\ x[4(n-i)-1] \\ x[4(n-i)-2] \\ x[4(n-i)-3] \end{bmatrix}$$

as on p.6

$$\begin{bmatrix} d[4n+4] \\ d[4n+3] \\ d[4n+2] \\ d[4n+1] \end{bmatrix} = \begin{bmatrix} 0_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot \sum_{i=0}^{L_p} \begin{bmatrix} P_0^i & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & P_1^i & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & P_2^i & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & P_3^i & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & P_4^i & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & P_5^i & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & P_6^i & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & P_7^i \end{bmatrix} \cdot \begin{bmatrix} x[4(n-i)+4] \\ x[4(n-i)+3] \\ x[4(n-i)+2] \\ x[4(n-i)+1] \\ x[4(n-i)] \\ x[4(n-i)-1] \\ x[4(n-i)-2] \\ x[4(n-i)-3] \end{bmatrix}$$

with $P_i^j = \dots$

Frequency Domain Adaptive Filtering

(compare to p.18-19)

The update formulas are given as follows

2) Filter update (example $L_B=N=4$, similar for other N)

$$\begin{aligned}
 \begin{bmatrix} p_0(z^4) \\ p_1(z^4) \\ p_2(z^4) \\ p_3(z^4) \end{bmatrix}_{n+1} &= \begin{bmatrix} p_0(z^4) \\ p_1(z^4) \\ p_2(z^4) \\ p_3(z^4) \end{bmatrix}_n + \mu \cdot \sum_{i=0}^{L_B} z^{-4i} \begin{bmatrix} x[4(n-i)+4] & x[4(n-i)+3] & x[4(n-i)+2] & x[4(n-i)+1] \\ x[4(n-i)+3] & x[4(n-i)+2] & x[4(n-i)+1] & x[4(n-i)] \\ x[4(n-i)+2] & x[4(n-i)+1] & x[4(n-i)] & x[4(n-i)-1] \\ x[4(n-i)+1] & x[4(n-i)] & x[4(n-i)-1] & x[4(n-i)-2] \end{bmatrix} \begin{bmatrix} e[4n+4] \\ e[4n+3] \\ e[4n+2] \\ e[4n+1] \end{bmatrix} \\
 &= \begin{bmatrix} p_0(z^4) \\ p_1(z^4) \\ p_2(z^4) \\ p_3(z^4) \end{bmatrix}_n + \mu \cdot \begin{bmatrix} e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 & 0 & 0 & 0 \\ 0 & e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 & 0 & 0 \\ 0 & 0 & e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 & 0 \\ 0 & 0 & 0 & e[4n+4] & e[4n+3] & e[4n+2] & e[4n+1] & 0 \end{bmatrix} \sum_{i=0}^{L_B} z^{-4i} \begin{bmatrix} x[4(n-i)+4] \\ x[4(n-i)+3] \\ x[4(n-i)+2] \\ x[4(n-i)+1] \\ x[4(n-i)] \\ x[4(n-i)-1] \\ x[4(n-i)-2] \\ x[4(n-i)-3] \end{bmatrix} \\
 \stackrel{\text{as on p.6}}{=} &\begin{bmatrix} p_0(z^4) \\ p_1(z^4) \\ p_2(z^4) \\ p_3(z^4) \end{bmatrix}_n + \mu \cdot \begin{bmatrix} 0_{4 \times 4} & I_{4 \times 4} \end{bmatrix} \cdot F^{-1} \cdot \begin{bmatrix} E_0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & E_1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & E_2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & E_3 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & E_4 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & E_5 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & E_6 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & E_7 \end{bmatrix} \cdot F \cdot \sum_{i=0}^{L_B} z^{-4i} \begin{bmatrix} x[4(n-i)+4] \\ x[4(n-i)+3] \\ x[4(n-i)+2] \\ x[4(n-i)+1] \\ x[4(n-i)] \\ x[4(n-i)-1] \\ x[4(n-i)-2] \\ x[4(n-i)-3] \end{bmatrix}
 \end{aligned}$$

Frequency Domain Adaptive Filtering

This is referred to as **PB-FDAF**

(‘Partitioned Block Frequency Domain Adaptive Filtering’)

- PB-FDAF is functionally equivalent to Block-LMS
- **Complexity** $\approx 3+2 \cdot L_P$ (I)FFT’s (size 2N))
per block of N output samples (check!)

Example: L=1024, N=128, then

$$\frac{\text{cost LMS}}{\text{cost PB-FDAF}} \approx 6$$

- **Processing delay/latency** is $O(N)$.

Example: L=1024, N=128 and $f_s=8000\text{Hz}$, then delay is 32 ms
(used in commercial acoustic echo cancellers)